

# ALGEBRAIC GEOMETRY CODES IN THE SUM-RANK METRIC

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## Linear codes and codes in the Hamming metric

$k$  a field (think about  $k = \mathbb{F}_q$ ),  $\mathcal{H}$  a  $k$ -linear vector space endowed with a metric

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**Reed–Solomon (RS) codes:**



The diagram shows a horizontal line representing the finite field  $\mathbb{F}_q$ . On this line, there are points labeled  $x_1$ ,  $x_2$ , an ellipsis  $\dots$ , and  $x_n$ . Arrows point from each of these points down to the corresponding element in the tuple  $(P(x_1), P(x_2), \dots, P(x_n))$  within the set definition of  $RS_\delta(\mathbf{x})$ .

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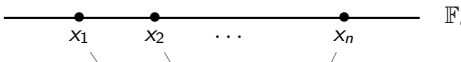
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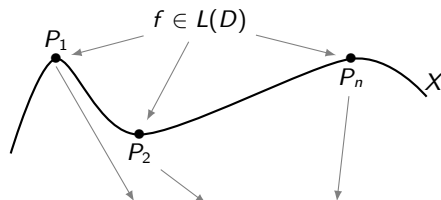
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✓ **Good parameters:**  $n + 1 - g \leq \delta + d \leq n + 1$

✓ **Longer codes**

## General definitions

$\underline{V} = (V_1, \dots, V_s)$  s-uple of  $k$ -vector spaces

$(n_i = \dim_k V_i)$

$$\mathcal{H} = \text{End}_k(\underline{V}) \quad := \quad \text{End}_k(V_1) \times \cdots \times \text{End}_k(V_s)$$

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$n_i = 1 \ \forall i \quad \rightsquigarrow \quad$  codes of length  $s$  in the **Hamming metric**

## Particular case and Singleton bound

$\ell$  = finite extension of  $k$  of degree  $r$

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### Singleton bound

The  $\ell$ -parameters of  $\mathcal{C}$  satisfy

$$d + \delta_\ell \leq n_\ell + 1.$$

Codes with parameters attaining this bound are called **Maximum Sum-Rank Distance (MSRD)**.

## Ore polynomials and Linearized Reed–Solomon codes

$\ell$  field,  $\Phi : \ell \rightarrow \ell$  a ring homomorphism,  $\ell^{\Phi=1} = k$ ,  $[\ell : k] = r$ ,

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As in the Hamming case, we can **try to overcome the problem using algebraic curves**

**Main idea:** consider Ore polynomials with coefficients in the function field of a curve



# Divisors and Riemann–Roch spaces: classical theory

## Definition

Let  $X$  be a nice curve,  $K$  its function field. A *divisor* on  $X$  is a formal finite sum

$$D = \sum_{\mathfrak{p} \in X} n_{\mathfrak{p}} \mathfrak{p} \quad \text{with } n_{\mathfrak{p}} \in \mathbb{Z} \text{ almost all zero.}$$

The group of divisors on  $X$  is denoted by  $\text{Div}(X)$ .

$D \in \text{Div}(X)$  is *positive*,  $D \geq 0$ , if  $n_{\mathfrak{p}} \geq 0 \forall \mathfrak{p}$ . The *degree* of  $D$  is  $\deg_X(D) = \sum_{\mathfrak{p} \in X} n_{\mathfrak{p}} \deg_X(\mathfrak{p})$ .

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## Riemann–Roch theorem

Let  $K_X$  denotes a canonical divisor on  $X$ . For any divisor  $D \in \text{Div}(X)$  we have

$$\dim_k L_X(D) = \deg_X(D) + 1 - g_X + \dim_k L_X(K_X - D),$$

$= 0$  when  $\deg_X(D) > 2g_X - 2$ .

