

# Generic aspects of entanglement in high-dimensional quantum systems

Some applications of asymptotic geometric analysis in quantum information theory

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- 1 Introduction
- 2 Quantifying the typical amount of entanglement/correlations in multipartite pure states
- 3 Typical strength of entanglement criteria for bipartite mixed states

## Overview of the lecture's topic

**Goal of this lecture:** Understand how *asymptotic geometric analysis* can be useful to tackle problems arising in *quantum information theory*.

[ The reference: *Alice and Bob meet Banach*, by G. Aubrun and S. Szarek. ]

**Question:** What is asymptotic geometric analysis?

Use of probabilistic techniques in the study of Banach spaces of high but finite dimension.

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Why “high dimension”?

Quantum system composed of 1 particle: described by a complex Hilbert space  $H$ .

Quantum system composed of  $M$  such particles: the associated space is  $H^{\otimes M}$ .

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Why “probabilistic techniques”?

- 1 Identify typical properties of high dimensional quantum systems.
- 2 Prove the existence of quantum systems having certain properties, using random constructions.

**In this talk:** We will focus on point (1), with a particular interest for multipartite quantum systems and properties related to entanglement.

## Two seminal works from 2006:

- 1 Aspects of generic entanglement, by P. Hayden, D. Leung and A. Winter.  
Goal: Quantify the typical amount of entanglement in bipartite quantum states.
- 2 Tensor products of convex sets and the volume of separable states on  $N$  qudits, by G. Aubrun and S. Szarek.  
Goal: Study the typical performance of conditions which are checked in practice to guarantee the entanglement of bipartite quantum states.

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**Note:** “typical” = “with probability going to 1 (exponentially) as the underlying dimension grows”. So there are usually two steps in the argument:

- 1 Identify the average behavior of the property under consideration.
- 2 Show that this average behavior is generic in high dimension.  
*Concentration of measure phenomenon:* a sufficiently ‘well-behaved’ function has an exponentially small probability of deviating from its average as the dimension grows.



## Reminder: separability vs entanglement in multipartite quantum systems

$H_1, \dots, H_M$  complex Hilbert spaces. In this talk: of finite, but usually large, dimension.  
→  $H_i \equiv \mathbf{C}^{d_i}$  with  $d_i \gg 1$ , for  $1 \leq i \leq M$ .

### Definition [Separability and entanglement]

A state  $\rho$  on  $H_1 \otimes \dots \otimes H_M$  is called *separable* if it is a convex combination of product states, i.e.

↳ positive semidefinite operator with trace 1 on  $H_1 \otimes \dots \otimes H_M$

$$\rho = \sum_{k=1}^r \lambda_k \rho_1^k \otimes \dots \otimes \rho_M^k, \text{ with } \begin{cases} \lambda_k \geq 0, 1 \leq k \leq r, \sum_{k=1}^r \lambda_k = 1 \\ \rho_i^k \text{ state on } H_i, 1 \leq k \leq r, 1 \leq i \leq M \end{cases} .$$

Otherwise it is called *entangled*.

[ Note: If  $\rho$  is a pure state, i.e.  $\rho = |\psi\rangle\langle\psi|$  for some unit vector  $\psi \in H_1 \otimes \dots \otimes H_M$ , then  $\rho$  is separable iff  $\psi = \psi_1 \otimes \dots \otimes \psi_M$  for some unit vectors  $\psi_i \in H_i, 1 \leq i \leq M$ . ]

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**Fact:** If a multipartite quantum system is in a separable state, there is no intrinsically quantum correlation between its subsystems. It thus does not provide any advantage over a classical system in information processing tasks.

→ Characterizing and quantifying the entanglement of multipartite quantum states is an important issue in practice.

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## 2 Quantifying the typical amount of entanglement/correlations in multipartite pure states

- Typical amount of entanglement in uniformly distributed multipartite pure states
- What about more 'physically relevant' multipartite pure states?

## 3 Typical strength of entanglement criteria for bipartite mixed states

- Entanglement detection of bipartite states
- Estimating the 'size' of sets of states in high-dimensional bipartite systems
- Typical properties of random high-dimensional bipartite states

# Characterizing and quantifying pure state entanglement through tensor norms

Let  $\psi \in \mathbf{H}_1 \otimes \cdots \otimes \mathbf{H}_M$  be a pure state, i.e.  $\|\psi\|_2 = 1$ .

↳ Euclidean norm

Its *injective norm* is:  $\|\psi\|_\varepsilon := \sup \{ |\langle \varphi_1 \otimes \cdots \otimes \varphi_M | \psi \rangle| : \varphi_i \in \mathbf{H}_i, \|\varphi_i\|_2 = 1 \}$ .

Its *projective norm* is:  $\|\psi\|_\pi := \inf \left\{ \sum_{k=1}^r |\alpha_k| : \chi_i^k \in \mathbf{H}_i, \|\chi_i^k\|_2 = 1, \psi = \sum_{k=1}^r \alpha_k \chi_1^k \otimes \cdots \otimes \chi_M^k \right\}$ .

