

Practical Encryption from Isogenies between Elliptic Products

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SIDH Attacks in a Nutshell

Given $(\varphi(P), \varphi(Q))$ under a secret isogeny $\varphi: E \rightarrow E'$ of degree d , where $\langle P, Q \rangle = E[n]$, it is possible to recover φ .

Trapdoor Function

Triple of algorithms $(\text{KeyGen}, f, f^{-1})$

- $\text{KeyGen}(\lambda) \xrightarrow{\$} (\text{sk}, \text{pk})$
- $f(\text{pk}, x) \rightarrow y$
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For all pk and x , $f^{-1}(\text{sk}, f(\text{pk}, x)) = x$.

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One-way

Given pk and y , finding x st $f(\text{pk}, x) = y$ is hard.

FESTA Trapdoor

$$\begin{pmatrix} P_b \\ Q_b \end{pmatrix} E_0$$

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$$\begin{pmatrix} P_b \\ Q_b \end{pmatrix} E_0 \xrightarrow{\varphi_A} E_A$$

$$\text{sk} = (\mathbf{A}, \varphi_A)$$

FESTA Trapdoor

$$\begin{array}{ccc} \begin{pmatrix} P_b \\ Q_b \end{pmatrix} & & \begin{pmatrix} R_A \\ S_A \end{pmatrix} = \mathbf{A} \begin{pmatrix} \varphi_A(P_b) \\ \varphi_A(Q_b) \end{pmatrix} \\ E_0 & \xrightarrow{\varphi_A} & E_A \end{array}$$

$$\begin{pmatrix} \alpha & \beta \\ \gamma & \delta \end{pmatrix} \begin{pmatrix} P_1 \\ P_2 \end{pmatrix} = \begin{pmatrix} [\alpha]P_1 + [\beta]P_2 \\ [\gamma]P_1 + [\delta]P_2 \end{pmatrix}$$

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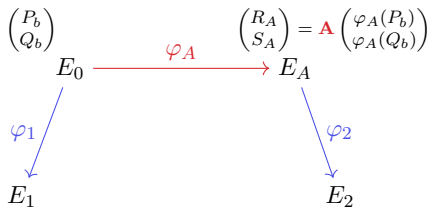
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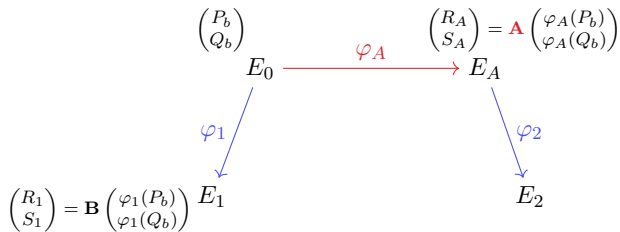
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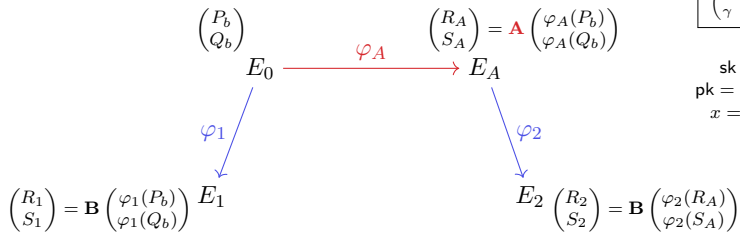
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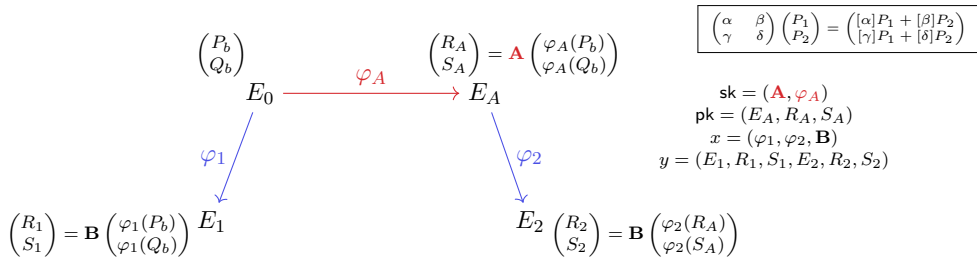
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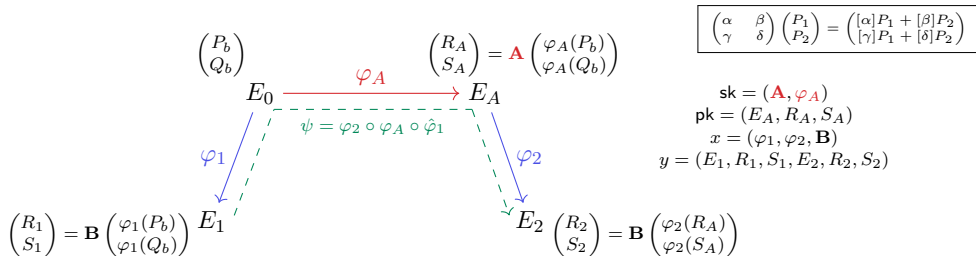
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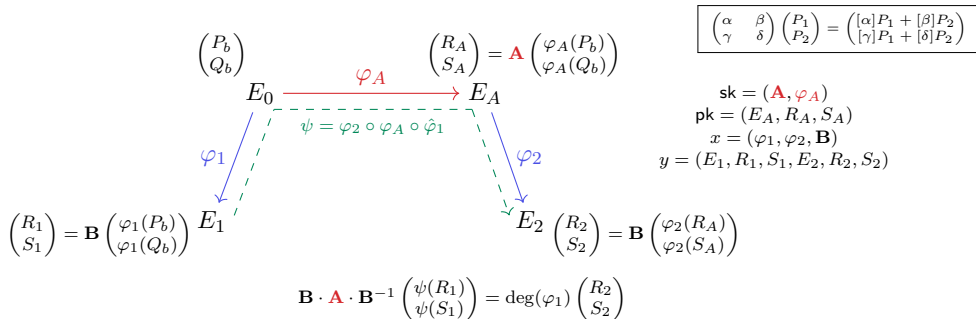
FESTA Trapdoor