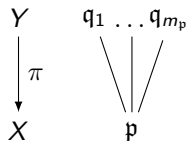
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$\pi$  a Galois cover with cyclic Galois group of order  $r$

$L := k(Y)$  the fields of functions of  $Y$ ,  $\text{Gal}(L/K) = \langle \Phi \rangle$

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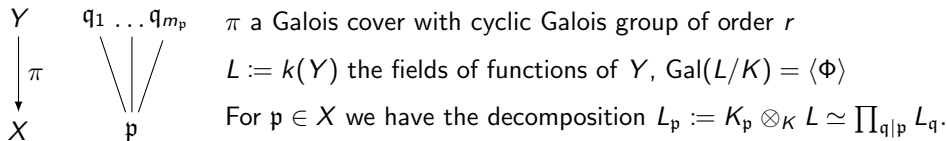


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$L := k(Y)$  the fields of functions of  $Y$ ,  $\text{Gal}(L/K) = \langle \Phi \rangle$

For  $\mathfrak{p} \in X$  we have the decomposition  $L_{\mathfrak{p}} := K_{\mathfrak{p}} \otimes_K L \simeq \prod_{\mathfrak{q}|\mathfrak{p}} L_{\mathfrak{q}}$ .

## Our setting

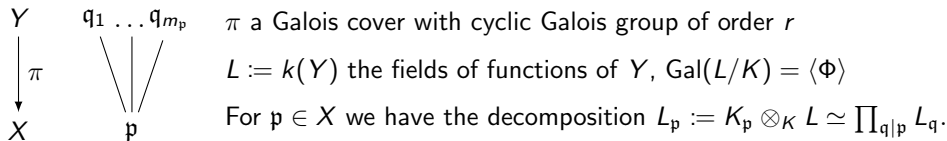


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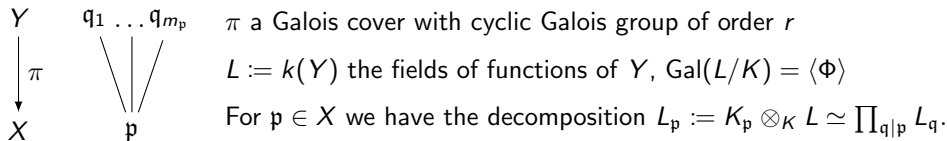
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- ☐ principal divisors associated to  $f \in D_{L,x}$   
 $\rightsquigarrow$  need to define a valuation
- ☐ Riemann–Roch spaces of  $D_{L,x}$
- ☐ a Riemann–Roch theorem
- ☐ equivalent of “evaluate at a rational point”



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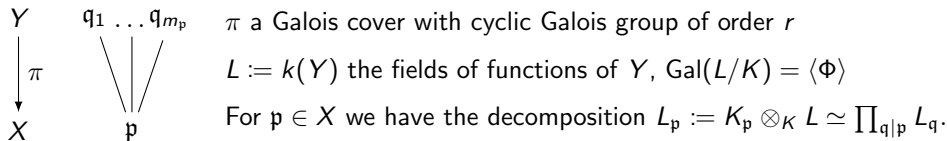
Define the **valuation map**  $w_{q,j,x} : D_{L_p,x} \rightarrow \frac{1}{r}\mathbb{Z} \sqcup \{\infty\}_{(1 \leq j \leq m_p)}$ : for  $f = f_0 + f_1 T + \dots + f_{r-1} T^{r-1}$ ,

$$w_{q,x}(f) = \min_{0 \leq i < r} \left( \frac{v_q(f_i)}{e_q} + i \cdot \frac{v_p(x)}{r} \right),$$

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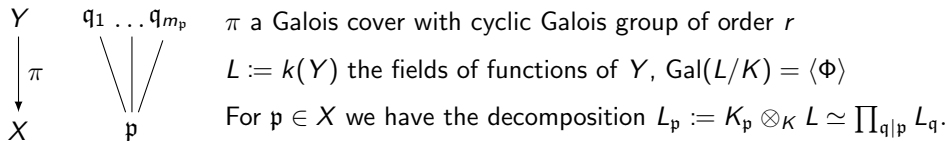
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For  $\mathfrak{p} \in X$ ,  $e_{\mathfrak{p}} w_{\mathfrak{q},x}(f) \in \frac{1}{b_{\mathfrak{p}}}\mathbb{Z}$  where  $b_{\mathfrak{p}}$  is the denominator of  $\rho_{\mathfrak{p}} = \frac{e_{\mathfrak{p}} \cdot v_{\mathfrak{p}}(x)}{r}$  after reduction

