

Approximate simulation and bisimulation relations for fuzzy automata

Stefan Stanimirović (Joint work with Ivana Micić and Miroslav Ćirić)
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The idea of bisimulations

- ▶ *Simulations* and *bisimulations* - notions used in various settings: modal logic, concurrency theory, set theory, automata theory (fuzzy, weighted, probabilistic, nondeterministic automata).
- ▶ They are usually defined as **binary relations** over the domains of two structures, which can be automata.
- ▶ Two states of automata are in a relation iff they can perform the same actions to reach bisimilar states.
- ▶ When the two structures are the same, the largest bisimulation is an **equivalence relation**, with applications in factorization, separating expressive powers of modal and related logics, concept learning in description logics, etc.

Weighted automata (over semirings):

- ▶ P. Buchholz. “Bisimulation relations for weighted automata”. In: *Theoretical Computer Science* 393.1 (2008), pp. 109–123
- ▶ N. Damjanović, M. Ćirić, and J. Ignjatović. “Bisimulations for weighted automata over an additively idempotent semiring”. In: *Theoretical Computer Science* 534 (2014), pp. 86–100
- ▶ S. Stanimirović, A. Stamenković, and M. Ćirić. “Improved algorithms for computing the greatest right and left invariant Boolean matrices and their application”. In: *Filomat* 33.9 (2019), pp. 2809–2831

Fuzzy automata (over lattice-ordered structures):

- ▶ M. Ćirić et al. “Bisimulations for fuzzy automata”. In: *Fuzzy Sets and Systems* 186.1 (2012), pp. 100–139
- ▶ Y. Cao, G. Chen, and E. E. Kerre. “Bisimulations for Fuzzy-Transition Systems”. In: *IEEE Transactions on Fuzzy Systems* 19.3 (2011), pp. 540–552
- ▶ L. A. Nguyen and D. X. Tran. “Computing Fuzzy Bisimulations for Fuzzy Structures under the Gödel Semantics”. In: *IEEE Transactions on Fuzzy Systems* (2020)

Motivation for approximate bisimulations

- ▶ There are cases when there is no bisimulation between two automata, but they are “**more or less**” bisimilar. (Bisimilar **to some degree**.)
- ▶ There are cases when no factorization of an automaton can be made via bisimulations, but we want to factorize the automaton, even if the factorized automaton is not strictly language-equivalent to the starting automaton, but equivalent **to some degree**. (Especially of interest for large automata.)
- ▶ The presence of the *vagueness* in this problem leads to the observation of *approximate bisimulations* for **fuzzy automata**.

Heyting algebras, fuzzy sets and fuzzy relations

- ▶ $\mathcal{H} = (H, \wedge, \vee, \rightarrow, 0, 1)$ - **Heyting Algebra**, if:

(H1) $(H, \wedge, \vee, 0, 1)$ is a bounded distributive lattice,

(H2) For every $x, y, z \in H$, $x \wedge y \leq z$ iff $x \leq y \rightarrow z$.

($x \rightarrow y$ is called a **relative pseudo-complement** of x in y .)

$x \leftrightarrow y = (x \rightarrow y) \wedge (y \rightarrow x)$ - **biimplication**.

- ▶ **Gödel structure**: $H = [0, 1]$, $x \wedge y = \min(x, y)$, $x \vee y = \max(x, y)$,

$$x \rightarrow y = \begin{cases} 1, & x \leq y \\ y, & \text{otherwise} \end{cases}, \quad x \leftrightarrow y = \begin{cases} 1, & x = y, \\ \min(x, y), & \text{otherwise} \end{cases}.$$

- ▶ $\alpha : A \rightarrow H$ - a **fuzzy subset** of A .

H^A - the set of all fuzzy subsets of A .

- ▶ $\varphi : A \times A \rightarrow H$ - a **fuzzy relation** on A .

$H^{A \times A}$ - the set of all fuzzy relations on A .

The degree of subsethood and equality

- ▶ $\alpha, \beta \in H^A$ - two fuzzy sets.
- ▶ $S(\alpha, \beta) \in H$ - the **degree of subsethood of α in β** :

$$S(\alpha, \beta) = \bigwedge_{a \in A} \alpha(a) \rightarrow \beta(a).$$

$S(\alpha, \beta)$ - a truth degree of “for each $a \in A$: if a belongs to α then a belongs to β .”

- ▶ $E(\alpha, \beta) \in H$ - the **degree of equality of α and β** :

$$E(\alpha, \beta) = \bigwedge_{a \in A} \alpha(a) \leftrightarrow \beta(a).$$

$E(\alpha, \beta)$ - a truth degree of “for each $a \in A$: a belongs to α iff a belongs to β .”

