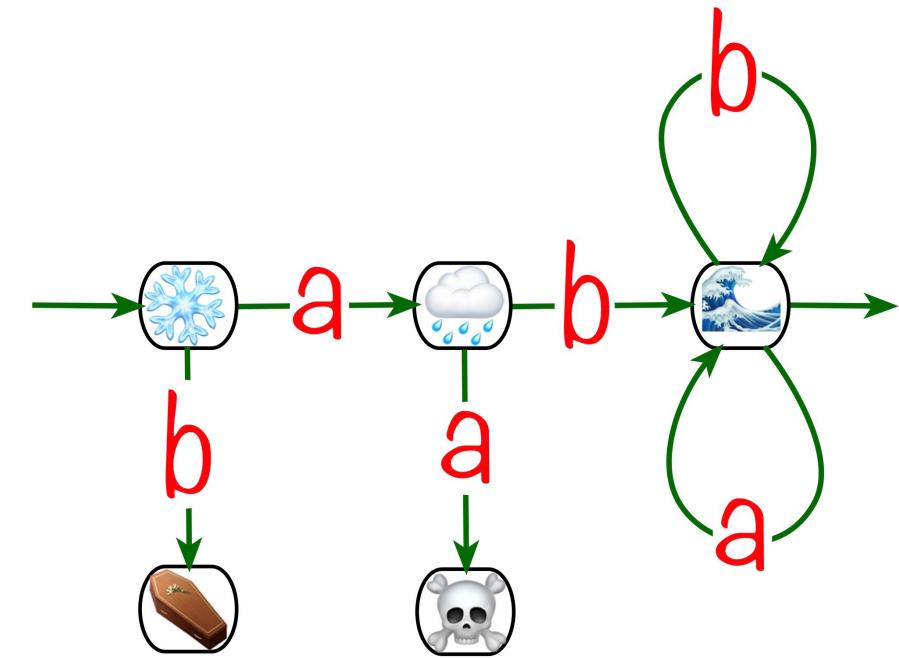
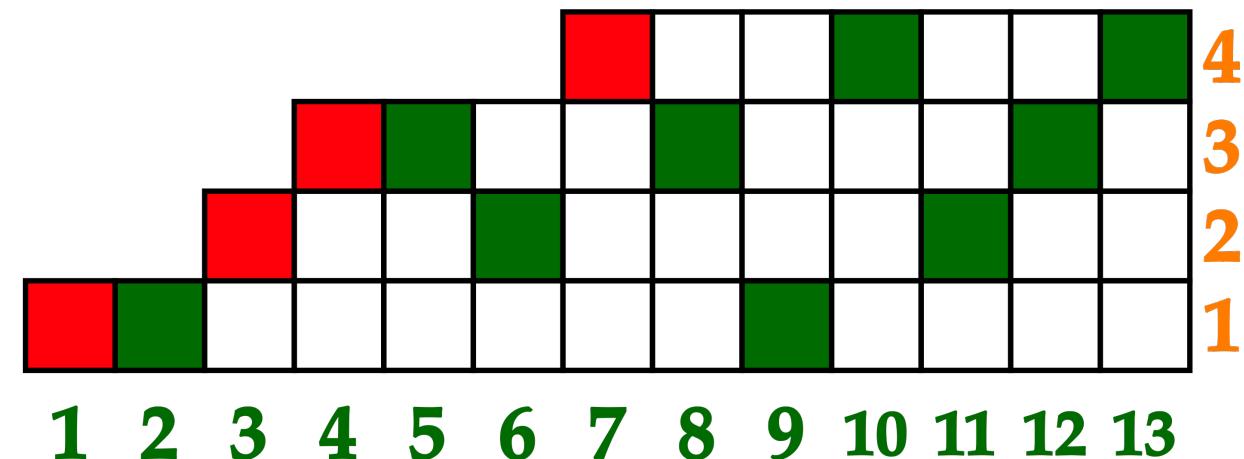


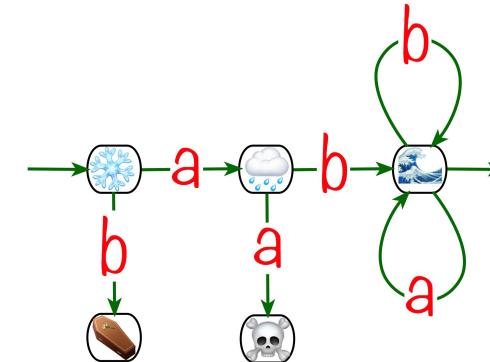
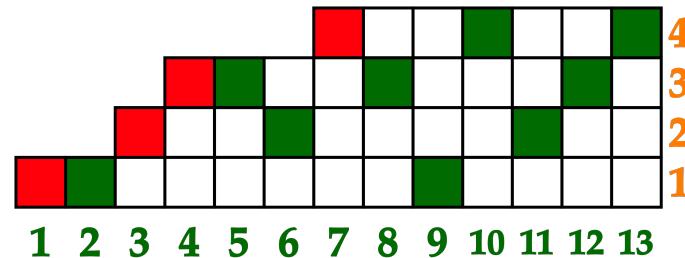
# The impatient collector

Anis Amri, Philippe Chassaing



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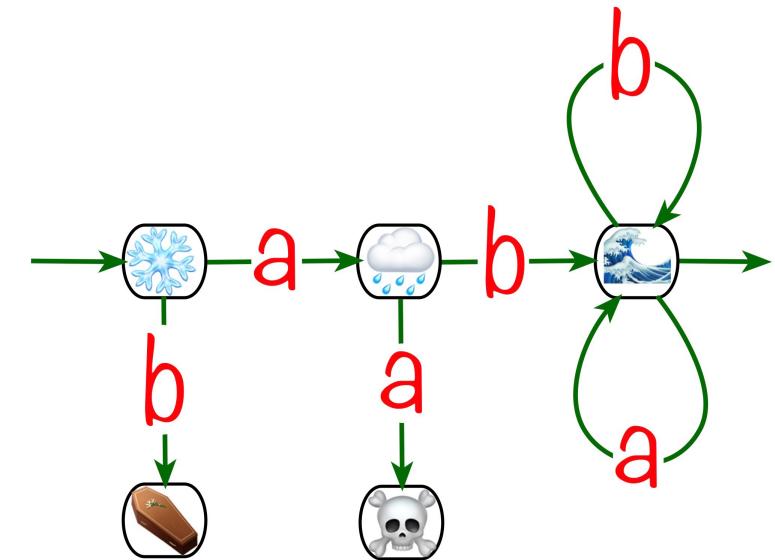
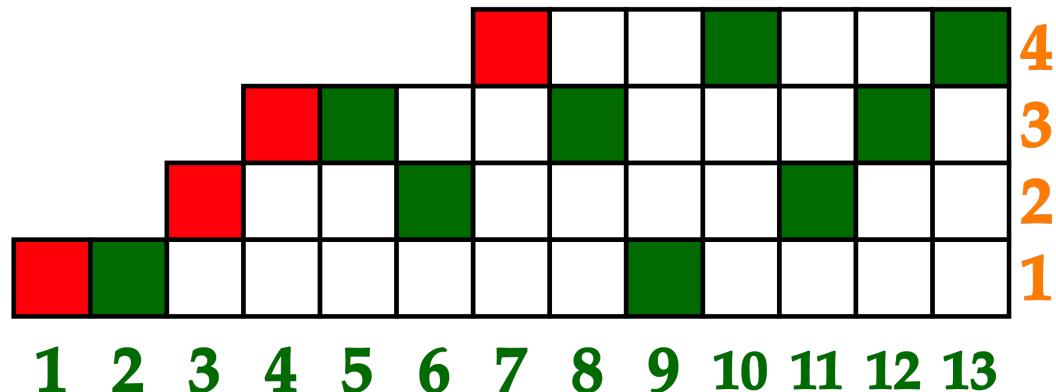


following papers by :

- Cyril Nicaud, Frédérique Bassino, Julien Clément, Andrea Sportiello, et al.
- Devroye and Cai
- Addario-Berry et al.