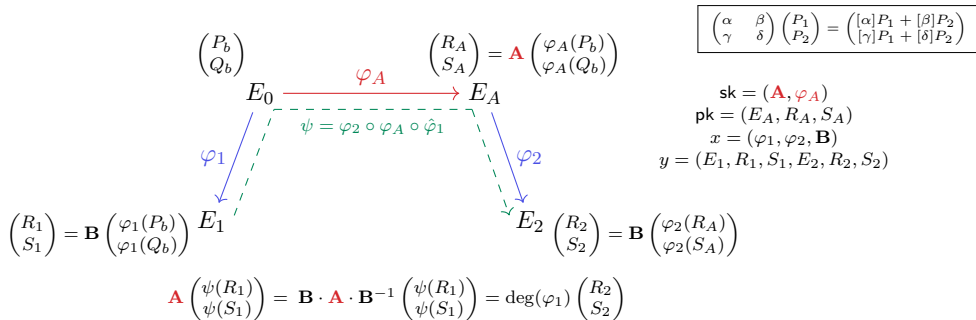
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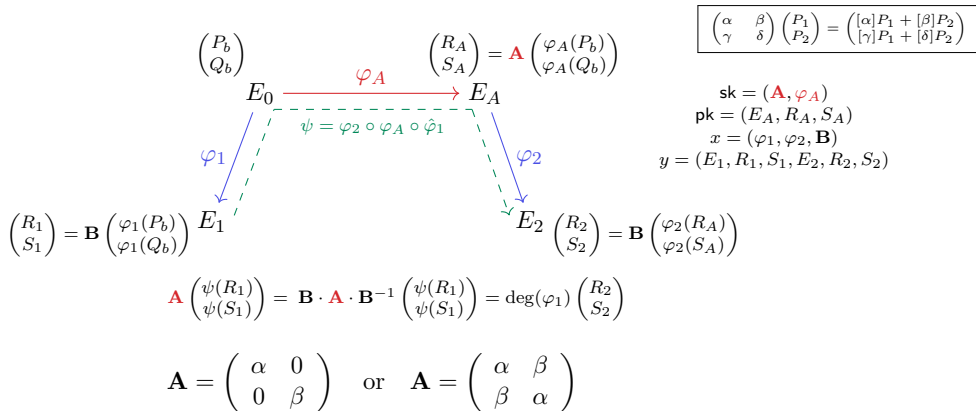
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Kani's Criterion