Approximate simulations

- ▶ $\mathcal{A} = (A, \sigma^A, \delta^A, \tau^A)$ and $\mathcal{B} = (B, \sigma^B, \delta^B, \tau^B)$ - fuzzy automata;
- ▶ $\lambda \in H$ - a scalar from a complete Heyting algebra (**an approximation degree**).
- ▶ A fuzzy relation $\varphi \in H^{A \times B}$ is:

λ -approximate forward simulation, if:

$$\begin{aligned} S(\sigma^A, \sigma^B \circ \varphi^{-1}) &\geq \lambda, \\ S(\varphi^{-1} \circ \delta_x^A, \delta_x^B \circ \varphi^{-1}) &\geq \lambda \\ &\text{for every } x \in X, \\ S(\varphi^{-1} \circ \tau^A, \tau^B) &\geq \lambda, \end{aligned}$$

λ -approximate backward simulation, if:

$$\begin{aligned} S(\tau^A, \varphi \circ \tau^B) &\geq \lambda, \\ S(\delta_x^A \circ \varphi, \varphi \circ \delta_x^B) &\geq \lambda \\ &\text{for every } x \in X, \\ S(\sigma^A \circ \varphi, \sigma^B) &\geq \lambda. \end{aligned}$$

Approximate bisimulations

► A fuzzy relation $\varphi \in H^{A \times B}$ is:

λ -approximate forward bisimulation, if: φ and φ^{-1} are λ -approximate forward simulations,

λ -approximate forward-backward bisimulation (λ -FB), if:

$$\begin{aligned} E(\sigma^A, \sigma^B \circ \varphi^{-1}) &\geq \lambda, \\ E(\varphi^{-1} \circ \delta_x^A, \delta_x^B \circ \varphi^{-1}) &\geq \lambda \\ &\text{for every } x \in X, \\ E(\varphi^{-1} \circ \tau^A, \tau^B) &\geq \lambda, \end{aligned}$$

λ -approximate backward bisimulation, if: φ and φ^{-1} are λ -approximate backward simulations,

λ -approximate backward-forward bisimulation (λ -BF), if:

$$\begin{aligned} E(\tau^A, \varphi \circ \tau^B) &\geq \lambda, \\ E(\delta_x^A \circ \varphi, \varphi \circ \delta_x^B) &\geq \lambda \\ &\text{for every } x \in X, \\ E(\sigma^A \circ \varphi, \sigma^B) &\geq \lambda. \end{aligned}$$

An example

Let $\mathcal{A} = (A, \sigma^A, \delta^A, \tau^A)$ and $\mathcal{B} = (B, \sigma^B, \delta^B, \tau^B)$ be fuzzy automata over the Gödel structure \mathcal{G} and $X = \{x, y\}$ with $|A| = 3, |B| = 2$ and

$$\begin{aligned}\sigma^A &= [0 \quad 0 \quad 1], & \tau^A &= [1 \quad 1 \quad 1]^T, \\ \delta_x^A &= \begin{bmatrix} 1 & 0.3 & 0.4 \\ 0.5 & 1 & 0.3 \\ 0.4 & 0.6 & 0.7 \end{bmatrix}, & \delta_y^A &= \begin{bmatrix} 0.5 & 0.6 & 0.2 \\ 0.6 & 0.3 & 0.4 \\ 0.7 & 0.7 & 1 \end{bmatrix}, \\ \sigma^B &= [0.8 \quad 0.8], & \tau^B &= [1 \quad 1]^T, \\ \delta_x^B &= \begin{bmatrix} 1 & 0.6 \\ 0.6 & 0.7 \end{bmatrix}, & \delta_y^B &= \begin{bmatrix} 0.6 & 0.6 \\ 0.7 & 1 \end{bmatrix}.\end{aligned}$$

An example (cont.)

- ▶ For $\lambda = 1$ - no simulations nor bisimulations;
- ▶ $\lambda = 0.8$:

$$\varphi^{\text{FS}} = \begin{bmatrix} 1 & 0.7 \\ 1 & 0.7 \\ 0.6 & 1 \end{bmatrix}, \quad \varphi^{\text{BS}} = \begin{bmatrix} 1 & 0.7 \\ 1 & 0.7 \\ 0.7 & 1 \end{bmatrix}.$$

- ▶ $\lambda = 0.7$:

$$\psi^{\text{FS}} = \begin{bmatrix} 1 & 1 \\ 1 & 1 \\ 0.6 & 1 \end{bmatrix}, \quad \psi^{\text{BS}} = \begin{bmatrix} 1 & 1 \\ 1 & 1 \\ 1 & 1 \end{bmatrix}, \quad \psi^{\text{BFB}} = \begin{bmatrix} 1 & 0.6 \\ 1 & 0.6 \\ 1 & 1 \end{bmatrix}.$$

Theorem

Let \mathcal{A} and \mathcal{B} be fuzzy automata, $\lambda \in H$, and $\varphi \in H^{A \times B}$. Then:

1. If φ is **λ -approximate simulation**, then

$$S([\mathcal{A}], [\mathcal{B}]) \geq \lambda,$$

2. If φ is **λ -approximate bisimulation**, then

$$E([\mathcal{A}], [\mathcal{B}]) \geq \lambda.$$

Theorem

Let \mathcal{A} and \mathcal{B} be fuzzy automata such that there exists **at least one** λ -approximate (bi)simulation between \mathcal{A} and \mathcal{B} . Then there exists **the greatest** λ -approximate (bi)simulation, which is also a **partial fuzzy function**.