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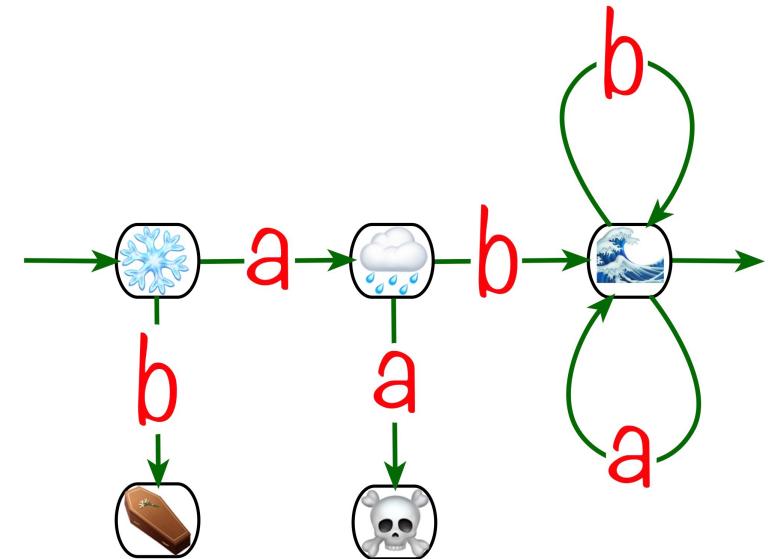
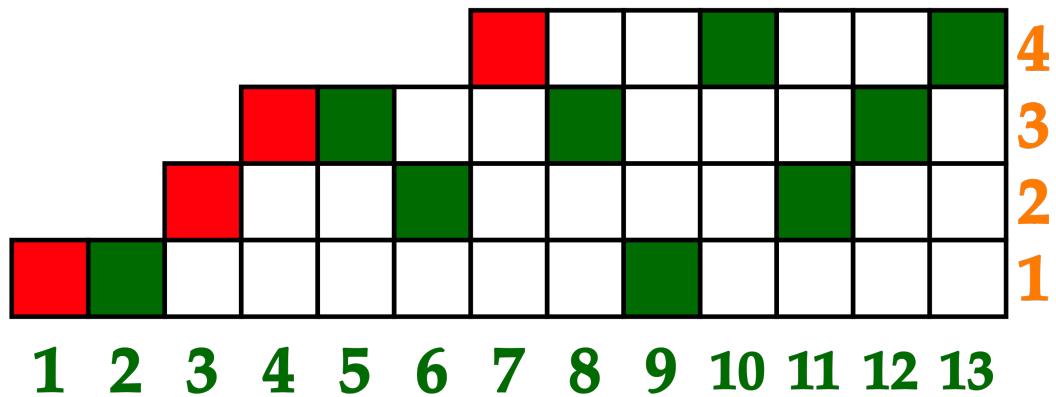
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Byproduct 1: study of a variation of the coupon collector problem.

# The impatient collector

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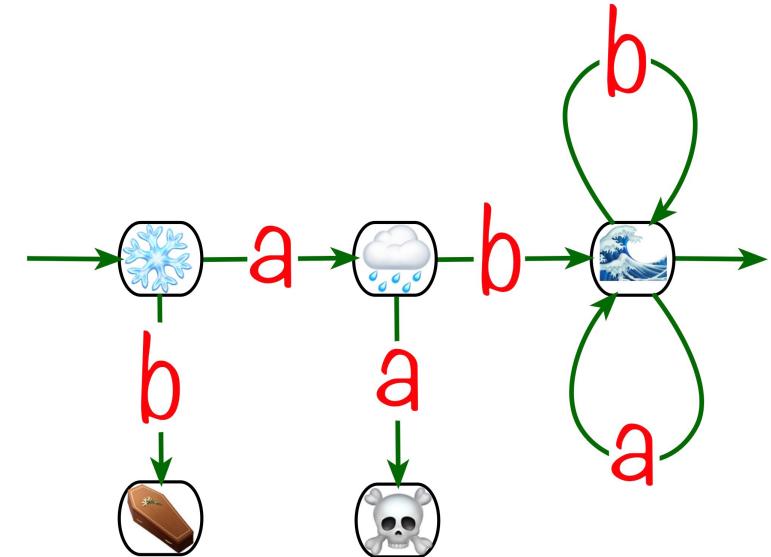
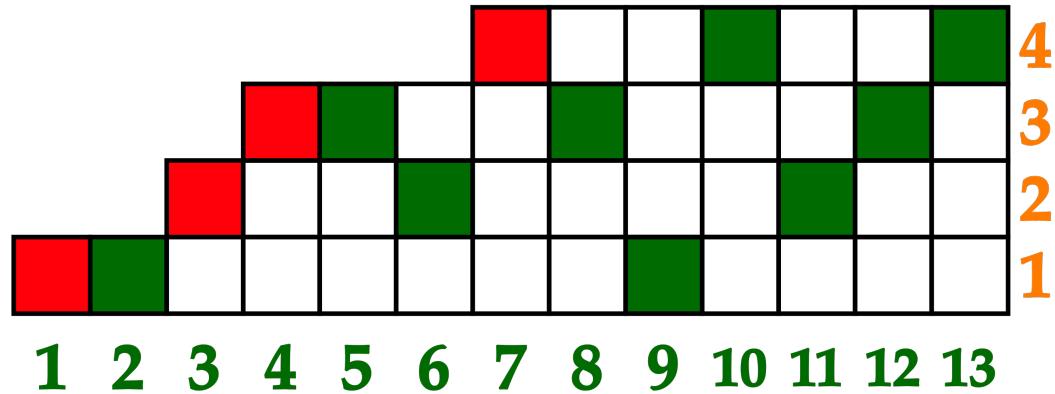


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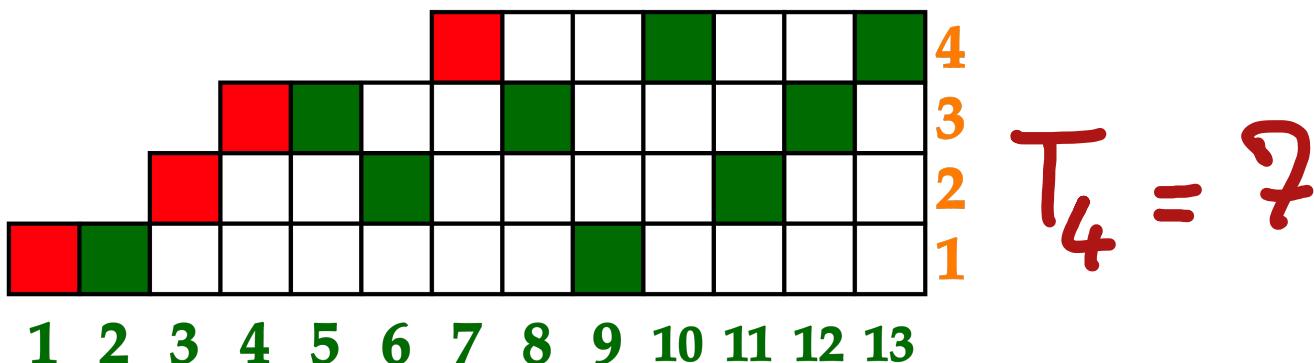


Byproduct 1: study of a variation of the coupon collector problem.

Byproduct 2: refined asymptotic analysis of the Stirling numbers of the 2<sup>nd</sup> kind.

Byproduct 3: asymptotic analysis of large automata

# The unhurried collector



$n$  coupons (here  $n=4$ )

$\omega = \omega_1 \omega_2 \omega_3 \dots \in [1, n]^N$   $\omega_i$ 's are i.i.d. uniform.

$\omega_{[a,b]} = \omega_a \omega_{a+1} \dots \omega_b$  is a factor of  $\omega$ .

$T_n(\omega)$ : the smallest  $k$  s.t.

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$$\mathbb{E}[T_n] = n + l_n \approx n \ln(n).$$

Erdős Rényi 1961 If  $T_n = n \ln n + n Z_n$ , then:

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## ON A CLASSICAL PROBLEM OF PROBABILITY THEORY

by

P. ERDŐS and A. RÉNYI

We consider the following classical "urn-problem". Suppose that there are  $n$  urns given, and that balls are placed at random in these urns one after the other. Let us suppose that the urns are labelled with the numbers  $1, 2, \dots, n$  and let  $\xi_j$  be equal to  $k$  if the  $j$ -th ball is placed into the  $k$ -th urn. We suppose that the random variables  $\xi_1, \xi_2, \dots, \xi_N, \dots$  are independent, and  $\mathbf{P}(\xi_j = k) = \frac{1}{n}$  for  $j = 1, 2, \dots$  and  $k = 1, 2, \dots, n$ . By other words each

ball may be placed in any of the urns with the same probability and the choices of the urns for the different balls are independent. We continue this process so long till there are at least  $m$  balls in every urn ( $m = 1, 2, \dots$ ). What can be said about the number of balls which are needed to achieve this goal?

