## Kleene and Büchi for Weighted Forest Languages over M-Monoids

Frederic Dörband<sup>1</sup> Technische Universität Dresden, February 10, 2020

We recall forests as finite tuples of trees. This syntactic extension of trees already occurred in  $[^{2,3,4}]$ . We define weighted forest automata (wfa) as an extension of weighted tree automata (wta). Our definition generalizes the unweighted forest automata introduced in  $[^1]$  to the weighted case, while also generalizing the weighted forest automata introduced in  $[^3]$  from commutative semirings to M-monoids. In comparison with  $[^2]$ , our paper is a generalization from semirings to M-monoids, but also a restriction from hedge languages to forest languages.

A wfa is syntactically similar to a wta, however the semantics of wfa allow for entire forests to be processed (instead of single trees). The essence of our paper is a theorem stating that our automaton model generates only so-called rectangular weighted forest languages. That is, for every recognizable weighted forest language  $\varphi$ , there are recognizable weighted tree languages (the so-called rectangular components of  $\varphi$ ), whose horizontal concatenation equals  $\varphi$ .

Using this rectangularity property, we then infer both, a Kleene-like result and a Büchi-like result for weighted forest languages. On the one hand, this generalizes the Kleene-like result for weighted tree automata from  $[^5]$  to the case of forests and the Kleene-like result for unweighted forest automata  $[^1]$  to the case of M-monoids. On the other hand, this generalizes the Büchi-like result for weighted tree automata from  $[^6]$  to the case of forests.

For our Kleene-like result, we introduce rational forest expressions, which are horizontal products of rational tree expressions and for our Büchi-like result, we introduce forest M-expressions, which are horizontal products of tree M-expressions. We ultimately ask the question, whether the definitions for rational forest expressions and forest Mexpressions can be done inductively. However, we see that straightforward inductive definitions are not possible due to ambiguity in the choice of rectangular components.

<sup>&</sup>lt;sup>1</sup>Research of this author is supported by the DFG Research Training Group 1763 "QuantLA"

<sup>&</sup>lt;sup>2</sup>L. Straßburger. "A Kleene theorem for forest languages". In: Language and automata theory and applications. Vol. 5457. Lecture Notes in Comput. Sci. Springer, Berlin, 2009, pp. 715–727. URL: https://doi.org/10.1007/978-3-642-00982-2\_61.

<sup>&</sup>lt;sup>3</sup>C. Mathissen. Weighted Automata and Weighted Logics over Tree-like Structures. Citeseer, 2009.

<sup>&</sup>lt;sup>4</sup>F. Dörband. "A Kleene Theorem for Weighted Forest Automata". MA thesis. Technische Universität Dresden, 2019.

<sup>&</sup>lt;sup>5</sup>Z. Fülöp, A. Maletti, and H. Vogler. "A Kleene theorem for weighted tree automata over distributive multioperator monoids". In: *Theory Comput. Syst.* 44 (2009), pp. 455–499.

<sup>&</sup>lt;sup>6</sup>Zoltán Fülöp, Torsten Stüber, and Heiko Vogler. "A Büchi-like theorem for weighted tree automata over multioperator monoids". In: *Theory of Computing Systems* 50.2 (2012), pp. 241–278.