

Convergence Bounds: Differences and Similarities

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PinT 2022, 11–15 July 2022

Setting and the Parareal algorithm

Consider the ODE (coming from the method-of-lines),

$$\frac{\partial u(t)}{\partial t} = L u(t) + f(t), \quad t \in (0, T], \quad u(0) = u_0 .$$

After time-discretization, we obtain the time-stepping procedure

$$u_{n+1} = \Phi u_n + f_n, \quad n = 0, \dots, n_t .$$

For example, for Forward Euler,

$$\Phi = 1 + \Delta t L .$$

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The [Parareal algorithm](#) is given by the iterations

$$\begin{cases} U_0^k = u_0 & k = 0, \dots, K \\ U_{n+1}^0 = G U_n^0 + f_n & n = 0, \dots, n_t \\ U_{n+1}^{k+1} = F U_n^k + G U_n^{k+1} - G U_n^k + f_n & k = 0, \dots, K, n = 0, \dots, n_t \end{cases}$$

where the coarse operator G is a [cheap approximation](#) of the fine operator $F = \Phi^m$.

From the iteration

$$U_{n+1}^{k+1} = F U_n^k + G U_n^{k+1} - G U_n^k + f_n, \quad n = 1, \dots, n_t,$$

the error, $\epsilon_n^k := u_n - U_n^k$, can be computed as

$$\epsilon_{n+1}^{k+1} = F \epsilon_n^k + G \epsilon_n^{k+1} - G \epsilon_n^k.$$

From the iteration

$$U_{n+1}^{k+1} = F U_n^k + G U_n^{k+1} - G U_n^k + f_n, \quad n = 1, \dots, n_t,$$

the **error**, $\epsilon_n^k := u_n - U_n^k$, can be computed as

$$\epsilon_{n+1}^{k+1} = F \epsilon_n^k + G \epsilon_n^{k+1} - G \epsilon_n^k.$$

In turn, it can be bounded as

$$\|\epsilon_{n+1}^{k+1}\| =: e_{n+1}^{k+1} \leq \underbrace{\|F - G\|}_{=\alpha} e_n^k + \underbrace{\|G\|}_{=\beta} e_n^{k+1}.$$

We thus only need to solve the iteration

$$e_{n+1}^{k+1} = \alpha e_n^k + \beta e_n^{k+1}.$$

F, G



Reformulation



α, β



Recurrence solving



Bounding

$$\epsilon_{n+1}^{k+1} = (F - G) \epsilon_n^k + G \epsilon_n^{k+1}$$

Bounding

$$e_{n+1}^{k+1} = \alpha e_n^k + \beta e_n^{k+1}$$

The iteration $\mathbf{e}_{n+1}^k = \alpha \mathbf{e}_n^{k-1} + \beta \mathbf{e}_n^k$ can be written in matrix form as

$$\begin{pmatrix} I & & & & \\ -\beta & I & & & \\ & & \ddots & & \\ & & & \ddots & \\ & & & & -\beta & I \end{pmatrix} \begin{pmatrix} \mathbf{e}_0^k \\ \mathbf{e}_1^k \\ \vdots \\ \mathbf{e}_{n_t}^k \end{pmatrix} = \begin{pmatrix} 0 & & & & \\ \alpha & 0 & & & \\ & \ddots & & & \\ & & \ddots & & \\ & & & \alpha & 0 \end{pmatrix} \begin{pmatrix} \mathbf{e}_0^{k-1} \\ \mathbf{e}_1^{k-1} \\ \vdots \\ \mathbf{e}_{n_t}^{k-1} \end{pmatrix}.$$

Lemma (Recurrence solving)

Assuming that α and β are scalars, the error at step k is given by

$$\mathbf{e}^k = M(\beta) (I_{n_t} \otimes \alpha) \mathbf{e}^{k-1} = \dots = M(\beta)^k (I_{n_t} \otimes \alpha^k) \mathbf{e}^0,$$

where the Toeplitz matrix M is defined as

$$M(\beta) = \begin{pmatrix} 0 & & & & \\ I & 0 & & & \\ \beta & I & 0 & & \\ \vdots & \ddots & \ddots & \ddots & \\ \beta^{n_t-1} & \dots & \beta & I & 0 \end{pmatrix}.$$