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## Facts:

- The  $\varepsilon$  and  $\pi$  norms are *dual norms*:  $\|\psi\|_\varepsilon = \sup_{\|\varphi\|_\pi \leq 1} |\langle \varphi | \psi \rangle|$  and  $\|\psi\|_\pi = \sup_{\|\varphi\|_\varepsilon \leq 1} |\langle \varphi | \psi \rangle|$ .
- The  $\varepsilon$  and  $\pi$  norms are, respectively, minimal and maximal norms among *tensor norms* (i.e. norms which factorize on product vectors:  $\|\psi_1 \otimes \cdots \otimes \psi_M\| = \|\psi_1\| \cdots \|\psi_M\|$ ).  
In particular:  $\|\psi\|_\varepsilon \leq \|\psi\|_2 \leq \|\psi\|_\pi$ .
- A pure state  $\psi$  is separable iff  $\|\psi\|_\varepsilon = \|\psi\|_\pi = 1$ .  
→ If  $\|\psi\|_\varepsilon \ll 1$  or  $\|\psi\|_\pi \gg 1$ , then  $\psi$  is 'very' entangled.

## Particular case of bipartite pure states

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We can identify  $|\psi\rangle = \sum_{k,l=1}^d \psi_{kl} |kl\rangle \in \mathbf{C}^d \otimes \mathbf{C}^d$  with  $M_\psi = \sum_{k,l=1}^d \psi_{kl} |k\rangle\langle l| \in \mathcal{M}_d(\mathbf{C})$ .

Then clearly,  $\|\psi\|_2 = \|M_\psi\|_2$ .

And the *Schmidt decomposition* of  $\psi$  corresponds to the *singular value decomposition* of  $M_\psi$ :

$$|\psi\rangle = \sum_{k=1}^r \sqrt{\lambda_k} |e_k f_k\rangle \longleftrightarrow M_\psi = \sum_{k=1}^r \sqrt{\lambda_k} |e_k\rangle\langle f_k|,$$

with  $r \leq d$  the *Schmidt rank* of  $\psi$ ,  $\sum_{k=1}^r \lambda_k = 1$ ,  $\{e_k\}_{k=1}^r, \{f_k\}_{k=1}^r$  orthonormal sets in  $\mathbf{C}^d$ .

So  $\|\psi\|_\varepsilon = \max_{1 \leq k \leq r} \sqrt{\lambda_k} = \|M_\psi\|_\infty$  and  $\|\psi\|_\pi = \sum_{k=1}^r \sqrt{\lambda_k} = \|M_\psi\|_1$ .



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—> Checking bipartite pure state separability is easy.

Quantitatively, for all  $\psi \in \mathbf{C}^d \otimes \mathbf{C}^d$  s.t.  $\|\psi\|_2 = 1$ ,  $\frac{1}{\sqrt{d}} \leq \|\psi\|_\varepsilon \leq 1$  and  $1 \leq \|\psi\|_\pi \leq \sqrt{d}$ .

But no such simple characterization for  $M > 2$  (no equivalent of the Schmidt decomposition).

## Geometric measure of entanglement

### Definition [Geometric measure of entanglement (Wei/Goldbart)]

Let  $\psi \in \mathbb{H}_1 \otimes \cdots \otimes \mathbb{H}_M$  be a pure state. Its *geometric measure of entanglement (GME)* is

$$E(\psi) := -\log \sup \{ |\langle \varphi_1 \otimes \cdots \otimes \varphi_M | \psi \rangle|^2 : \varphi_i \in \mathbb{H}_i, \|\varphi_i\|_2 = 1 \}.$$

By definition,  $E(\psi) = -2 \log \|\psi\|_\varepsilon$ . So  $E(\psi) = 0$  iff  $\psi$  is separable.

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**Remark:** The definition of the GME can be extended to mixed states on  $H_1 \otimes \cdots \otimes H_M$  (but it is not an entanglement measure anymore), as

$$\begin{aligned} E(\rho) &:= -\log \sup \{ \langle \varphi_1 \otimes \cdots \otimes \varphi_M | \rho | \varphi_1 \otimes \cdots \otimes \varphi_M \rangle : \varphi_i \in H_i, \|\varphi_i\|_2 = 1 \} \\ &= -\log \sup \{ \text{Tr}(\rho \sigma) : \sigma \text{ separable state on } H_1 \otimes \cdots \otimes H_M \}. \end{aligned}$$

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**Fact:** For any unit vector  $\psi \in (\mathbf{C}^d)^{\otimes M}$ ,  $\|\psi\|_\varepsilon \geq \frac{1}{\sqrt{d^{M-1}}}$ , i.e.  $E(\psi) \leq (M-1) \log d$ .

This can be checked recursively, starting from the bipartite case:

For any unit vector  $\psi \in \mathbf{C}^{d_1} \otimes \mathbf{C}^{d_2}$ ,  $\|\psi\|_\varepsilon \geq \frac{1}{\sqrt{d}}$ , where  $d := \min(d_1, d_2)$ .

↳ maximal possible Schmidt rank of  $\psi$

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**Question:** Are multipartite pure states generically ‘very’ or ‘little’ entangled?