We denote the number in question (which is of course a random variable) by  $v_m(n)$ . The "dixie cup"-problem considered in [1] is clearly equivalent with the above problem. In [1] the mean value  $\mathbf{M}(v_m(n))$  of  $v_m(n)$  has been evaluated (here and in what follows  $\mathbf{M}()$  denotes the mean value of the random variable in the brackets) and it has been shown that

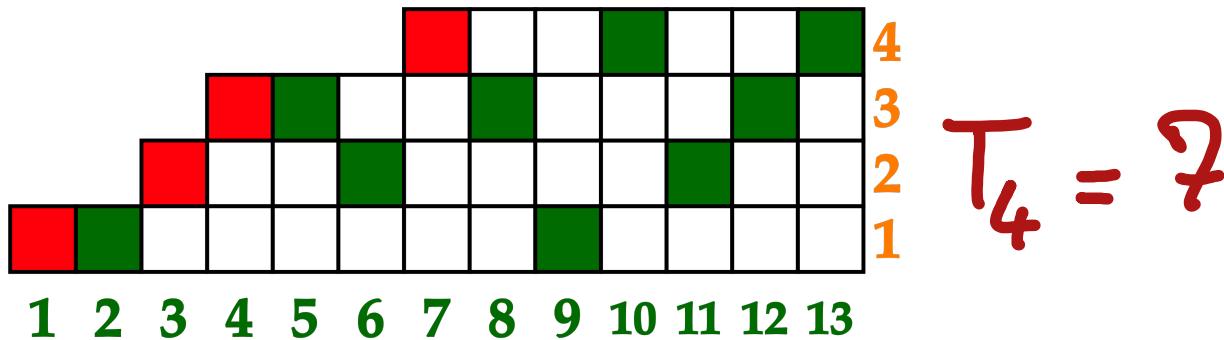
$$(1) \quad \mathbf{M}(v_m(n)) = n \log n + (m-1) n \log \log n + n \cdot C_m + o(n)$$

where  $C_m$  is a constant, depending on  $m$ . (The value of  $C_m$  is not given in [1]).

In the present note we shall go a step further and determine asymptotically the probability distribution of  $v_m(n)$ ; we shall prove that for every real  $x$  we have

$$(2) \quad \lim_{n \rightarrow +\infty} \mathbf{P}\left(\frac{v_m(n)}{n} < \log n + (m-1) \log \log n + x\right) = \exp\left(-\frac{e^{-x}}{(m-1)!}\right).$$

# The unhurried collector



$$T_4 = 7$$

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on aura donc à très-peu près pour l'expression du nombre  $i$  de tirages, après lesquels la probabilité que tous les numéros seront sortis est  $\frac{1}{k}$ ,

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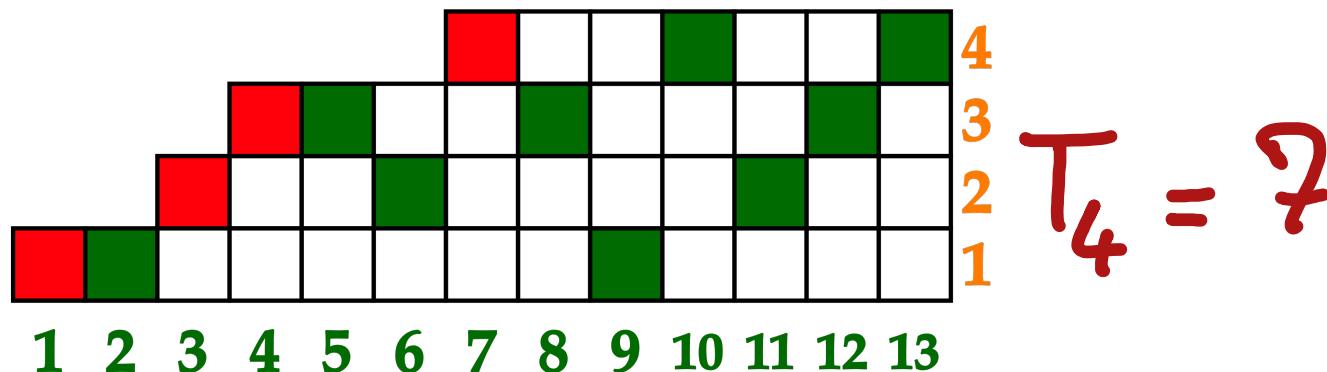
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□

cites sources such as Bernstein or Renyi (1960's)

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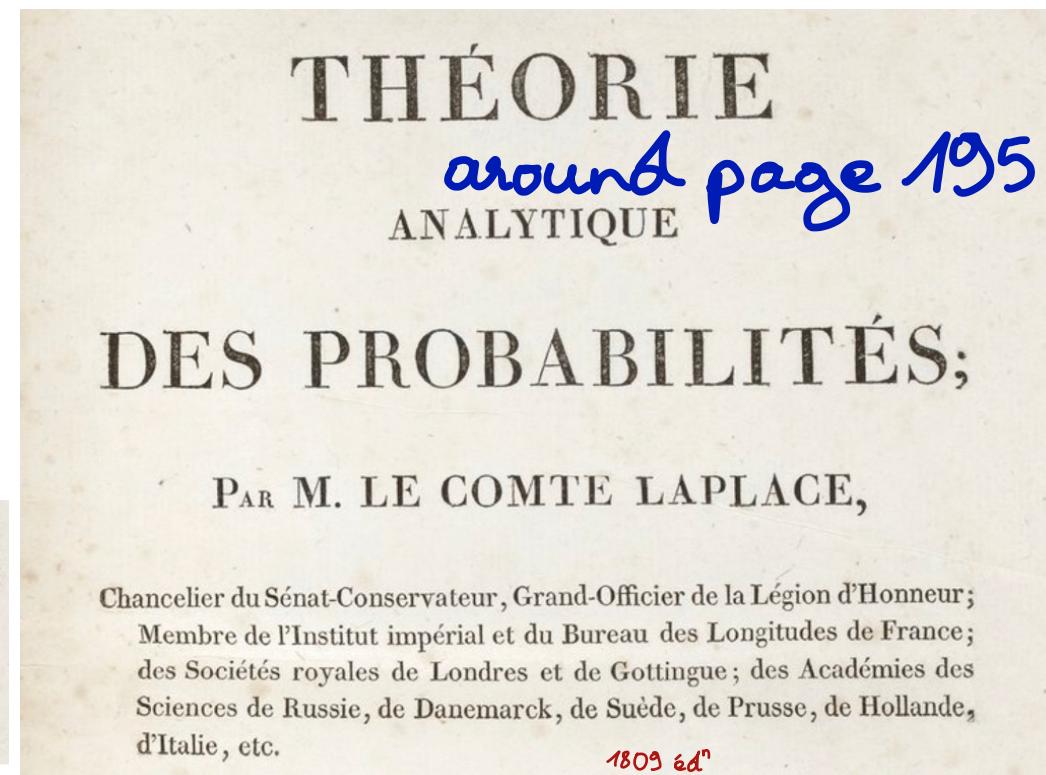
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quantile  $f_n$  of  $T_n$

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quantile  $f_n$  of Gumbel



## The completion curve

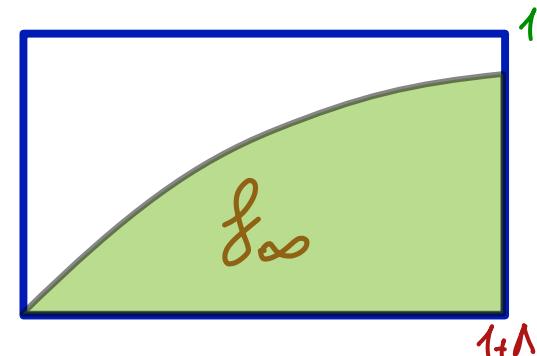
$\omega = \omega_1, \omega_2, \omega_3, \dots$  is "uniform" on  $[1, n]^{\mathbb{N}}$

$$\zeta_n(t, \omega) = \frac{1}{n} \# \{ \omega_k, 1 \leq k \leq t_n \}$$

is the completion curve.