- $\varphi_{N_1}: E_0 \rightarrow E_1, \varphi_{N_2}: E_0 \rightarrow E_2$ s.t. $\gcd(N_1, N_2) = 1$
- $K := \langle ([N_2]\varphi_{N_1}(P), [N_1]\varphi_{N_2}(P)), ([N_2]\varphi_{N_1}(Q), [N_1]\varphi_{N_2}(Q)) \rangle$, where $\langle P, Q \rangle = E_0[N_1 + N_2]$

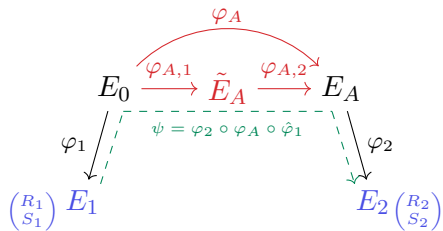
Theorem

The $(N_1 + N_2, N_1 + N_2)$ -polarised isogeny Φ with kernel K has matrix form

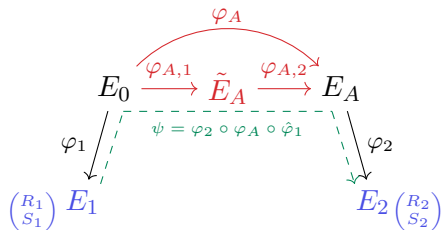
$$\begin{pmatrix} \hat{\varphi}_{N_1} & -\hat{\varphi}_{N_2} \\ g_{N_2} & \hat{g}_{N_1} \end{pmatrix},$$

where g_{N_i} are N_i -isogenies such that $\varphi_{N_2} \circ \hat{\varphi}_{N_1} = g_{N_1} \circ g_{N_2}$.

Concrete Inversion

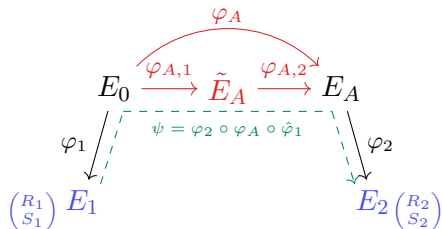


Concrete Inversion



- $d_{A,i} = \deg(\varphi_{A,i})$, $d_i = \deg(\varphi_i)$
- $m_1^2 d_{A,1} d_1 + m_2^2 d_{A,2} d_2 = 2^b$
- $\varphi_{N_1} = [m_1] \circ \varphi_1 \circ \widehat{\varphi}_{A,1}$,
 $\varphi_{N_2} = [m_2] \circ \varphi_2 \circ \varphi_{A,2}$
- $d_i > 2^{2\lambda}$, $d_{A,1} \cdot d_{A,2} > 2^{2\lambda}$

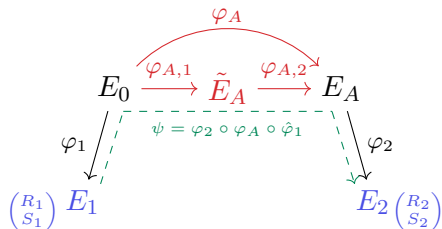
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$$K := \left\langle \begin{array}{l} ([m_2 d_{A,2} d_2] R_1, [d_1 m_1] R'_2), \\ ([m_2 d_{A,2} d_2] S_1, [d_1 m_1] S'_2) \end{array} \right\rangle, \quad \begin{pmatrix} R'_2 \\ S'_2 \end{pmatrix} = \mathbf{A}^{-1} \begin{pmatrix} R_2 \\ S_2 \end{pmatrix}$$

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$$\Phi = \begin{pmatrix} [m_1] \circ \varphi_{A,1} \circ \widehat{\varphi}_1 & -[m_2] \circ \widehat{\varphi}_{A,2} \circ \widehat{\varphi}_2 \\ [m_2] \circ g_{d_{A,2}} & [m_1] \circ \widehat{g}_{d_{A,1} d_1} \end{pmatrix}$$

<https://github.com/FESTA-PKE/FESTA-SageMath>

```
=====  
Running FESTA_128  
=====
```

```
=====  
Keygen took: 4.853 seconds  
=====
```

```
-----  
Compressed public key: 561 bytes  
-----
```

```
=====  
Encrypt took: 3.513 seconds  
=====
```

```
-----  
Compressed ciphertext: 1122 bytes  
-----
```

```
=====  
Decrypt took: 10.102 seconds  
=====
```

Thanks for your
attention!



Questions?