→ What is the typical value of the GME for a unit vector  $\psi \in (\mathbf{C}^d)^{\otimes M}$  sampled at random?

## GME of uniformly distributed multipartite pure states

Theorem [Typical  $\varepsilon$  norm of a random unit vector (Aubrun/Szarek)]

There exist constants  $c, C, \alpha_0 > 0$  s.t., for  $\psi \in (\mathbf{C}^d)^{\otimes M}$  a uniformly distributed unit vector,

$$\mathbf{P} \left( c \sqrt{\frac{M \log M}{d^{M-1}}} \leq \|\psi\|_\varepsilon \leq C \sqrt{\frac{M \log M}{d^{M-1}}} \right) \geq 1 - e^{-\alpha_0 d M \log M}.$$

**Consequence:** For  $\psi \in (\mathbf{C}^d)^{\otimes M}$  a uniformly distributed unit vector, when  $d$  or  $M$  is large,  $E(\psi) = (M-1) \log d - \log(M \log M) + O(1)$  with high probability.

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*Proof idea:* Observe that  $\psi \sim g / \|g\|_2$ , where  $g \in (\mathbf{C}^d)^{\otimes M}$  has independent complex Gaussian entries with mean 0 and variance 1.

- By the standard Gaussian concentration inequality:  $\mathbf{P} \left( \|g\|_2 \geq \sqrt{d^M} (1 \pm \varepsilon) \right) \leq e^{-d^M \varepsilon^2}$ .

- Set  $\mathcal{V} := \{ \varphi_1 \otimes \dots \otimes \varphi_M : \varphi_i \in \mathbf{C}^d, \|\varphi_i\|_2 = 1 \}$ , so that  $\mathbf{E} \|g\|_\varepsilon = \mathbf{E} \sup_{\varphi \in \mathcal{V}} |\langle \varphi | g \rangle|$ .

To estimate the latter quantity, use results about suprema of Gaussian processes.

Upper bound: 'small' covering subset of  $\mathcal{V}$ . Lower bound: 'large' separated subset of  $\mathcal{V}$ .

Conclusion:  $\mathbf{E} \|g\|_\varepsilon$  is of order  $\sqrt{dM \log M}$ .

Then show that  $g \mapsto \|g\|_\varepsilon$  also concentrates around its average.

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**Question:** Are ‘interesting’ multipartite pure states really captured by the uniform distribution?



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## 2 Quantifying the typical amount of entanglement/correlations in multipartite pure states

- Typical amount of entanglement in uniformly distributed multipartite pure states
- What about more 'physically relevant' multipartite pure states?

## 3 Typical strength of entanglement criteria for bipartite mixed states

- Entanglement detection of bipartite states
- Estimating the 'size' of sets of states in high-dimensional bipartite systems
- Typical properties of random high-dimensional bipartite states

## 'Physical' states of many-body quantum systems and tensor network states

*Curse of dimensionality* in many-body quantum systems: A system composed of  $M$   $d$ -dimensional subsystems has dimension  $d^M$ , which is exponential in  $M$ .

However, 'physically relevant' states of many-body quantum systems, such as *ground states of gapped local Hamiltonians*, are (conjectured to be) well approximated by so-called *tensor network states (TNS)*, which form a small subset of the global state space (Hastings, Landau/Vazirani/Vidick).

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**Tensor network state on  $(\mathbf{C}^d)^{\otimes M}$ :** Take a graph  $G$  with  $M$  vertices and  $L$  edges.

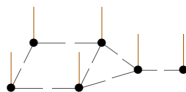
Put at each vertex  $v$  a tensor  $\chi_v \in \mathbf{C}^d \otimes (\mathbf{C}^q)^{\otimes \delta(v)}$  to get a tensor  $\hat{\chi}_G \in (\mathbf{C}^d)^{\otimes M} \otimes (\mathbf{C}^q)^{\otimes 2L}$ .

Contract together the indices of  $\hat{\chi}_G$  associated to a same edge to get a tensor  $\chi_G \in (\mathbf{C}^d)^{\otimes M}$ .

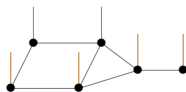
→ If  $\delta(v) \leq \delta$  for all  $v$ , then  $\chi_G$  is described by at most  $Mq^\delta d$  parameters, which is linear in  $M$ .



$G$  with 6 vertices and 7 edges



$$\hat{\chi}_G \in (\mathbf{C}^d)^{\otimes 6} \otimes (\mathbf{C}^q)^{\otimes 14}$$



$$\chi_G \in (\mathbf{C}^d)^{\otimes 6}$$

$d$ -dimensional indices: *physical* indices.  $q$ -dimensional indices: *bond* indices.

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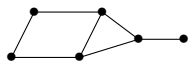
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**Tensor network state on  $(\mathbf{C}^d)^{\otimes M}$ :** Take a graph  $G$  with  $M$  vertices and  $L$  edges.