Proposition. (unhurried case)

$$\lim_n \mathbb{E} [\zeta_n(t)] = 1 - e^{-t} \\ \equiv \zeta_\infty(t).$$



# The completion curve

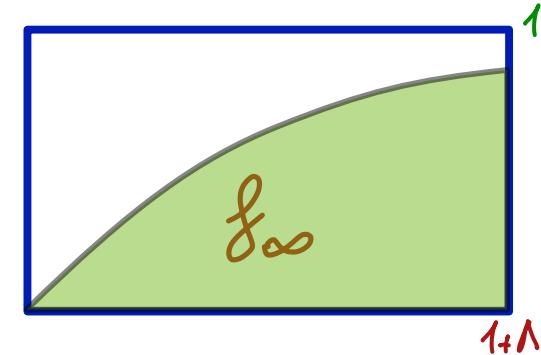
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Proposition. (unhurried case)

$$\lim_n \mathbb{E} [\zeta_n(t)] = 1 - e^{-t} \stackrel{=} \zeta_\infty(t). \quad \leftarrow \quad \mathbb{E} [\ ] = 1 - (1 - \frac{1}{n})^{[tn]}$$



Question (impatient case)

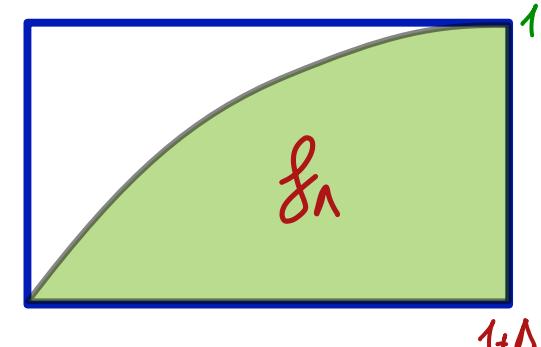
Given that  $T_n \leq (1+\lambda)n, (\lambda > 0)$

asymptotic behaviour of  $\zeta_n(t)$  ??

i.e. description of

$$\zeta(\lambda, t) = \lim_n \mathbb{E} [\zeta_n(t) \mid T_n \leq (1+\lambda)n] ??$$

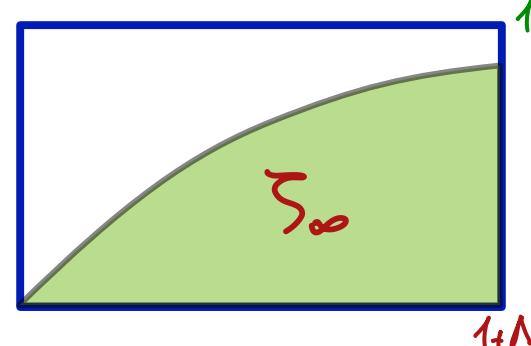
Note:  $\mathbb{P}(T_n \leq (1+\lambda)n) \approx C(\lambda) e^{-n J(\lambda)}$  in which  $J(\lambda) \xrightarrow{+∞} 0$ .



## The completion curve

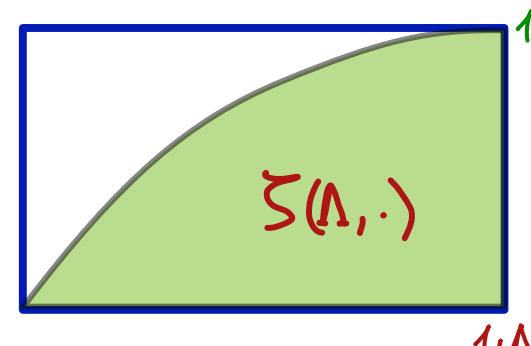
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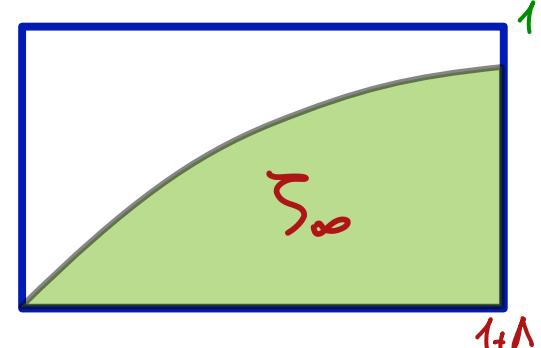


\* exists,

# The completion curve

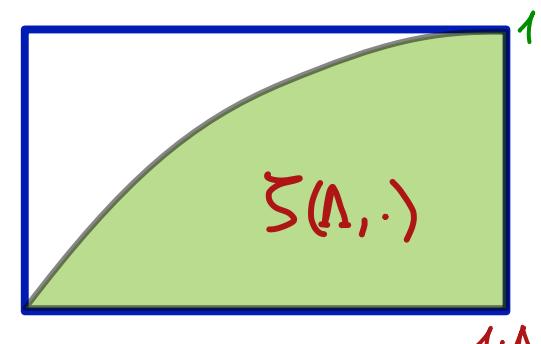
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Question: asymptotic behaviour of  $\zeta_n(t) ??$

$$\zeta(\lambda, t) = \lim_n \mathbb{E} [\zeta_n(t) \mid T_n < (1+\lambda)n]$$



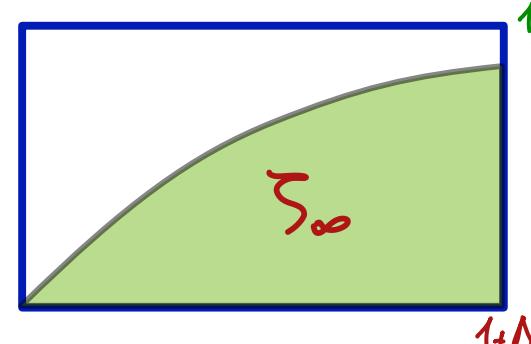
\* exists,

\* is solution of an ODE,

# The completion curve

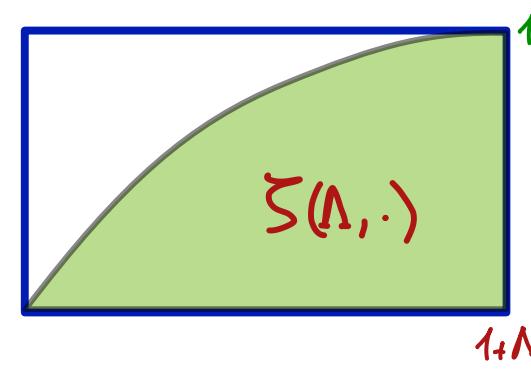
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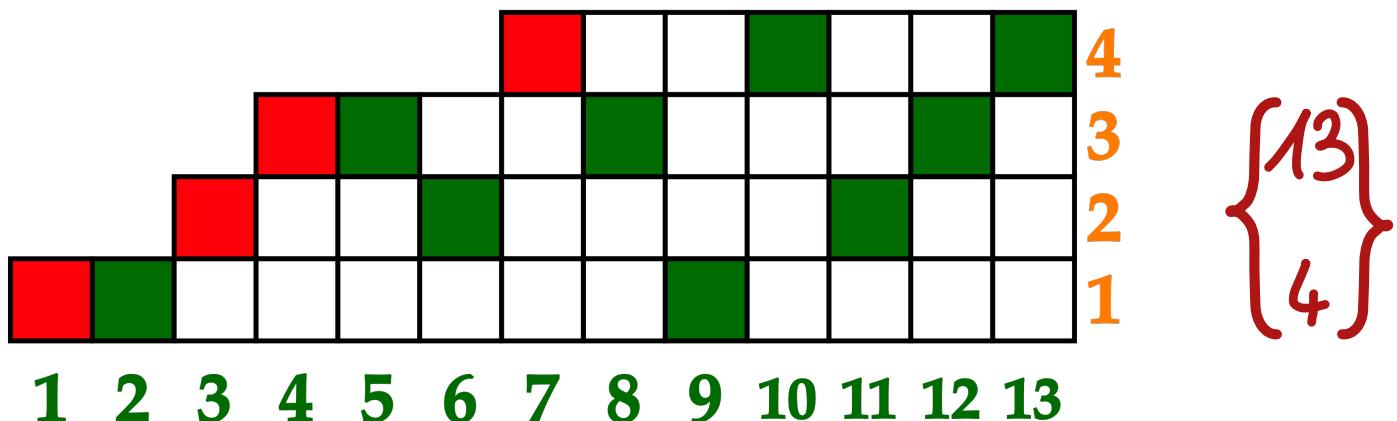
- \* exists,
- \* is solution of an ODE,
- \* satisfies  $\zeta(\lambda, t) = \lim_n \zeta_n(t)$  in probability ...

