

Subtitles for

DEUX OU TROIS
CHOSSES QUE JE
SAIS D'ELLE

(LA LOGIQUE)

JEAN-YVES GIRARD

I — THE END OF SYSTEMS

	Analytic	Synthetic
Explicit	Constat	Usine
Implicit	Performance	Usage

LL : central role of usine (proof-nets).

	Type – free	Typed
Explicit	Data	Cut – free
Implicit	Programs	Deductive

In technical terms.

1 – BHK, HERBRAND, GENTZEN

- Logical system : *oxymoron*, logic should not be sectarian.
LL : not yet another system, since *system-free!*
- *Gentzen* subformula property.
Proof of A : involves the sole A , not full system.
- *Herbrand's* theorem, prefiguration of *proof-nets* (usine).
 $\forall x \exists y A[x, y] \neq \exists y \forall x A[x, y]$: explained by *tests*.
 $A[x, t]$: make $x =: f(y)$ in case $\exists \forall$.
- *BHK*, proofs as functions, meaning is *use* (usage), not checkable (usine).
Proof of $A \Rightarrow B$: function from proofs of A to proofs of B .
Performative : from constat to constat (cut-elimination).
- *Open* architecture ; independent parts compatible.
Usage : Curry-like (existentialist, with cut).
Usine : Church-like (not quite essentialist, cut-free).
Adequation : external, proof of normalisation.

2 – EQUALITY

- « **Leibniz equality** » : any property of a is a property of b .

$$a = b := \forall X (X(a) \multimap X(b))$$

- **Quid** of « left of = » ?

Shocking : not *semantic*, does not preserve... equality!

Systems : avoid aporia by selecting properties *compatible* with equality.

- Transitivity $\vdash X(a) \otimes \sim X(b), X(b) \otimes \sim X(c), \sim X(a) \wp X(c)$

Semantic mistake : remove all $X(\cdot)$ and everything works !

Rewrites as : $\vdash a \otimes \sim b, b \otimes \sim c, \sim a \wp c$

- **Individuals as propositions** suppose *linear* maintenance.

Intuitionistic : $\neg\neg((a \equiv b) \vee (b \equiv c) \vee (c \equiv a))$

Equality : $a \multimap b \ \& \ b \multimap a : X(x) := x, X(x) := \sim x$.

- Leibniz : logical *epicycle*.

Metaphysics : individuals and their « **properties** » sort of viruses.

- Non linearity : need system to avoid *miscegenation*.

3 – ADDITIVES

- Second order $A \oplus B := \forall X ((A \multimap X) \Rightarrow ((B \multimap X) \Rightarrow X))$

$$\begin{array}{c} \vdots \\ \vdots \\ \vdots \\ A \end{array} \frac{}{A \oplus B} (l \oplus I) \quad \begin{array}{c} \vdots \\ \vdots \\ \vdots \\ B \end{array} \frac{}{A \oplus B} (r \oplus I) \quad \begin{array}{c} [A] \\ \vdots \\ \vdots \\ A \oplus B \end{array} \quad \begin{array}{c} [B] \\ \vdots \\ \vdots \\ C' \end{array} \quad \begin{array}{c} \vdots \\ \vdots \\ \vdots \\ C'' \end{array} \frac{}{C} (\oplus E)$$

- No subformula property, *extraneous* C may hide a cut.
Prawitz : *commutative* conversions against the performative paradigm.
- LL solution (new !) : delogicalise C', C'', C , into plain *locations* a, b, c .
Switching L/R : connect c with a or b .
Hidden cuts : normalised during « compilation » (correctness).
Extra locations : a, b, c part of formula $A \oplus B$.
Rule $(l \oplus I)$: proof of A and link $\llbracket a, c \rrbracket$
Rule $(\&)$: proof of A, a , proof of B, b and link(s) $\llbracket c, \Gamma \rrbracket$

4 – EMPTY TYPES

- Nightmare of *realisability* (cheap BHK), e.g., $\neg A \vee \neg\neg A$.
Invisible proofs : excluded since *false* but logically correct.
Example : switchings of $\sim A$.
Question : difference between *legit proof* and plain counter-switching.
- Second order : synthetic (subjective) features due to *existence*.
Proof of $\exists X A$: involves switchings of the witness.
- Euler-Poincaré invariant for $\llbracket t_1, \dots, t_k, u_1, \dots, u_\ell \rrbracket$:
Objective : if $\ell = 0$, then $|\llbracket t_1, \dots, t_k \rrbracket| = 2 - k$.
Subjective : if $\ell \neq 0$, then $|\llbracket t_1, \dots, t_k, u_1, \dots, u_\ell \rrbracket| = -k$.
Visibility : \mathcal{P} *true* when $|\mathcal{P}| \geq 0$.
- Truth preserved by rules, *including cut*, and normalisation.
Derealistic : applies to proofs, not to formulas ; no semantics !
Hidden proofs : essential to architecture like hidden files (*.emacs*).
Absurdity : $0 := !(\neg \wp \neg) \otimes \neg$. Not empty.

5 – ARITHMETIC

- Multiplicative constants $1, \perp$ essentialist fantasy of *category theory*.

Don't work : no correctness criterion, hence *impossible*.

Actual constants : \neg, \exists (katakanas *fu* et *wo*).

Atoms \neg objective, \exists subjective. Self-dual $\neg = \sim\neg, \exists = \sim\exists$.

Absurdity : $0 := !(\neg \exists \exists) \otimes \exists$.

- Desaxiomatisation of Peano *arithmetic*.