F, G



Reformulation



α, β



Recurrence solving



Bounding

$$\epsilon_{n+1}^{k+1} = (F - G) \epsilon_n^k + G \epsilon_n^{k+1}$$

Bounding

$$e_{n+1}^{k+1} = \alpha e_n^k + \beta e_n^{k+1}$$

$$e^k = M(\beta)^k (I_{n_t} \otimes \alpha^k) e^0$$

Bounding

We then want to bound

$$\mathbf{e}^k = M(\beta)^k (I_{n_t} \otimes \alpha^k) \mathbf{e}^0 .$$

Linear Bound [Lemma 4.4, Gander, Vandewalle, 2007]

$$\|M(\beta)^k\|_\infty \leq \|M(\beta)\|_\infty^k = \left(\frac{1 - |\beta|^{n_t}}{1 - |\beta|} \right)^k .$$

Superlinear Bound [Lemma 4.3, Gander, Vandewalle, 2007]

$$\begin{aligned} \|M(\beta)^k\|_\infty &= \sum_{i=0}^{n_t-k} \binom{i+k-1}{k-1} |\beta|^i \\ &= \frac{1}{(k-1)!} \sum_{i=0}^{n_t-k} \left[\prod_{l=1}^{k-1} (i+l) \right] |\beta|^i \end{aligned}$$

Bounding

We then want to bound

$$\mathbf{e}^k = M(\beta)^k (I_{n_t} \otimes \alpha^k) \mathbf{e}^0 .$$

Linear Bound [Lemma 4.4, Gander, Vandewalle, 2007]

$$\|M(\beta)^k\|_\infty \leq \|M(\beta)\|_\infty^k = \left(\frac{1 - |\beta|^{n_t}}{1 - |\beta|} \right)^k .$$

Explicit Superlinear Bound [Lemma 4.4, Gander, Vandewalle, 2007]

If $|\beta| < 1$, then

$$\|M(\beta)^k\|_\infty \leq \binom{n_t}{k} = \frac{1}{k!} \prod_{l=0}^{k-1} (n_t - l) .$$

General structure

F, G



Reformulation



α, β



Recurrence solving



Bounding

$$\epsilon_{n+1}^{k+1} = (F - G) \epsilon_n^k + G \epsilon_n^{k+1}$$

Bounding

$$e_{n+1}^{k+1} = \alpha e_n^k + \beta e_n^{k+1}$$

$$e^k = M(\beta)^k (I_{n_t} \otimes \alpha^k) e^0$$

Linear, (explicit) superlinear bounds for $\|M(\beta)^k\|$

Formulation for eigenvalues

We said that $\alpha = \|F - G\|$ and $\beta = \|G\|$, but they do not have to be.

Assume that F and G are **simultaneously diagonalizable**, then there exists a matrix U such that $F = U\Lambda U^{-1}$ and $G = UH U^{-1}$, where $\Lambda = \text{diag}(\lambda_1^m, \dots, \lambda_{n_x}^m)$ and $H = \text{diag}(\mu_1, \dots, \mu_{n_x})$.

$$\epsilon_{n+1}^{k+1} = (F - G)\epsilon_n^k + G\epsilon_n^{k+1}$$

becomes the **decoupled system**,

$$\tilde{\epsilon}_{n+1}^{k+1} = (\Lambda - H)\tilde{\epsilon}_n^k + H\tilde{\epsilon}_n^{k+1}.$$

We are left with bounding equations

$$\tilde{\epsilon}_{n+1}^{k+1}(\omega) = \underbrace{(\lambda_\omega^m - \mu_\omega)}_{=\alpha} \tilde{\epsilon}_n^k(\omega) + \underbrace{\mu_\omega}_{=\beta} \tilde{\epsilon}_n^{k+1}(\omega).$$

F, G



Reformulation



α, β



Recurrence solving



Bounding

$$\epsilon_{n+1}^{k+1} = (F - G) \epsilon_n^k + G \epsilon_n^{k+1}$$

Bounding, Simultaneous diagonalization

$$e_{n+1}^{k+1} = \alpha e_n^k + \beta e_n^{k+1}$$

$$e^k = M(\beta)^k (I_{n_t} \otimes \alpha^k) e^0$$

Linear, (explicit) superlinear bounds for $\|M(\beta)^k\|$

Parareal and MGRIT bounds

Recall: $\mathbf{e}^k = M(\beta)^k (I_{n_t} \otimes \alpha^k) \mathbf{e}^0$.

→ Do these bounds correspond to some we already know?

- If $\alpha = \|F - G\|$ and $\beta = \|G\|$, (“norm bounds”)
 - LINEAR: Gander, Vandewalle (2007).
 - EXPLICIT SUPERLINEAR: Gander, Vandewalle (2007), Gander, Hairer (2008).
 - SUPERLINEAR: Gander, Lunet, Ruprecht, de Steck (2022).
- If $\alpha = |\lambda_\omega^m - \mu_\omega|$ and $\beta = |\mu_\omega|$, (“eigenvalue bounds”)
 - LINEAR: Dobrev, Kolev, Petersson, Schroder (2017).
 - (EXPLICIT) SUPERLINEAR: do not exist yet.