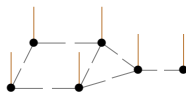
Put at each vertex  $v$  a tensor  $\chi_v \in \mathbf{C}^d \otimes (\mathbf{C}^q)^{\otimes \delta(v)}$  to get a tensor  $\hat{\chi}_G \in (\mathbf{C}^d)^{\otimes M} \otimes (\mathbf{C}^q)^{\otimes 2L}$ .

Contract together the indices of  $\hat{\chi}_G$  associated to a same edge to get a tensor  $\chi_G \in (\mathbf{C}^d)^{\otimes M}$ .

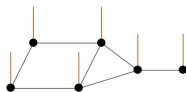
→ If  $\delta(v) \leq \delta$  for all  $v$ , then  $\chi_G$  is described by at most  $Mq^\delta d$  parameters, which is linear in  $M$ .



$G$  with 6 vertices and 7 edges



$$\hat{\chi}_G \in (\mathbf{C}^d)^{\otimes 6} \otimes (\mathbf{C}^q)^{\otimes 14}$$



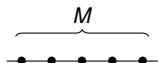
$$\chi_G \in (\mathbf{C}^d)^{\otimes 6}$$

$d$ -dimensional indices: *physical* indices.  $q$ -dimensional indices: *bond* indices.

If the underlying graph  $G$  is 1-dimensional (line or circle),  $\chi_G$  is a *matrix product state (MPS)*.

## A simple model of random translation-invariant MPS

$M$  particles on a circle



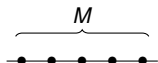
$$|\chi\rangle = \sum_{i=1}^d \sum_{a,a'=1}^q g_{iaa'} |iaa'\rangle$$

Pick a tensor  $\chi \in \mathbf{C}^d \otimes (\mathbf{C}^q)^{\otimes 2}$  whose entries are independent complex Gaussians with mean 0 and variance  $1/dq$ .  
 Repeat it on all sites and contract neighboring  $q$ -dimensional indices.  
 $\rightarrow$  Obtained tensor  $\chi_M \in (\mathbf{C}^d)^{\otimes M}$ : *random translation-invariant MPS with periodic boundary conditions.*

$$|\chi_M\rangle = \sum_{i_1, \dots, i_M=1}^d \left( \sum_{a_1, \dots, a_M=1}^q g_{i_1 a_M a_1} \cdots g_{i_M a_{M-1} a_M} \right) |i_1 \cdots i_M\rangle$$

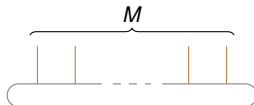
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Associated transfer operator:

$$T : \mathbf{C}^q \otimes \mathbf{C}^q \rightarrow \mathbf{C}^q \otimes \mathbf{C}^q,$$

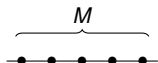
obtained by contracting the  $d$ -dimensional indices of  $\chi$  and  $\bar{\chi}$ .



$$T = \sum_{i=1}^d \left( \sum_{a,a',b,b'=1}^q g_{iaa'} \bar{g}_{ibb'} |ab\rangle \langle a'b'| \right) =: \sum_{i=1}^d G_i \otimes \bar{G}_i$$

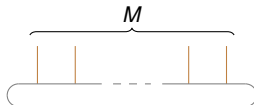
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Pick a tensor  $\chi \in \mathbf{C}^d \otimes (\mathbf{C}^q)^{\otimes 2}$  whose entries are independent complex Gaussians with mean 0 and variance  $1/dq$ . Repeat it on all sites and contract neighboring  $q$ -dimensional indices.  $\rightarrow$  Obtained tensor  $\chi_M \in (\mathbf{C}^d)^{\otimes M}$ : random translation-invariant MPS with periodic boundary conditions.

$$|\chi\rangle = \sum_{i=1}^d \sum_{a,a'=1}^q g_{iaa'} |iaa'\rangle$$



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**Remark:** The parameter  $q$  quantifies the amount of bipartite entanglement: Across any bipartite cut preserving the ordering of subsystems,  $\chi_M$  has Schmidt rank at most  $q^2 \ll d^{M/2}$ .

Now what about genuinely multipartite entanglement?

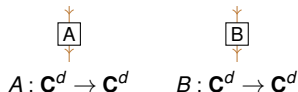
$\hookrightarrow$  area vs volume law

$\rightarrow$  If  $q = 1$ ,  $\chi_M = \chi^{\otimes M}$  is separable. But what can we say for  $q \gg 1$ ?

## Correlations in an MPS

Let  $A, B$  be 1-site observables, i.e. observables on  $\mathbf{C}^d$ .

**Goal:** Quantify the correlations between the outcomes of  $A$  and  $B$ , when performed on 'distant' sites.






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$$A: \mathbf{C}^d \rightarrow \mathbf{C}^d \quad B: \mathbf{C}^d \rightarrow \mathbf{C}^d$$


Compute the value on the MPS  $\chi_M$  of the observable  $A_1 \otimes I_k \otimes B_1 \otimes I_{M-k-2}$ , i.e.

