

An Algorithm for Testing the Half-plane Property of Matroids

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Joint work with Mario Kummer

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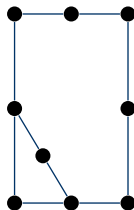
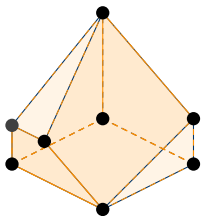
Journées Nationales de Calcul Formel

03.03.2022

1 Hyperbolic Polynomials and Spectrahedral Cones

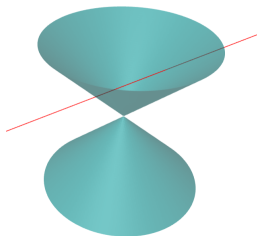
2 Connection to Matroids

3 An Algorithm for the Half-Plane Property of Matroids



Hyperbolic Polynomials

Definition: A homogeneous polynomial $h \in \mathbb{R}[x_1, \dots, x_n]$ is called hyperbolic with respect to $e \in \mathbb{R}^n$ if $h(e) \neq 0$ and for all $v \in \mathbb{R}^n$, $h(et - v)$ in $\mathbb{R}[t]$ has only real roots.

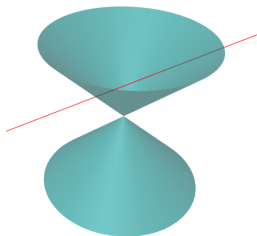


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The hyperbolicity cone of h at e is

$$C_h(e) = \{v \in \mathbb{R}^n : h(et - v) = 0 \implies t \in \mathbb{R}_{\geq 0}\}.$$



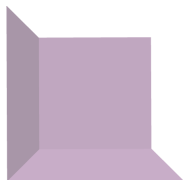
Hyperbolic Polynomials

$$h(et - v) = h(e_1 t - v_1, \dots, e_n t - v_n)$$

$$C_h(e) = \left\{ v \in \mathbb{R}^n : h(et - v) = 0 \implies t \in \mathbb{R}_{\geq 0} \right\}$$

$$h = x_1 x_2 x_3, \quad e = (1, 1, 1)$$

$$h = (t - v_1)(t - v_2)(t - v_3)$$



$$C_h(e) = \mathbb{R}_{\geq 0}^3$$

$$h = \det(X), \quad X = \begin{pmatrix} x_1 & x_2 \\ x_2 & x_3 \end{pmatrix}, \quad e = I$$

$$\det \begin{pmatrix} v_1 - t & v_2 \\ v_2 & v_3 - t \end{pmatrix}$$



$$C_h(e) = \left\{ v \in \mathbb{R}^3 : \begin{pmatrix} v_1 & v_2 \\ v_2 & v_3 \end{pmatrix} \succeq 0 \right\}$$

Cone of PSD 2×2 matrices

Determinantal Representability

Definition: A homogeneous polynomial $f \in \mathbb{R}[x_1, \dots, x_n]$ is said to have a determinantal representation if there are PSD matrices A_1, \dots, A_n such that

$$f = \lambda \det(x_1 A_1 + \dots + x_n A_n)$$

for some $\lambda \in \mathbb{R}$.

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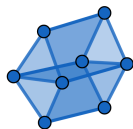
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f is weakly determinantal \iff f is hyperbolic with respect to all $e \in \mathbb{R}_{>0}^n$



Counter example: The basis generating polynomial of the Vamós matrix

Spectrahedral Cones

Definition: A convex cone C is called spectrahedral if

$$C = \{v \in \mathbb{R}^n : A(v) = v_1 A_1 + \dots + v_n A_n \succeq 0\}$$

where A_1, \dots, A_n are real symmetric $d \times d$ matrices.

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- Spectrahedral cones are hyperbolicity cones. (consider $h = \det(A_1 x_1 + \dots + A_n x_n)$).

Note: For the rest of the talk “hyperbolic” refers to hyperbolic with respect to every point in the positive orthant.

Spectrahedral Representability

Question: Given a homogeneous polynomial $h \in \mathbb{R}[x_1, \dots, x_n]$ that is hyperbolic. When is $C_h(e)$ spectrahedral?

h has a determinantal representation $\implies C_h$ is spectrahedral.

Theorem (Helton-Vinnikov, 2007) Let $h \in \mathbb{R}[x_1, \dots, x_n]$ be hyperbolic. The hyperbolicity cone C_h is spectrahedral if and only if there exists a hyperbolic polynomial g with $C_h \subset C_g$ such that $h \cdot g$ has a determinantal representation.

$$C_{h \cdot g} = C_h \cap C_g$$

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Question: Are all hyperbolicity cones spectrahedral?