Questions:

→ How do those bounds compare to each other?

→ What are the advantages of each of these bounds?

Example: Dahlquist equation $\partial_t u = (-1 + i) u$

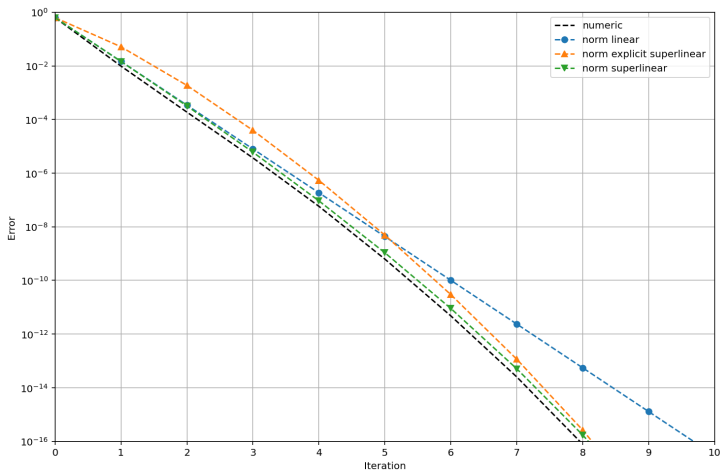


Figure: Norm bounds for the Dahlquist equation for $T = 5$.

Example: Dahlquist equation $\partial_t u = (-1 + i) u$

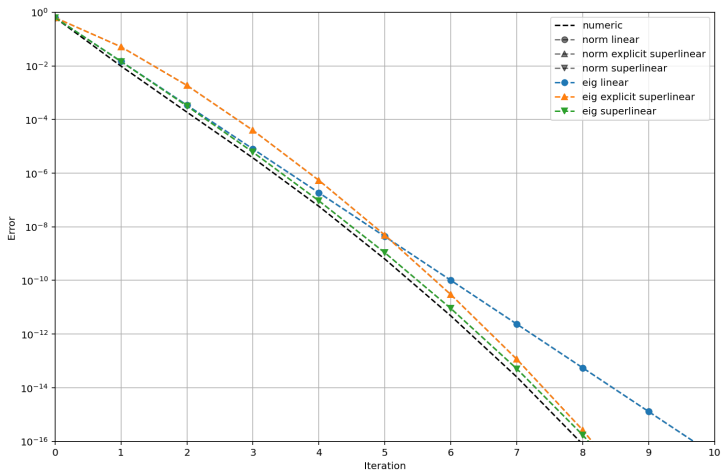
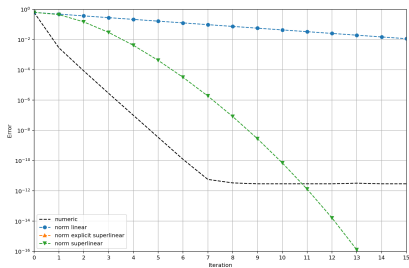
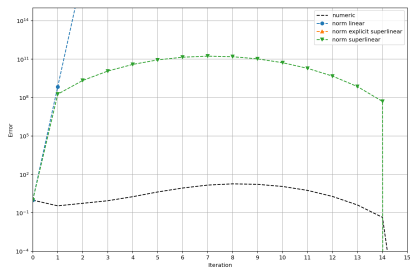


Figure: Eigenvalue bounds for the Dahlquist equation for $T = 5$.

Example: heat and advection equations

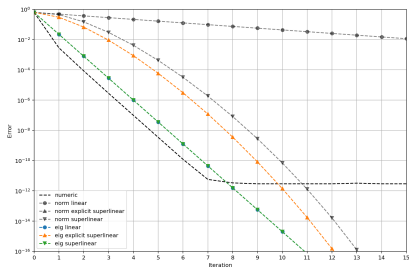


(a) Heat

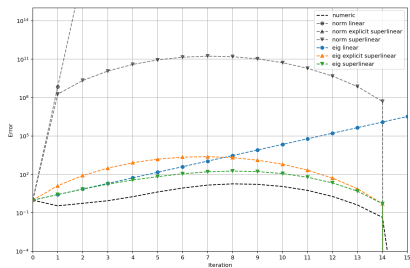


(b) Advection

Example: heat and advection equations



(a) Heat



(b) Advection

Generalization to Multigrid Reduction in Time

For Multigrid Reduction in Time (MGRIT) with *FCF*-relaxation, the error propagator is given by

$$E = B_{\Delta}^{-1}(B_{\Delta} - A_{\Delta})(I - A_{\Delta}),$$

that is,

$$E = \left(\sum_{j=0}^{n_t-1} U_{-j-1} \otimes \Psi^j \right) (I \otimes (\Phi^m - \Psi)) (U_{-1} \otimes \Phi^m).$$