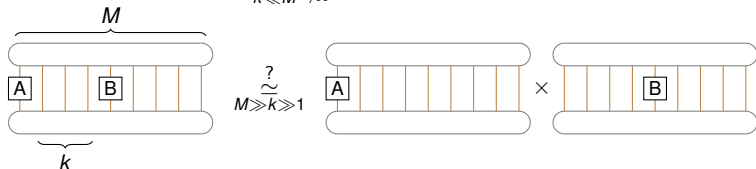
$$v_\chi(A, B, k) := \frac{\langle \chi_M | A_1 \otimes I_k \otimes B_1 \otimes I_{M-k-2} | \chi_M \rangle}{\langle \chi_M | \chi_M \rangle}.$$

Compare it to the product of the values on  $\chi_M$  of  $A_1 \otimes I_{M-1}$  and  $I_{k+1} \otimes B_1 \otimes I_{M-k-2}$ , i.e.

$$v_\chi(A)v_\chi(B) := \frac{\langle \chi_M | A_1 \otimes I_{M-1} | \chi_M \rangle \langle \chi_M | I_{k+1} \otimes B_1 \otimes I_{M-k-2} | \chi_M \rangle}{\langle \chi_M | \chi_M \rangle^2}.$$

Correlations in the MPS  $\chi_M$ :  $\gamma_\chi(A, B, k) := |v_\chi(A, B, k) - v_\chi(A)v_\chi(B)|$ .

**Question:** Do we have  $\gamma_\chi(A, B, k) \xrightarrow[k \ll M \rightarrow \infty]{} 0$ ? And if so, at which speed?



## Exponential decay of correlations in random translation-invariant MPS

Clearly, separability implies no correlation between 1-site observables:

If  $\chi_M = \chi^{\otimes M}$ , then  $\gamma_\chi(A, B, k) = 0$  for any  $k \leq M$  and any observables  $A, B$  on  $\mathbf{C}^d$ .

**Intuition:** In an MPS  $\chi_M$ , the correlations between 1-site observables decay exponentially with the distance separating the sites, i.e. there exist  $C(\chi), \tau(\chi) > 0$  s.t., for any  $k \ll M$  and any observables  $A, B$  on  $\mathbf{C}^d$ ,

$$\gamma_\chi(A, B, k) \leq C(\chi) e^{-\tau(\chi)k} \|A\|_\infty \|B\|_\infty.$$

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### Theorem [Typical correlation length of a random MPS (Lancien/Pérez-García)]

There exist constants  $C, c_0 > 0$  s.t., for  $\chi_M \in (\mathbf{C}^d)^{\otimes M}$  a random translation-invariant MPS,

$$\mathbf{P} \left( \xi(\chi) \leq \frac{C}{\log d} \right) \geq 1 - e^{-c_0 q}.$$

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*Proof idea:* Let  $\lambda_1(T), \lambda_2(T)$  be the two largest eigenvalues of the transfer operator  $T$  and set  $\varepsilon(T) := |\lambda_2(T)|/|\lambda_1(T)|$ . Then,  $\gamma_\chi(A, B, k) \leq C(T)\varepsilon(T)^k \|A\|_\infty \|B\|_\infty$ . So  $\xi(\chi) = 1/|\log \varepsilon(T)|$ .

We can then prove that  $\mathbf{P} \left( |\lambda_1(T)| \geq 1 - \frac{C}{\sqrt{d}} \text{ and } |\lambda_2(T)| \leq \frac{C}{\sqrt{d}} \right) \geq 1 - e^{-c_0 q}$ .

↳ spectral analysis for a non-normal random matrix with tensor product structure

- The amount of correlations in a random MPS is generically small. Is it also the case for the amount of multipartite entanglement?  
→ Can we estimate the GME of a random MPS? Work in progress with I. Nechita...
- What about more complicated models of random MPS, where the random 1-site tensor has some symmetries?
- Can the generic amount of correlations and multipartite entanglement be computed in TNS with a more complicated geometry?  
Typically small correlation length can be proven for random TNS on a 2-dimensional regular lattice (Lancien/Pérez-García), but everything else remains essentially open.

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**Known:** The problem of deciding whether a given multipartite quantum state is entangled or separable (and even just approximate versions of it) is in general computationally hard (Gharibian).  
—→ Solution in practice: Look for necessary conditions to separability, which are easier to check than separability itself, aka *entanglement criteria*.



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Here, we focus on the bipartite case. And we look at two such necessary conditions to separability, which can be efficiently checked and are thus widely used in practice:

- being *positive under partial transposition (PPT)*,
- being *k-extendible*.

### Definition [Partial transposition]

The *partial transposition* (on B) of a state  $\rho_{AB}$  on  $A \otimes B$  is defined as

$$\Gamma_{AB}(\rho_{AB}) := I_A \otimes T_B(\rho_{AB}),$$

where  $I$  denotes the identity map and  $T$  denotes the transposition map.

### Theorem [Necessary condition to separability (Peres)]

On a bipartite Hilbert space  $A \otimes B$ , if a state is separable, then it is positive under partial transposition (PPT).