Second order : multiplication and recurrence à la *Dedekind*

$$x \in \mathbb{N} := \forall X (\forall y (X(y) \multimap X(y + \bar{1})) \Rightarrow (X(\bar{0}) \multimap X(x)))$$

Individuals : $\bar{1} := \neg, x + y := x \otimes y, x - y := y \multimap x$.

Axioms 3 and 4 : e.g., $x + \bar{1} \equiv y + \bar{1} \multimap x \equiv y$ prouvables.

- Works thanks to *derealistic* truth.

True and false : $\bar{1} := \neg$ true, $\overline{\bar{1}} := \neg \exists \neg \exists \neg$ false...

But : $\bar{1} + \overline{\bar{1}} \equiv \bar{0} (\equiv \neg \exists \neg)$ true.

$\bar{1} \otimes \overline{\bar{1}}$ true : contradicts tarskian pleonasm and *any* truth table.

II — LOGIC AND OPERATORS

The actual analytics.

6 – Gol

- Analytic : beyond discussion (*computation*). Quid of analytic *space* ?
 Pure λ -calculus robust and *archaic* : no linearity.
 Operator algebras : C^* , von Neumann.
- Upper part of *multiplicative* proof-net seen as :
 Permutation : leads to *Geometry of Interaction*.
 Partition : adapted to « Schrödinger's cut. »
- Gol : cut-free proof as *partial symmetry*, $u^3 = u, u = u^*$
 Application of « function » F to « argument » a (cut-elimination).
 $[F]a := (1 - a^2)(1 - Fa)^{-1}F(1 - a^2)$
 Provided $1 - Fa$ invertible.
- The ultimate BHK. But takes care of the sole *usage* (proofs as functions).
 Termination : Fa nilpotent, e.g., $(Fa)^{N+1} = 0$:
 $(1 - a^2)(F + FaF + FaFaF + \dots + F(aF)^N)(1 - a^2)$
 Algebra : finitistic (prestellar), to be defined.

7 – THE ALGEBRA \mathbb{U}

- Fixed functional language ; irrelevant provided contains $g(\cdot, \cdot)$ and a .

Monomials : tu^* with *same* bound variables.

Examples : $f(x)g(x, x)^*$ but $f(x)g(x, y)^*$ excluded.

Composition : $tu^* \cdot vw^* := t\theta w\theta^*$ if u, v *unifiable* with *mg*u θ .

Otherwise $tu^* \cdot vw^* := 0$.

Example : $xx^* \cdot tu^* = tu^*$.

Involution : $(tu^*)^* := ut^*$; *contravariant* compositionwise.

- Implemented on *Hilbert space* ℓ^2 spanned by closed terms.

$tu^*[a] := b$ when $tu^* \cdot aa_0^* := ba_0^*$; works since *same* variables.

$tu^*[a] := 0$ when $tu^* \cdot aa_0^* := 0$.

Involution implemented as *adjoint*. Indeed $t, u : \ell^2 \otimes \dots \otimes \ell^2 \mapsto \ell^2$.

- Algebra \mathbb{U} spanned by the tu^* admits a stellar norm $\|uu^*\| = \|u\|^2$.

Hence a *greatest* one ; the one to select in case.

- Keep \mathbb{U} *algebraic* including complex coefficients.

8 – NORMALISATION, REVISITED

- Suggested by quantum physics (Schrödinger's cut).

Projections : $u = u^* = u^2$ (correspond to Danos-Regnier *partitions*).

Implemented as closed *subspace* $|u|$ of Hilbert space ℓ^2 .

- **Replace** partial symmetries, typically :

$$\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} = f(x)g(x)^* + g(x)f(x)^* \text{ with :}$$

$$\begin{pmatrix} 1/2 & 1/2 \\ 1/2 & 1/2 \end{pmatrix} = \sqrt{2}/2(f(x) + g(x))\sqrt{2}/2(f(x)^* + g(x)^*).$$

- **Intersection** $u \cap v$ as « limit » of decreasing sequence

$$u \geq uvu \geq uvuvu \geq \dots \geq (uv)^n u \geq \dots$$

Provided $\text{Sp}(uv) \subset \{\lambda_1, \dots, \lambda_n, 1\}$ ($0 \leq \lambda_i < 1$)

$$uv \cdot (1 - \lambda_1)^{-1}(uv - \lambda_1) \cdot \dots \cdot (1 - \lambda_n)^{-1}(uv - \lambda_n) \cdot u$$

Dually : $u \uplus v := 1 - ((1 - u) \cap (1 - v))$ closed linear span of *union*.

- **Application of « function » F to « argument » a .**

$$[F]a := (F \uplus a) \cap (1 - a), \text{ provided } F \cap a = 0.$$

Graph-like : $x \in |a|, (-x) \oplus y \in |F| \mapsto y \in [F]a$.

III — ENVOI

What logic is about.

9 – A PROMETHEAN PROJECT

- **Legitimate** doubts about :
 - Principles** : church, politics, axiomatics (= army).
 - Realism** : semantics is *reasoning* about a « **reality** » we are part of.
- Pure reason : what can we *sure* of ?
 - Aristotle** : syllogistics refers to itself.
- Mistreated by scientism and obsolete *foundational* projects.
 - Axiomatic realism** : language refers to reality ($\forall \Rightarrow \exists$).
 - Systems** : contingent truths ; same reality, different spectacles.
- Linearity.
 - Not yet another system** : linear implies system-free.
- *Open* architecture ; independent parts compatible.
 - Usage** : Curry-like (existentialist, with cut).
 - Usine** : Church-like (not quite essentialist, cut-free).
 - Adequation** : external, proof of normalisation. Forget foundations !