In our notation,

$$E = M(G)(I \otimes (F - G))(U_{-1} \otimes F).$$

Generalization to Multigrid Reduction in Time

For Multigrid Reduction in Time (MGRIT) with $F(CF)^\nu$ -relaxation, the error propagator is given by

$$E = B_\Delta^{-1}(B_\Delta - A_\Delta)(I - A_\Delta)^\nu,$$

that is,

$$E = \left(\sum_{j=0}^{n_t-1} U_{-j-1} \otimes \Psi^j \right) (I \otimes (\Phi^m - \Psi)) (U_{-\nu} \otimes \Phi^{m\nu}).$$

In our notation,

$$E = M(G)(I \otimes (F - G))(U_{-\nu} \otimes F^\nu).$$





Conclusion

Summary:

- General structure for Parareal and MGRIT bounds.
- The new superlinear bound captures convergence the best.
- Noticed that:
 - Eigenvalue bounds are more **accurate**.
 - Norm bounds are **general**.

Future Work:

- Accurate convergence bounds for spatial coarsening.
- Generalization of this bounding strategy to other smoothers.
- Convergence bounds for variable time-step sizes.

-  V. A. Dobrev et al. “Two-Level Convergence Theory for Multigrid Reduction in Time (MGRIT)”. In: *SIAM J. Sci. Comput.* 39.5 (Jan. 2017). Publisher: Society for Industrial and Applied Mathematics, S501–S527.
-  M. J. Gander and E. Hairer. “Nonlinear convergence analysis for the parareal algorithm”. In: *Domain decomposition methods in science and engineering XVII*. Springer, 2008, pp. 45–56.
-  M. J. Gander and S. Vandewalle. “Analysis of the Parareal Time-Parallel Time-Integration Method”. In: *SIAM J. Sci. Comput.* 29.2 (Jan. 2007), pp. 556–578.
-  M. J. Gander et al. *A unified analysis framework for iterative parallel-in-time algorithms*. en. Number: arXiv:2203.16069 arXiv:2203.16069 [cs, math]. Mar. 2022.

Why eigenvalue bounds are better than norm bounds?

If we express **norm bound** in terms of eigenvalues, the linear bound is given by

$$\|E_\omega\|_\infty \leq \sup_\omega \frac{1 - |\mu_\omega|^{n_t}}{1 - |\mu_\omega|} \sup_\omega |\lambda_\omega^m - \mu_\omega| ,$$

whereas the linear **eigenvalue bound** is given by

$$\|E_\omega\|_\infty \leq \sup_\omega \left\{ \frac{1 - |\mu_\omega|^{n_t}}{1 - |\mu_\omega|} |\lambda_\omega^m - \mu_\omega| \right\} .$$

Bounds with coarse initialization

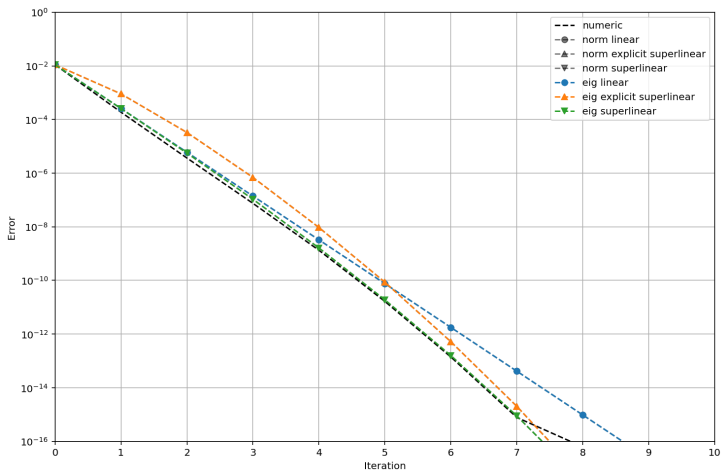


Figure: Bounds with coarse initialization for the Dahlquist equation $\partial_t u = (-1 + i)u$ with the same parameters as on slide 12.