### Remarks:

- This is obvious since  $\Gamma_{AB}(\sigma_A \otimes \tau_B) = \sigma_A \otimes T_B(\tau_B)$ .
- NSC for separability on  $\mathbf{C}^2 \otimes \mathbf{C}^2$  or  $\mathbf{C}^2 \otimes \mathbf{C}^3$  (Horodecki's). In higher dimensions, there exist PPT entangled states.
- Special instance in the class of separability relaxations built on:  
 $\rho_{AB}$  is separable iff for any positive map  $\Lambda_B$ ,  $I_A \otimes \Lambda_B(\rho_{AB})$  is positive (Horodecki's).

## The $k$ -extendibility criterion

### Definition [ $k$ -extendibility]

Let  $k \geq 2$ . A state  $\rho_{AB}$  on  $A \otimes B$  is  $k$ -extendible (w.r.t.  $B$ ) if there exists a state  $\rho_{AB^k}$  on  $A \otimes B^{\otimes k}$  which is invariant under any permutation of the  $B$  subsystems and s.t.  $\rho_{AB} = \text{Tr}_{B^{k-1}}[\rho_{AB^k}]$ .

### Theorem [Necessary and sufficient condition to separability (Doherty/Parrilo/Spedalieri)]

On a bipartite Hilbert space  $A \otimes B$ , a state is separable iff it is  $k$ -extendible for all  $k \geq 2$ .

### Remarks:

- “ $\rho_{AB}$  separable  $\Rightarrow \rho_{AB}$   $k$ -extendible for all  $k$ ” is obvious since  $\sigma_A \otimes \tau_B = \text{Tr}_{B^{k-1}}[\sigma_A \otimes \tau_B^{\otimes k}]$ .
- “ $\rho_{AB}$   $k$ -extendible for all  $k \Rightarrow \rho_{AB}$  separable” relies on the *quantum de Finetti theorem* (Christandl/König/Mitchison/Renner).
- $\rho_{AB}$   $k$ -extendible  $\Rightarrow \rho_{AB}$   $k'$ -extendible for all  $k' \leq k$ .  
—  $\rightarrow$  Hierarchy of NC for separability, which an entangled state is guaranteed to stop passing at some point (but one cannot tell when a priori).

## Quantifying the strength of separability relaxations

**Problem:** When relaxing the separability constraint to one which is easier to check, how 'rough' is the approximation?

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**Known:** There exist states which are PPT or  $k$ -extendible, and nevertheless 'very' entangled (i.e. far away from the set of separable states in some standard or operational distance measure).  
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### Two possible quantitative strategies:

- Estimate the size of the set of states, either satisfying a given entanglement criterion or being indeed separable.  
—→ Information on how much bigger than the separable set the relaxed set is.
- Characterize when certain random states are with high probability, either violating a given entanglement criterion or indeed entangled.  
—→ Information on how powerful the separability test is to detect entanglement.

## Interlude: asymptotic spectrum of GUE and Wishart matrices

### Definitions [Gaussian Unitary Ensemble and Wishart matrices]

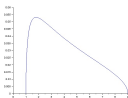
- $G$  is an  $n \times n$  GUE matrix if  $G = (H + H^*)/\sqrt{2}$  with  $H$  an  $n \times n$  matrix whose entries are independent complex Gaussians with mean 0 and variance 1.
- $W$  is an  $(n, s)$ -Wishart matrix if  $W = HH^*$  with  $H$  an  $n \times s$  matrix whose entries are independent complex Gaussians with mean 0 and variance 1.

### Definitions [Semicircular and Marčenko-Pastur distributions]

- $d\mu_{SC(m, \sigma^2)}(x) := \frac{1}{2\pi\sigma^2} \sqrt{4\sigma^2 - (x - m)^2} \mathbf{1}_{[m-2\sigma, m+2\sigma]}(x) dx$ .
  - $d\mu_{MP(\lambda)}(x) := \begin{cases} f_\lambda(x) dx & \text{if } \lambda \geq 1 \\ (1 - \lambda) \delta_0 + \lambda f_\lambda(x) dx & \text{if } \lambda < 1 \end{cases}$ , where
- $$f_\lambda(x) := \frac{\sqrt{(\lambda_+ - x)(x - \lambda_-)}}{2\pi\lambda x} \mathbf{1}_{[\lambda_-, \lambda_+]}(x), \text{ with } \lambda_\pm := (\sqrt{\lambda} \pm 1)^2.$$



$\mu_{SC(0,1)}$



$\mu_{MP(4)}$

**Fact:** For any Hermitian matrix  $M$  on  $\mathbf{C}^n$ , denote by  $\mu_M := \frac{1}{n} \sum_{i=1}^n \delta_{\lambda_i(M)}$  its spectral distribution.

- $(G_n)_{n \in \mathbf{N}}$  sequence of  $n \times n$  GUE matrices:  $(\mu_{G_n/\sqrt{n}})_{n \in \mathbf{N}}$  converges strongly to  $\mu_{SC(0,1)}$ .
- $(W_n)_{n \in \mathbf{N}}$  sequence of  $(n, \lambda n)$ -Wishart matrices:  $(\mu_{W_n/\lambda n})_{n \in \mathbf{N}}$  converges strongly to  $\mu_{MP(\lambda)}$ .