Generalized Lax Conjecture

Conjecture: Every hyperbolicity cone is spectrahedral.

Every hyperbolic program can be written as a semi-definite program.

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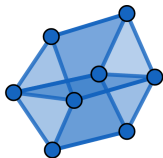
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- The conjecture is true for matching polynomials of simple graphs. (Amini, 2019).

C spectrahedral $\implies C$ is a hyperbolicity cone for some h .

h has a determinantal representation $\implies C_h$ is spectrahedral.

C_h is spectrahedral $\iff \exists g$ hyperbolic with $C_h \subset C_g$ such that $h \cdot g$ has a determinantal representation.

The basis generating polynomial of the Vamos matroid is hyperbolic, but not weakly determinantal!



Connection to Matroids

Definition: A matroid M is $E = [n]$ with a collection \mathcal{B} of its subsets (bases) satisfying

If $B_1, B_2 \in \mathcal{B}$ and $x \in B_1 \setminus B_2$,

then $\exists y \in B_2 \setminus B_1$ such that $(B_1 \setminus \{x\}) \cup \{y\} \in \mathcal{B}$.

The basis generating polynomial of M is $h_M := \sum_{B \in \mathcal{B}} \prod_{i \in B} x_i$.

$$\begin{array}{cccc} 1 & 2 & 3 & 4 \\ \left(\begin{array}{cccc} 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{array} \right) \end{array}$$

$$E = \{1, 2, 3, 4\}$$

$$\mathcal{B} = \{\{2, 3, 4\}, \{2, 1, 4\}, \{1, 3, 4\}\}$$

$$h_M = x_2 x_3 x_4 + x_2 x_1 x_4 + x_1 x_3 x_4$$

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h_M is homogeneous and multiaffine

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Definition: A homogeneous polynomial $f \in \mathbb{R}[x_1, \dots, x_n]$ is said to have the half-plane property if there exists an open half-plane $\mathcal{H} \subset \mathbb{C}$ with $0 \in \partial\mathcal{H}$ such that $f(x_1, \dots, x_n) \neq 0$ for $x_1, \dots, x_n \in \mathcal{H}$.

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Theorem(Choe et. al., 2004): Support of a homogeneous multiaffine polynomial $f \in \mathbb{R}[x_1, \dots, x_n]$ with the half-plane property is the collection of bases of some matroid M .

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Theorem(Choe et. al., 2004): Support of a homogeneous multiaffine polynomial $f \in \mathbb{R}[x_1, \dots, x_n]$ with the half-plane property is the collection of bases of some matroid M .

Let M be a matroid with the basis generating polynomial $h_M \in \mathbb{R}[x_1, \dots, x_n]$.

h_M is weakly determinantal $\implies h_M$ has the half-plane property $\iff h_M$ is hyperbolic

Connection to Matroids

Questions:

- Do all matroids M have the half-plane property (i.e., h_M is hyperbolic)?

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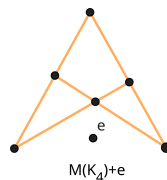
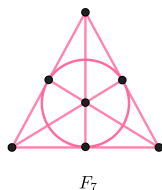
Theorem(Choe et. al., 2004): The half-plane property is closed under taking minors and direct sums of matroids.

Theorem(Kummer and S., 2021): Being weakly determinantal and having a spectrahedral hyperbolicity cone are closed under taking minors.

Classification of Matroids

Theorem(Choe et. al., 2004):

- All matroids on at most 6 elements have the half-plane property.
- Matroids that have rank or corank 2 have the half-plane property.
- Fano matroid F_7 , F_7^- , F_7^{--} , F_7^{-3} , $M(K_4) + e$, P_8 , P_8^- , P_8^{--} don't have the half-plane property.
- 6th root of unity matroids are weakly determinantal.



A Criteria for the Half-plane Property

Theorem (Brändén, 2007 - Wagner and Wei, 2009):

Let h_M be the basis generating polynomial of a matroid M . The following are equivalent:

- h_M has the half-plane property.
- For all $1 \leq i, j \leq n$, the Rayleigh difference

$$\Delta_{ij}(h_M) := \frac{\partial h_M}{\partial x_i}(x) \frac{\partial h_M}{\partial x_j}(x) - \frac{\partial^2 h_M}{\partial x_i \partial x_j}(x) h_M(x) \geq 0$$

for all $x \in \mathbb{R}^n$

(We call them SOS-Rayleigh if $\Delta_{ij}(h_M)$ is SOS for all i, j).