Theorem

Let \mathcal{A} be a fuzzy automaton and $\lambda \in H$. Then:

- ▶ There exists **the greatest** λ -approximate (bi)simulation on \mathcal{A} ,
- ▶ The greatest λ -approximate **simulation** on \mathcal{A} is a **fuzzy quasi-order** on A ,
- ▶ The greatest λ -approximate **bisimulation** on \mathcal{A} is a **fuzzy equivalence** on A .
- ▶ If φ is λ -approximate bisimulation fuzzy equivalence on \mathcal{A} , or λ -approximate backward simulation fuzzy quasi-order on \mathcal{A} , then

$$E([\mathcal{A}], [\mathcal{A}/\varphi]) \geq \lambda.$$

An example

Let \mathcal{H} be the Gödel structure, and let $\mathcal{A} = (A, \sigma^A, \delta^A, \tau^A)$ be a fuzzy automaton over \mathcal{H} and $X = \{x\}$ with $|A| = 6$ and

$$\begin{aligned}\sigma^A &= [0.2 \quad 0.7 \quad 0.4 \quad 0.9 \quad 0.7 \quad 0.7], \\ \tau^A &= [0.6 \quad 0.8 \quad 0.8 \quad 0.7 \quad 0.8 \quad 0.6]^T, \\ \delta_x^A &= \begin{bmatrix} 0.4 & 1 & 0 & 0.6 & 0.2 & 1 \\ 0.5 & 0.5 & 0.8 & 0.8 & 0.4 & 0.9 \\ 0.6 & 0.6 & 0.2 & 0.2 & 0.4 & 0.7 \\ 0 & 0.5 & 1 & 0.9 & 0.2 & 0.9 \\ 0.4 & 0.9 & 0.9 & 0.5 & 0.8 & 0.4 \\ 0.7 & 0.8 & 0.2 & 0.4 & 0 & 0.2 \end{bmatrix}. \\ [[\mathcal{A}]](u) &= \begin{cases} 0.7, & \text{if } u = e, \\ 0.8, & \text{otherwise.} \end{cases}\end{aligned}$$

An example (cont.)

- ▶ Case $\lambda = 1$: **No reduction** can be made via any type of bisimulation.
- ▶ Case $\lambda = 0.7$. Then the greatest **λ -approximate backward simulation**:

$$\psi = \psi^{\text{BS}} = \begin{bmatrix} 1 & 1 & 1 & 1 & 1 & 1 \\ 0.2 & 1 & 0.4 & 1 & 1 & 1 \\ 0.2 & 1 & 1 & 1 & 1 & 1 \\ 0.2 & 1 & 0.4 & 1 & 1 & 1 \\ 0.2 & 1 & 0.4 & 1 & 1 & 1 \\ 0.2 & 1 & 0.4 & 1 & 1 & 1 \end{bmatrix}$$

which means that \mathcal{A}/ψ has **3 different states**. It can be verified that $[[\mathcal{A}/\psi]](u) = 0.8$, for every $u \in X^*$.

An example (cont.)

- ▶ Case $\lambda = 0.7$. Then the greatest λ -*approximate forward bisimulation*

$$\varphi = \varphi^{\text{FB}} = \begin{bmatrix} 1 & 0.6 & 0.6 & 0.6 & 0.6 & 1 \\ 0.6 & 1 & 0.6 & 1 & 0.6 & 0.6 \\ 0.6 & 0.6 & 1 & 0.6 & 0.6 & 0.6 \\ 0.6 & 1 & 0.6 & 1 & 0.6 & 0.6 \\ 0.6 & 0.6 & 0.6 & 0.6 & 1 & 0.6 \\ 1 & 0.6 & 0.6 & 0.6 & 0.6 & 1 \end{bmatrix},$$

which means that \mathcal{A}/φ has **4 different states**. It can be verified that $[[\mathcal{A}/\varphi]](u) = 0.8$, for every $u \in X^*$.

- ▶ There are *no reductions* via other types of λ -approximate bisimulations.

- ▶ More details on approximate bisimulations:
S. Stanimirović, I. Micić, and M. Ćirić. “Approximate Bisimulations for Fuzzy Automata over Complete Heyting Algebras”. In: *IEEE Transactions on Fuzzy Systems* (2020)
- ▶ Future work:
 - ▶ Approximate (bi)simulations which are partial fuzzy functions and uniform fuzzy relations,
 - ▶ Faster algorithms for computing the greatest approximate (bi)simulations,
 - ▶ Approximate (bi)simulations for fuzzy/weighted automata over other structures.

Thank you for your attention!