[ Convergence in probability of the spectral distribution and of the extreme eigenvalues. ]

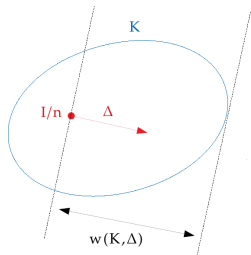
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## Definitions

Let  $K$  be a convex set of states on  $\mathbf{C}^n$  containing the maximally mixed state  $I/n$ .

- For  $\Delta$  a Hermitian on  $\mathbf{C}^n$  s.t.  $\|\Delta\|_2 = 1$ , the *width of  $K$  in the direction  $\Delta$*  is  $w(K, \Delta) := \sup_{\sigma \in K} \text{Tr}(\Delta(\sigma - \frac{I}{n}))$ .
- The *mean width of  $K$*  is the average of  $w(K, \cdot)$  over the Hilbert-Schmidt unit sphere of Hermitians on  $\mathbf{C}^n$ , equipped with the uniform probability measure.  
Equivalently:  $w(K) = \frac{1}{\gamma_n} \mathbf{E} w(K, G)$ , with  $G$  a GUE matrix on  $\mathbf{C}^n$  and  $\gamma_n := \mathbf{E} \|G\|_2 \underset{n \rightarrow +\infty}{\sim} n$ .



The mean width of a set of states is a measure of its size: for any 'reasonable'  $K$ ,  $w(K) \simeq \text{vrad}(K)$ , where  $\text{vrad}(K)$  is the *volume-radius* of  $K$  (i.e. the radius of the Euclidean ball with same volume as  $K$ ).

Computing it amounts to estimating the supremum of some Gaussian process.

## Mean width of the set of separable states

**Observation:** Denote by  $\mathcal{D}$  the set of all states on  $\mathbf{C}^n$ .  $\mathbf{E} \sup_{\sigma \in \mathcal{D}} \text{Tr}(G(\sigma - I/n)) = \mathbf{E} \|G\|_\infty$ . So by the semicircle law, the mean-width of  $\mathcal{D}$  is asymptotically  $2\sqrt{n}/\gamma_n$ , i.e.  $2/\sqrt{n}$ .

### Theorem [Mean width of the set of separable states (Aubrun/Szarek)]

Denote by  $\mathcal{S}$  the set of separable states on  $\mathbf{C}^d \otimes \mathbf{C}^d$ .

There exist universal constants  $c, C > 0$  s.t.  $\frac{c}{d^{3/2}} \leq w(\mathcal{S}) \leq \frac{C}{d^{3/2}}$ .

**Remark:** The mean width of the set of separable states is of order  $1/d^{3/2}$ , hence much smaller than the mean width of the set of all states (of order  $1/d$ ).

→ On high dimensional bipartite systems, most states are entangled.

*Proof idea:* The extreme points of  $\mathcal{S}$  are product pure states, of the form  $|\varphi_1\rangle\langle\varphi_1| \otimes |\varphi_2\rangle\langle\varphi_2|$ . So one can use the same arguments to estimate  $\mathbf{E} \sup_{\sigma \in \mathcal{S}} \text{Tr}(G\sigma)$  for  $G$  a GUE matrix on  $\mathbf{C}^d \otimes \mathbf{C}^d$ , as those used to estimate  $\mathbf{E} \|g\|_\varepsilon$  for  $g$  a standard Gaussian vector in  $\mathbf{C}^d \otimes \mathbf{C}^d$ .

## Mean width of the set of $k$ -extendible states

### Theorem [Mean width of the set of $k$ -extendible states (Lancien)]

Fix  $k \geq 2$  and denote by  $\mathcal{E}_k$  the set of  $k$ -extendible states on  $\mathbf{C}^d \otimes \mathbf{C}^d$ .

$$\text{Then, } w(\mathcal{E}_k) \underset{d \rightarrow \infty}{\sim} \frac{2/\sqrt{k}}{d}.$$

**Remark:** The mean width of the set of  $k$ -extendible states is of order  $1/d$ , hence much bigger than the mean width of the set of separable states.

—→ On high dimensional bipartite systems, the set of  $k$ -extendible states is a very rough approximation of the set of separable states.

*Proof strategy:*  $\sup_{\sigma \in \mathcal{E}_k} \text{Tr}(G(\sigma - I/d^2))$  can be expressed as  $\|\tilde{G}\|_\infty$  for some ‘modified’ GUE matrix  $\tilde{G}$ . So one has to estimate  $\mathbf{E}\|\tilde{G}\|_\infty$ .

This can be done by computing the  $p$ -order moments  $\mathbf{E} \text{Tr} \tilde{G}^p$ , and identifying the limiting spectral distribution (after rescaling by  $d/k$ ): a centered semicircular distribution  $\mu_{SC(0,k)}$ .

The latter has  $2\sqrt{k}$  as upper-edge.