# Stirling numbers of the 2nd kind $\left\{ \begin{matrix} m \\ l \end{matrix} \right\}$

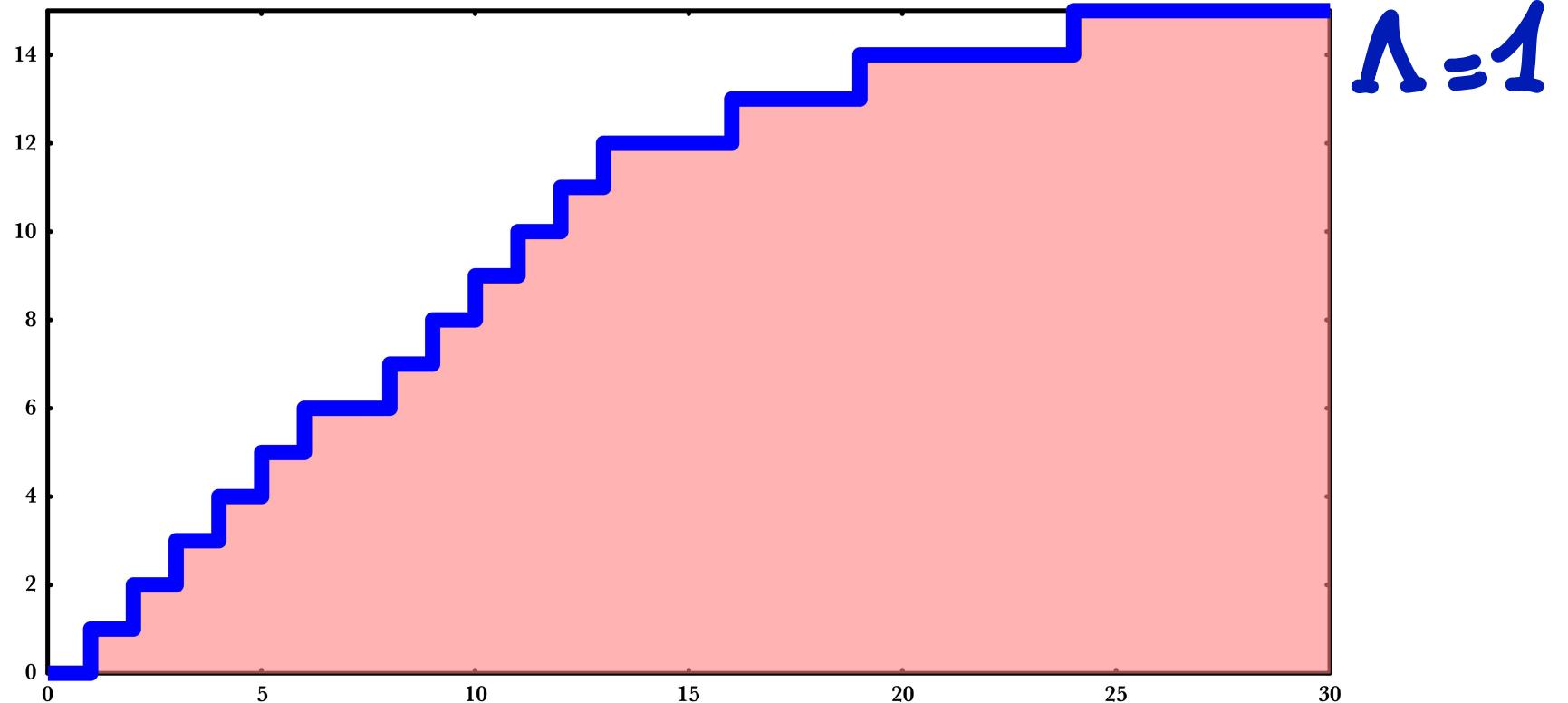
$\left\{ \begin{matrix} m \\ l \end{matrix} \right\}$  : number of partitions of  $[1, m]$  in  $l$  parts,  $l \leq m$ .

$\left\{ \begin{matrix} m \\ l \end{matrix} \right\} \times (l!)$  : number of surjections from  $[1, m]$  to  $[1, l]$ .

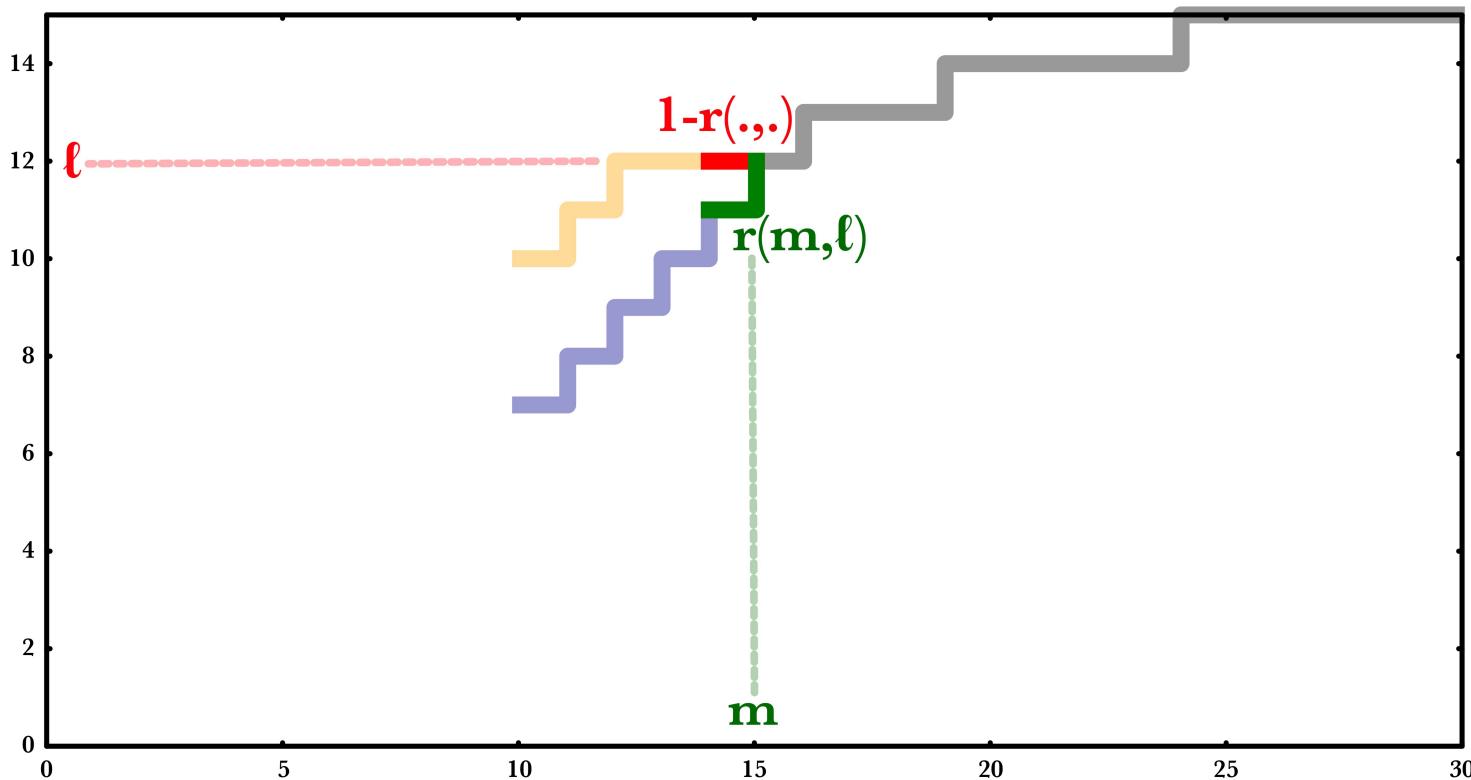
$\left\{ \begin{matrix} m \\ l \end{matrix} \right\} \times (l!) \times l^{-m} = P(T_l < m)$ .



# A Markov chain



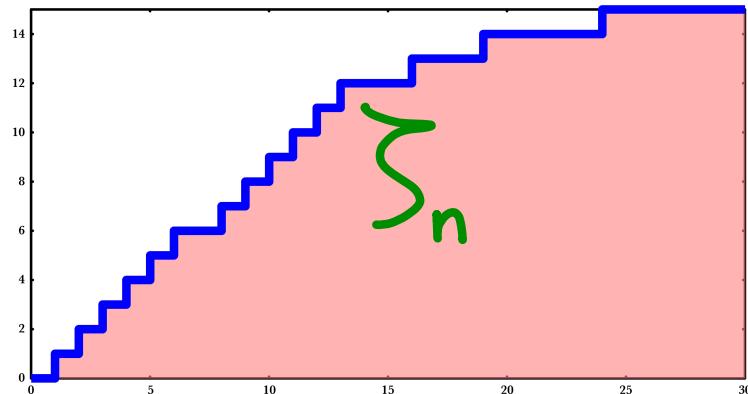
# A Markov chain



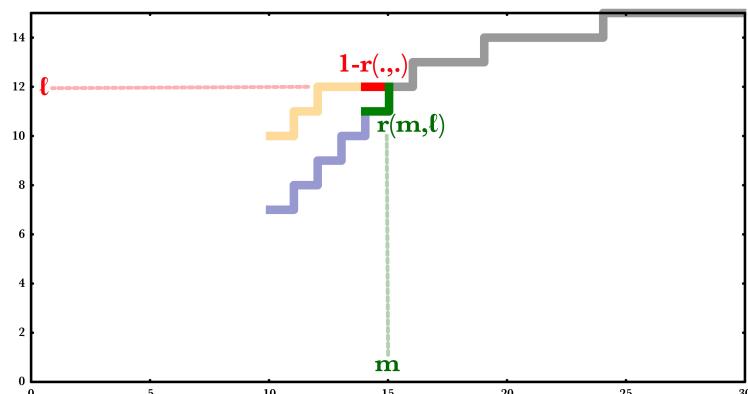
$r(m,l) = \frac{\binom{m-1}{l-1}}{\binom{m}{l}}$  is the probability of ↘