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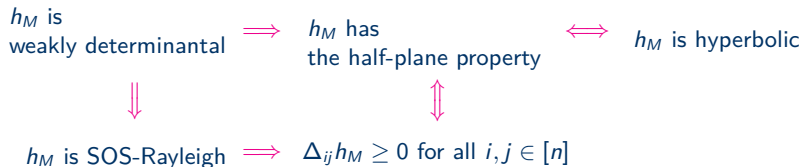
(We call them SOS-Rayleigh if $\Delta_{ij}(h_M)$ is SOS for all i, j).

- All of its proper minors have the half-plane property and for some $1 \leq i, j \leq n$, $\Delta_{ij}(h_M) \geq 0$ for all $x \in \mathbb{R}^n$.

A Criteria for Being Weakly Determinantal

Theorem(Kummer-Plaumann-Vinzant, 2015):

Let h_M be a basis generating polynomial of some matroid M . If h_M is weakly determinantal, then it is SOS-Rayleigh.



An Algorithm for the Half-plane Property

Input: A matroid M on ground set $E = [n]$ with the collection of bases \mathcal{B}
all of whose proper minors have the HPP

$$h_M := \sum_{B \in \mathcal{B}} \prod_{i \in B} x_i$$

$$J := \{(i, j) : 0 \leq i, j \leq n, i \neq j\}$$

For (i, j) in J Do

$$\Delta_{ij} := \frac{\partial h_M}{\partial x_i}(x) \frac{\partial h_M}{\partial x_j}(x) - \frac{\partial^2 h_M}{\partial x_i \partial x_j}(x) h_M(x)$$

solveSOS Δ_{ij}

SDP that attempts to find
a PSD Gram matrix G
with rational entries s.t.
 $m^T G m = \Delta_{ij}$

Use M2 package
"SumsOfSquares"

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 solveSOS Δ_{ij}

 If a PSD Gram matrix G
 with entries in \mathbb{Q} is found

Δ_{ij} is SOS
 M has the HPP – Stop

 else
 Continue

End

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End

Go to Julia for
finding negative points

For each (i, j) finds critical points of Δ_{ij}
using "HomotopyContinuation.jl"
and evaluates Δ_{ij} at the critical points

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solveSOS Δ_{ij}

If a PSD Gram matrix G
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else

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End

Go to Julia for
finding negative points

If a negative value is found

M doesn't have HPP

Else

Undetected

An Algorithm for SOS-Rayleigh

Input: A matroid M on ground set $E = [n]$ with the collection of bases \mathcal{B}
with the HPP

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`solveSOS` Δ_{ij}

If a PSD Gram matrix G
with entries in \mathbb{Q} is found

Δ_{ij} is SOS
Continue

else

Stop

Try to produce a
non-SOS certificate

Use M2 package
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End

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SDP that attempts to find
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for all matrices G_i that define
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solveSOS Δ_{ij}

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Try to produce a
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If certified

M isn't SOS-Rayleigh

Else

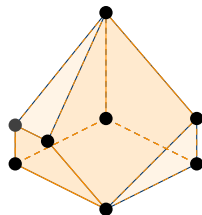
Undetected

End

→ **M is SOS-Rayleigh**

Matroids on 8 Elements of Rank 3 or 4

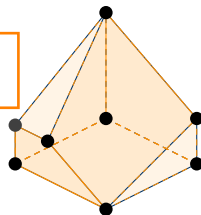
Properties	Count
Simple	685
Simple, connected and without the 10 forbidden minors	309
Having the HPP	287
SOS-Rayleigh	256
With the HPP and not SOS-Rayleigh	14
With the HPP and SOS-Rayleigh undetected	17
Without the HPP	22



Matroids on 8 Elements of Rank 3 or 4

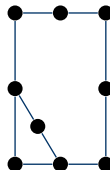
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In particular, they are not weakly determinantal.
Good candidates for searching for a counter example.



Matroids on 9 Elements of rank 3

Properties	Count
Simple	383
Simple, connected and without the 10 forbidden minors	119
Having the HPP	116
With the HPP and SOS-Rayleigh	106
With the HPP and SOS-Rayleigh undetected	10
Without the HPP	3



Matroids on 9 Elements of rank 4

Properties	Count
Simple	185982
Simple, connected and without the 35 excluded minors	6718
Having the HPP	4125
Candidates for having the HPP	819
Without the HPP	1218
HPP undetected	556

Merci pour votre attention!