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## Riemann's inequality for $\Lambda_{L,X}(E)$

For a divisor  $E = \sum_{q \in Y} n_q q \in \text{Div}_{\mathbb{Q}}(Y)$  the space  $\Lambda_{L,X}(E)$  is finite dimensional over  $k$  and

$$\dim_k \Lambda_{L,X}(E) \geq r \cdot \deg_Y(E) - r \cdot (g_Y - 1) - \frac{r^2}{2} \sum_{p \in X} \frac{b_p - 1}{b_p e_p} \deg_X(p).$$

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## Linearized Algebraic Geometry codes

(B., Caruso, 2023)

Let  $E = \sum_{q \in Y} n_q q \in \text{Div}_{\mathbb{Q}}(Y)$ . Chose  $x \in K$  and  $p_1, \dots, p_s$  rational places on  $X$  such that the hypotheses hold. Consider

$$\begin{array}{ccc} \alpha : & \Lambda_{L, x}(E) & \longrightarrow \prod_{i=1}^s \text{End}_k(V_{p_i}) \\ & f & \mapsto (\bar{\varepsilon}_{p_i}(f))_{1 \leq i \leq s}. \end{array}$$

The code  $\mathcal{C}(x; E; p_1, \dots, p_s)$  is defined as the image of  $\alpha$ .

## Code's parameters

We study the parameters of the  $k$ -linear code  $\mathcal{C}$  in  $\prod_{i=1}^s \text{End}_k(V_{\mathfrak{p}_i})$ .

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Assume  $\deg_Y(E) < sr$ . Assume the **previous hypotheses** and that  $D_{L,X}$  contains no nonzero **divisors**. Then, the **dimension**  $\delta$  and the **minimum distance**  $d$  of  $\mathcal{C}(X; E; \mathfrak{p}_1, \dots, \mathfrak{p}_s)$  satisfy

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**Singleton bound:**

$$rd + \delta \leq n + r$$

**We have:**

$$rd + \delta \geq n + r - \left( r \cdot g_Y + \frac{r^2}{2} \sum_{p \in X} \frac{b_p - 1}{b_p e_p} \deg_X(p) \right)$$

## The case of isotrivial covers

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$X = \mathbb{P}_k^1$ ,  $Y = \mathbb{P}_\ell^1$ ,  $E = \frac{\delta}{r} \cdot \infty \in \operatorname{Div}_{\mathbb{Q}}(Y) \rightsquigarrow$  linearized Reed–Solomon codes!

Our lower bounds  $\Rightarrow$  MSRD codes


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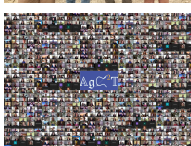
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AGC<sup>2</sup>T '17AGC<sup>2</sup>T '19AGC<sup>2</sup>T '21AGC<sup>2</sup>T '23

Merci de votre attention !

Questions?

[elena.berardini@math.u-bordeaux.fr](mailto:elena.berardini@math.u-bordeaux.fr)



## Remarks on the hypotheses

(H1) the algebra  $D_{L,x}$  has no nonzero zero divisor

(H2) for all places  $q$  above  $p$ , there exists  $u_q \in L_q^\times$  such that  $v_q(u_q) = \frac{e_p}{r} \cdot v_p(x)$  and

$$x = \prod_{q|p} N_{L_q/K_p}(u_q)$$

### Lemma

*The hypothesis (H1) holds as soon as there exists a place  $p \in X$  which is inert in  $Y$  and at which  $v_p(x)$  is coprime with  $r$ .*

### Lemma

*We assume that  $k$  is a finite field. Let  $p$  be a place of  $X$ . If  $p$  is unramified in  $Y$  and  $v_p(x)$  is divisible by  $r$ , then (H2) holds.*