$1 - r(m,l) = \frac{\binom{m-1}{l}}{\binom{m}{l}}$  is the probability of —

# A Markov chain



$r(m,l) = \frac{\binom{m-1}{l-1}}{\binom{m}{l}}$  is the probability of ↘



$1 - r(m,l) = \frac{\binom{m-1}{l} e}{\binom{m}{l} e}$  is the probability of —

The asymptotic shape  $\zeta(n, \cdot)$  is given by

the asymptotic behaviour of  $\binom{m}{l}$

# Asymptotics for $\binom{m}{\ell}$

$$\binom{m}{\ell} \sim \psi(m, \ell) = \frac{m!(e^\xi - 1)^\ell}{\ell! \xi^m \sqrt{2\pi m \left(1 - \frac{m}{\ell} e^{-\xi}\right)}}.$$

Good 1961

$$\lambda(m, \ell) = \lambda = \frac{m - \ell}{\ell}$$

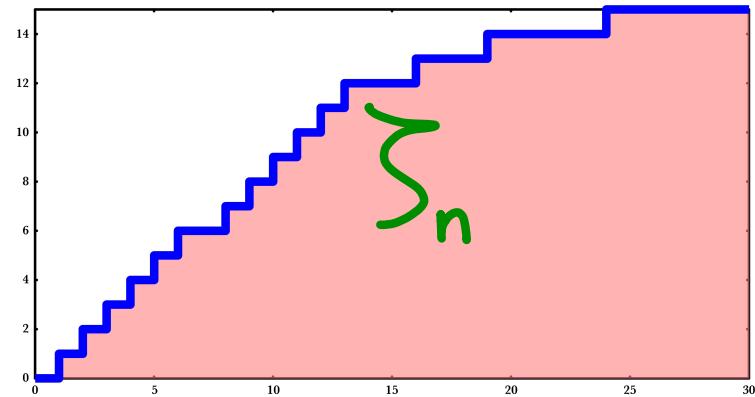
$$\xi(m, \ell) = \xi = 1 + \lambda + W_0 \left( - (1 + \lambda) e^{-1-\lambda} \right)$$

$$v = \frac{(\lambda + 1)(\xi - \lambda)}{2}$$

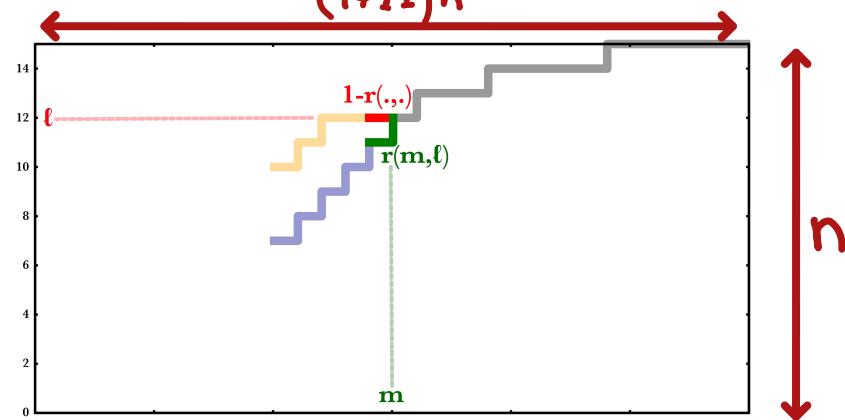
$W_0$ : Lambert function

$$\xi = (1 - e^{-\xi})(1 + \lambda)$$

# A Markov chain



$(1+\Lambda)n$



$$\left. \begin{aligned} \lambda &= \frac{m}{\ell} - 1 \\ \xi &= (1 - e^{-\xi})(1 + \lambda) \end{aligned} \right\} \rightarrow r(m, \ell) \simeq \exp(-\xi(\lambda)) = f(\lambda)$$

$\rightarrow \xi(\lambda, \cdot)$  is sol<sup>n</sup> of

$$r(m, \ell) = \frac{\binom{m-1}{\ell-1}}{\binom{m}{\ell}}$$

is the probability of ↗

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Good 1961

$$\begin{cases} y' = f\left(\frac{y}{2} - 1\right) \\ y(1 + \lambda) = 1 \end{cases}$$

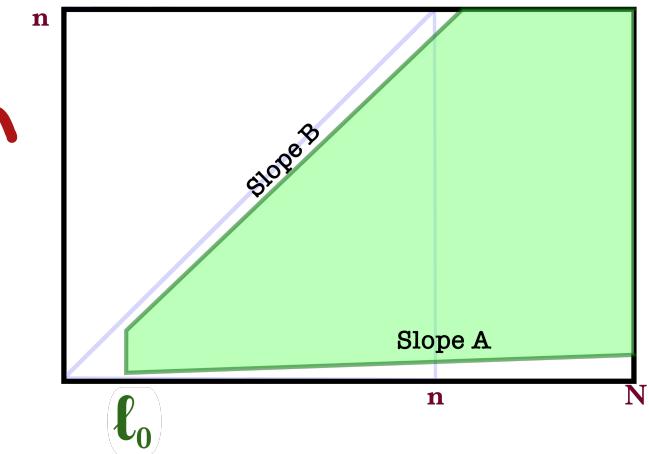
# Proof: Euler scheme + Azuma inequality

$$\left| \frac{\binom{m}{\ell} - \psi(m, \ell)}{\psi(m, \ell)} \right| \leq \frac{C_2(\ell_0, \delta)}{\ell}$$

Slope B =  $\frac{1}{1+\delta}$   
 Slope A =  $\frac{\delta}{1+\delta}$

for  $(m, \ell)$  in

$$|r(m, \ell) - \rho(\lambda)| \leq \frac{C_1(\ell_0, \delta)}{\ell}$$



**Theorem.** For any  $a > 0$ , and for  $n_0$  large enough, there exists  $C = C(n_0, a) > 0$  such that, for  $n \geq n_0$ ,

$$\mathbb{P}_{N,n} \left( \sup_{[a, 1+\Lambda(N, n)]} |\zeta_n - \zeta(\Lambda(N, n), \cdot)| \geq Cn^{-1/3} \right) \leq n^{1/3} e^{-\ln^2 n/2}.$$

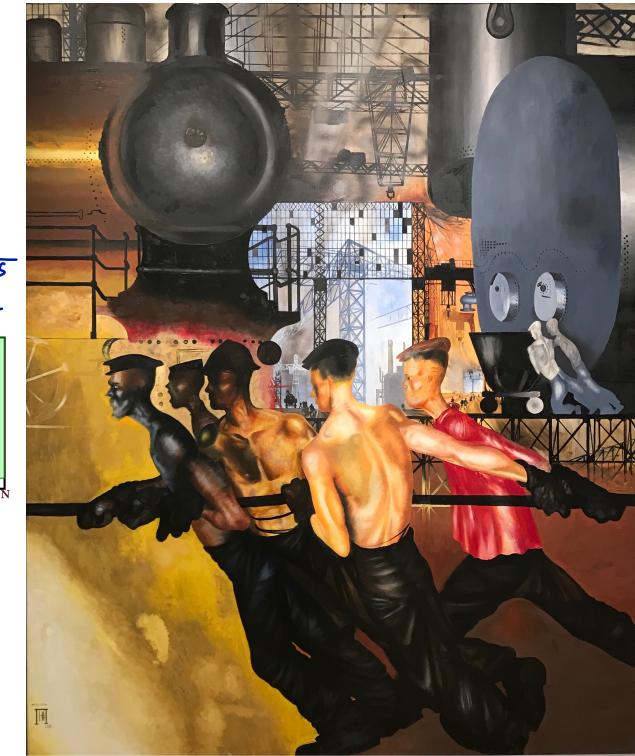
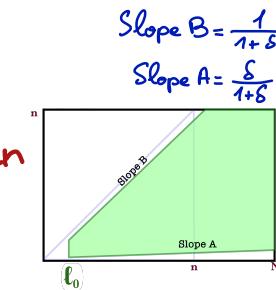
# Proof: Euler scheme + Azuma inequality & saddle point methods



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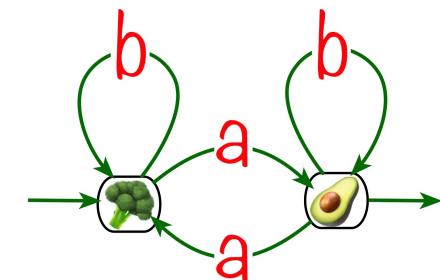


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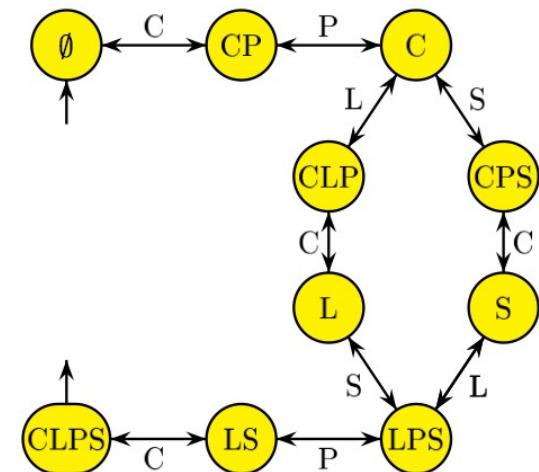
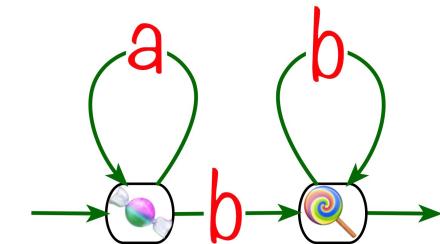
# Automata : some terminology

\* state: vertex



\* alphabet ( $A$ ): set of edges' labels.