**Note:** Similarly, the mean width of the set of PPT states is of order  $1/d$  (Aubrun/Szarek).

—→ On high dimensional bipartite systems, the set of PPT states is also a very rough approximation of the set of separable states.

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**Setting:** System of interest  $H$  coupled with an environment  $H'$ .

Given a global pure state  $\psi \in H \otimes H'$ , its *marginal* on  $H$   $\rho := \text{Tr}_{H'} |\psi\rangle\langle\psi|$  is a mixed state.

Set  $H \equiv \mathbf{C}^n$  and  $H' \equiv \mathbf{C}^s$ .

—→ Random mixed state model on the system  $\mathbf{C}^n$ , induced by the environment  $\mathbf{C}^s$ :

$\rho = \text{Tr}_{\mathbf{C}^s} |\psi\rangle\langle\psi|$  with  $\psi \in \mathbf{C}^n \otimes \mathbf{C}^s$  a uniformly distributed unit vector.

Equivalent description:  $\rho = \frac{W}{\text{Tr} W}$  with  $W$  a  $(n, s)$ -Wishart matrix.

**Question:** Given  $d \in \mathbf{N}$ , consider  $\rho$  a random state on  $\mathbf{C}^d \otimes \mathbf{C}^d$  induced by an environment  $\mathbf{C}^s$ , for some  $s \in \mathbf{N}$ .

For which values of  $s$  is  $\rho$  typically separable? PPT?  $k$ -extendible?

Hence 2 steps:  $\hookrightarrow$  with probability going to 1 (exponentially) as  $d$  grows

- Identify the range of  $s$  where  $\rho$  is, on average, separable / PPT /  $k$ -extendible.
- Show that the average behavior is generic in high dimension (concentration of measure phenomenon: a sufficiently 'well-behaved' function has an exponentially small probability of deviating from its average as the dimension grows).

### Theorem [Separability threshold for random induced states (Aubrun/Szarek/Ye)]

Let  $\rho$  be a random state on  $\mathbf{C}^d \otimes \mathbf{C}^d$  induced by  $\mathbf{C}^s$ . There exists a threshold  $s_0$  satisfying  $cd^3 \leq s_0 \leq Cd^3 \log^2 d$  for some universal constants  $c, C > 0$  s.t., if  $s < s_0$  then  $\rho$  is typically entangled, and if  $s > s_0$  then  $\rho$  is typically separable.

**Intuition:** If  $s \leq d^2$  then  $\rho$  is uniformly distributed on the set of states of rank at most  $s$ , therefore generically entangled. If  $s \gg d^2$  then  $\rho$  is expected to be close to  $Id/d^2$ , therefore separable.

→ Phase transition between these two regimes?

*Proof idea:* Convex geometry + Comparison of random matrix ensembles (majorization) + Concentration of measure in high dimension.

### Theorem [PPT threshold for random induced states (Aubrun)]

Let  $\rho$  be a random state on  $\mathbf{C}^d \otimes \mathbf{C}^d$  induced by  $\mathbf{C}^s$ .

If  $s < 4d^2$  then  $\rho$  is typically not PPT, and if  $s > 4d^2$  then  $\rho$  is typically PPT.

**Remark:** The threshold environment dimension at which random induced states are generically either PPT or not is of order  $d^2$ , hence much smaller than the one at which they are generically either separable or entangled (of order  $d^3$ ).

→ In the range  $d^2 \lesssim s \lesssim d^3$ , typical entanglement of random induced states is typically not detected by the PPT test.

*Proof strategy:* Everything boils down to characterizing when the partial transposition of a  $(d^2, s)$ -Wishart matrix  $W$  is positive.

This can be done by computing the  $p$ -order moments  $\mathbf{E} \text{Tr} \Gamma(W)^p$ , and identifying the limiting spectral distribution (after rescaling by  $s$ ): a non-centered semicircular distribution  $\mu_{SC(1, d^2/s)}$ .

The latter has positive support iff  $1 - 2\sqrt{d^2/s} \geq 0$  i.e. iff  $s \geq 4d^2$ .

**Note:** Similarly, the threshold environment dimension at which random induced states are generically either  $k$ -extendible or not is of order  $d^2$  (Lancien).

→ In the range  $d^2 \lesssim s \lesssim d^3$ , typical entanglement of random induced states is typically not detected by the  $k$ -extendibility test.

- On high dimensional bipartite systems, the volume of PPT or  $k$ -extendible states is more like the one of all states than like the one of separable states.  
For  $\rho$  a random state on  $\mathbf{C}^d \otimes \mathbf{C}^d$  induced by  $\mathbf{C}^s$ , if  $Cd^2 \leq s \leq cd^3$ , then with high probability  $\rho$  is entangled but this entanglement is not detected by the PPT or  $k$ -extendibility test.  
→ Asymptotic weakness of these necessary conditions for separability.
- Similar features are exhibited by all other known entanglement criteria.
- What about generalizations to the multipartite setting?  
First need: Entanglement criteria which are not designed only for bipartite states.  
Proposal: Construction of such criteria through tensor norms (Jivulescu/Lancien/Nechita).  
→ Can their typical performance be quantified?



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