# Probabilistic Methods <br> Part I. Lovász Local Lemma 

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ALEA Days, CIRM Marseille
13 March 2024

## Introduction

Probabilistic methods

- prove the existence of combinatorial objects
- using probabilistic tools and arguments
- First moment principles: linearity of expectation
- Second moment inequalities
- Lovász Local Lemma
- Entropy Compression
- Concentration inequalities


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## Outline of today's talk

- a warmup example
- hypergraph coloring problem
- statement of the Lovász Local Lemma
- application in hypergraph coloring
- application in acyclic graph coloring


## Warmup example

Given a graph on $n$ vertices and $m$ edges, what minimum size of a bipartite (spanning) subgraph can be guaranteed?

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The best we can hope for is $\sim \frac{m}{2}$ :

- a complete graph on $n$ vertices has $\binom{n}{2} \sim \frac{n^{2}}{2}$ edges
- a complete bipartite graph on $\left\lceil\frac{n}{2}\right\rceil+\left\lfloor\frac{n}{2}\right\rfloor$ vertices has $\sim \frac{n^{2}}{4}$ edges


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Randomized procedure

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- Remove monochromatic edges


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Then $\mathbb{E}\left(X_{e}\right)=\frac{1}{2}$, and by linearity of expectation,
$\mathbb{E}\left(\sum_{e \in E(G)} X_{e}\right)=\sum_{e \in E(G)} \mathbb{E}\left(X_{e}\right)=\frac{m}{2}$.

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$\mathbb{E}\left(\sum_{e \in E(G)} X_{e}\right)=\sum_{e \in E(G)} \mathbb{E}\left(X_{e}\right)=\frac{m}{2}$.
Therefore, there exists a coloring with at least $\frac{m}{2}$ bichromatic edges.

## Hypergraph coloring

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A hypergraph is $\underline{k \text {-regular }}$ if $|\{e \in E: v \in e\}|=k \forall v \in V$.

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If $\mathscr{A}:=\left(A_{e}, e \in E\right)$ were independent, we would have

$$
\mathbb{P}\left(\bigcap_{e \in E} \overline{A_{e}}\right)=\left(1-\frac{1}{2^{k-1}}\right)^{m}>0
$$

## Mutually independent events

## Definition

Let $A$ be an event and let $\mathscr{B}$ be a set of events in a probability space. We say that $A$ is mutually independent of $\mathscr{B}$ if

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for every set $S \subseteq \mathscr{B}$.
For example, in the context of random hypergraph coloring, $A_{e}$ is mutually independent of

$$
\left\{A_{e^{\prime}}: e \cap e^{\prime}=\emptyset\right\} .
$$

## Lovász Local Lemma

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Theorem (Lovász Local Lemma, Symmetric version)
Let $\mathscr{A}=\left\{A_{1}, A_{2}, \ldots, A_{n}\right\}$ be a set of events such that for each $i=1,2 \ldots, n$

- $\mathbb{P}\left(A_{i}\right) \leq p \quad$ and
- $\exists \mathscr{D}_{i} \subset \mathscr{A}$ of size at most $d$ such that $A_{i}$ is mutually independent of $\mathscr{A} \backslash \mathscr{D}_{i}$.

If

$$
e \cdot p \cdot(d+1) \leq 1
$$

then

$$
\mathbb{P}\left(\bigcap_{i=1}^{n} \overline{A_{i}}\right)>0 .
$$

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Theorem (LLL)
If $\mathbb{P}\left(A_{i}\right) \leq p, A_{i}$ is mutually independent of $\mathscr{A} \backslash \mathscr{D}_{i}$ with
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In the context of random coloring of a $k$-regular $k$-uniform hypergraph, $p=\frac{1}{2^{k-1}}$

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In the context of random coloring of a $k$-regular $k$-uniform hypergraph, $p=\frac{1}{2^{k-1}}$ and each $A_{e}$ is mutually independent of all but at most $k^{2}$ other edges, so $d=k^{2}$.

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There exists a coloring without a monochromatic edge whenever

$$
\frac{e}{2^{k-1}} \cdot k^{2} \leq 1
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Let $k \geq 9$. Then every $k$-regular $k$-uniform hypergraph is 2-colorable.

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Theorem (Alon and Bregman 1988, Henning and Yeo 2013)
Let $k \geq 4$. Then every $k$-regular $k$-uniform hypergraph is 2 -colorable.

## Acyclic graph coloring

Definition
Let $G=(V, E)$ be a graph. A coloring
$\varphi: V(G) \rightarrow\{1,2, \ldots, k\}$ is an acyclic coloring of $G$ if

- $\varphi(u) \neq \varphi(v) \quad \forall u v \in E(G), \quad$ ( $\varphi$ is a proper coloring)
- there is no bichromatic cycle in $G$.


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an acyclic coloring with 4 colors


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## Greedy bound

Can we bound $\chi_{a}(G)$ as a function of $\Delta(G)$, the maximum degree of $G$ ?

If we color every vertex with a color distinct from all the colors of its neighbors and the neighbors of its neighbors, surely we will not create any bichromatic cycle.

This is always possible provided we have at least

$$
\Delta+\Delta(\Delta-1)+1=\Delta^{2}+1
$$

colors. Hence,

$$
\chi_{a}(G) \leq \Delta^{2}+1
$$

for every graph $G$.

## Using Lovász Local Lemma

Can we bound $\chi_{a}(G)$ as a function of $\Delta(G)$, the maximum degree of $G$ ?

Theorem (Alon, McDiarmid, Reed 1991)
Let $G$ be a graph with maximum degree $\Delta$. Then

$$
\chi_{a}(G) \leq 50 \Delta^{4 / 3} .
$$

On the other hand, there are graphs for which

$$
\chi_{a}(G)=\Omega\left(\frac{\Delta^{4 / 3}}{(\log \Delta)^{1 / 3}}\right) .
$$

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Can we bound $\chi_{a}(G)$ as a function of $\Delta(G)$, the maximum degree of $G$ ?
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Let $C$ be a set of $K \geq 7 \Delta^{3 / 2}$ colors.
Randomized procedure : For each vertex $v$, let $F(v)$ be the set of colors forbidden at $v$ - the colors of the neighbors already colored, and let $C(v)=C \backslash F(v)$ be the set of available colors at $v$. Clearly, $|F(v)| \leq \Delta$.

- Choose an integer $i \leq K-\Delta$ uniformly randomly and color $v$ with $i$-th available color.
This procedure gives a proper coloring of $G$.


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Can we bound $\chi_{a}(G)$ as a function of $\Delta(G)$, the maximum degree of $G$ ?
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- Choose an integer $i \leq K-\Delta$ uniformly randomly and color $v$ with $i$-th available color.

Let $A_{P}$ be the event that a 4-vertex path $P=v_{1} v_{2} v_{3} v_{4}$ gets only two colors.

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\mathbb{P}\left(A_{P}\right) \leq \frac{1}{(K-\Delta)^{2}}
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The dependency degree is (less than)

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d<4 \cdot 4 \cdot \Delta^{3}=16 \Delta^{3}
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and so

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e p(d+1) \leq \frac{e\left(16 \Delta^{3}\right)}{\left(7 \Delta^{3 / 2}-\Delta\right)^{2}}<\frac{0.89 \Delta}{\left(\Delta^{1 / 2}-\frac{1}{7}\right)^{2}}<1
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## Conclusion

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Thank you for your attention!