\* starting state, final states



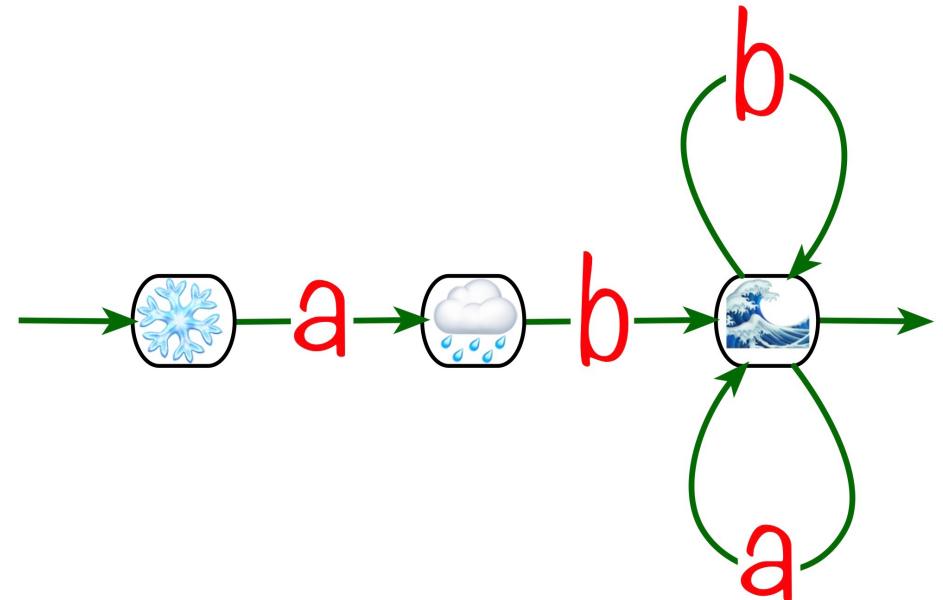
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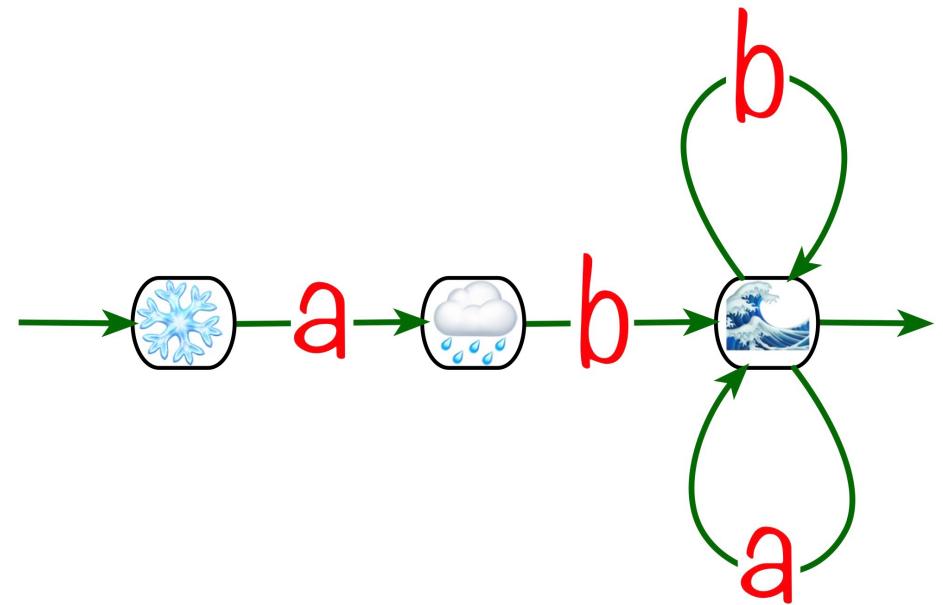
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\* complete automata: " " " exactly  
one outgoing edge ....



# Automata : some terminology

\* state: vertex

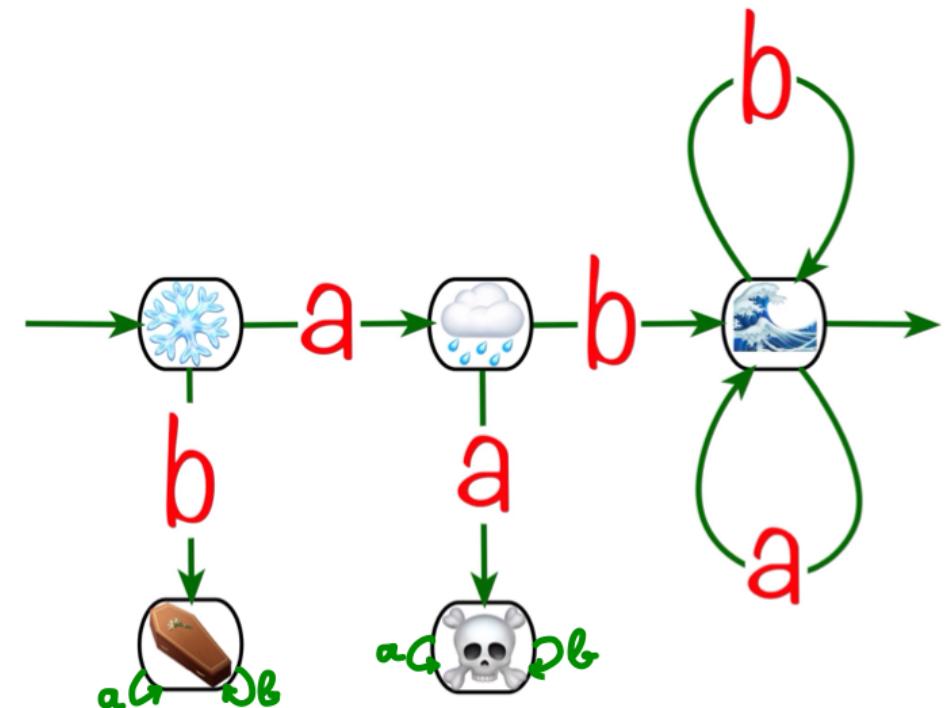
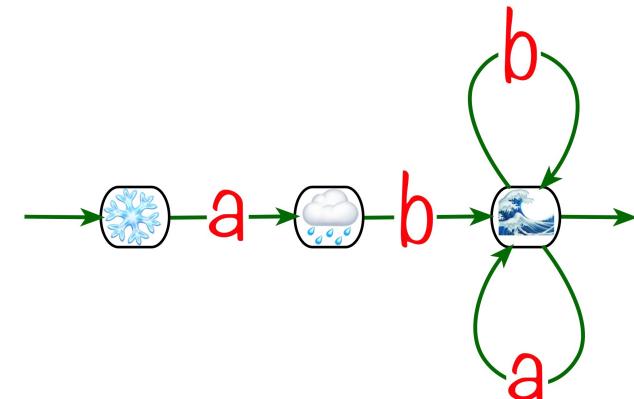
\* alphabet ( $A$ ): set of edges' labels.

\* starting state, final states

\* deterministic automata:  $\forall a \in A, \forall$  vertex,  
at most one outgoing edge labeled  $a$ .

\* complete automata: " " " exactly  
one outgoing edge ....

\* accessible automata: each state can be  
reached, from the starting state.



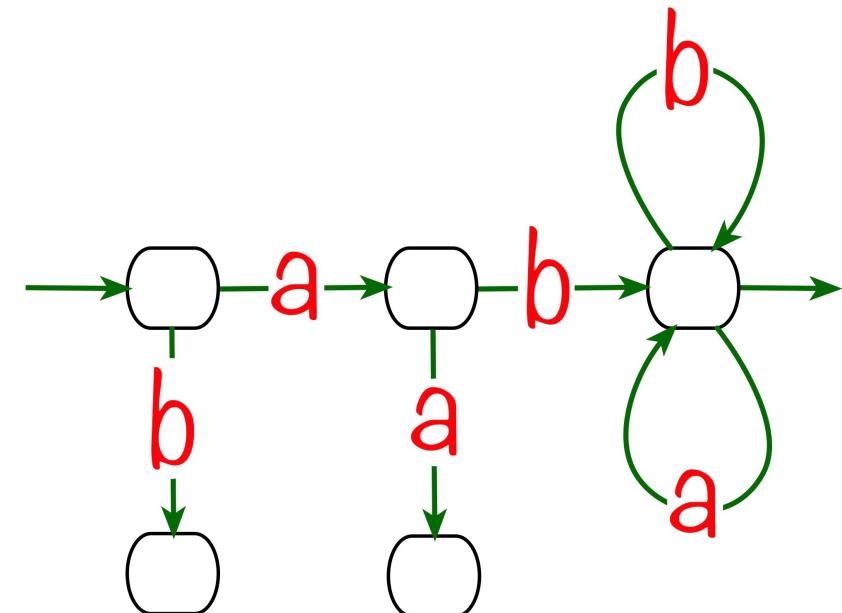
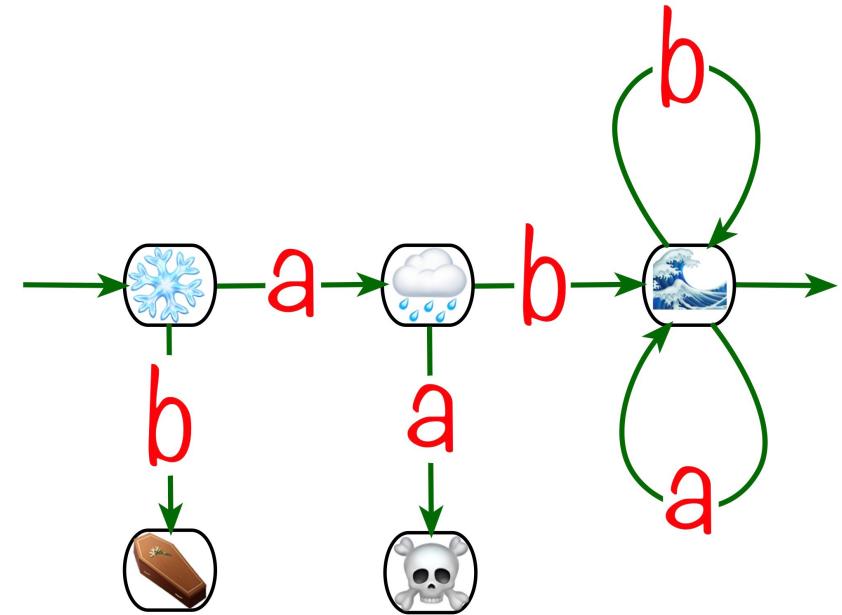
# Relation automata - coupon collector

← input for Moore's algorithm

$\mathcal{A}_{n,k}$ : set of the deterministic complete  
accessible automata (DCAA)  
with  $n$  vertices and a  $k$ -letters alphabet  $\mathcal{A}$ .

Koršunov: # $\mathcal{A}_{n,k} \approx \binom{k^{n+1}}{n} n! \times C(k,n)$

in which  $\lim_n C(k,n) = C(k) \in ]0,1[$ .



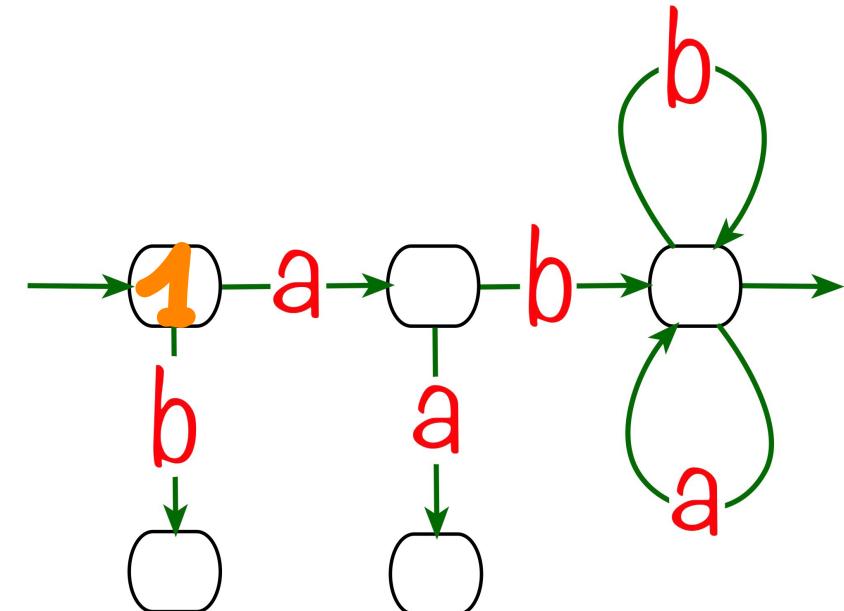
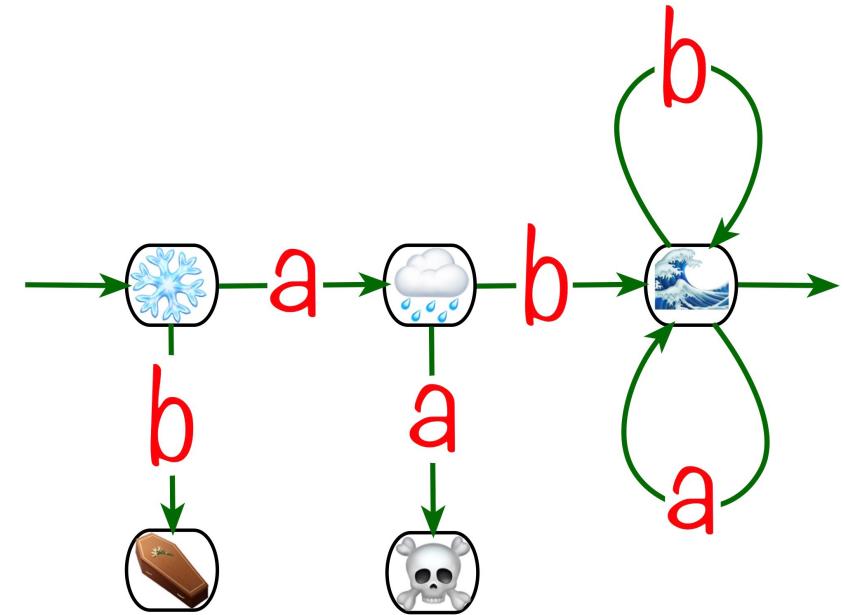
Provided an alphabetic order on  $\mathcal{A}$ , breadth-first search  
provides an ordering of the vertices (if accessibility).

# Relation automata - coupon collector

$\mathcal{Q}_{n,k}$ : set of the deterministic complete  
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Koršunov:  $\# \mathcal{Q}_{n,k} \approx \binom{k^{n+1}}{n} n! \times C(k,n)$

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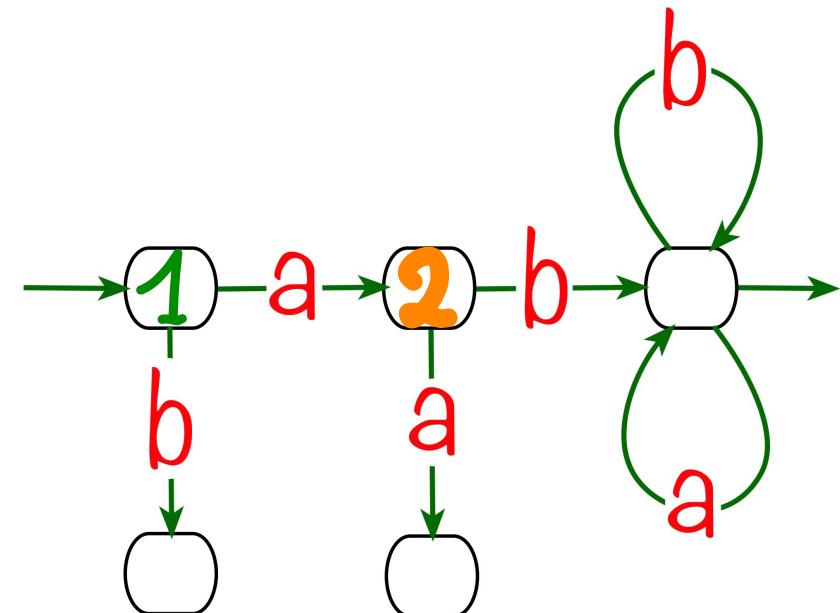
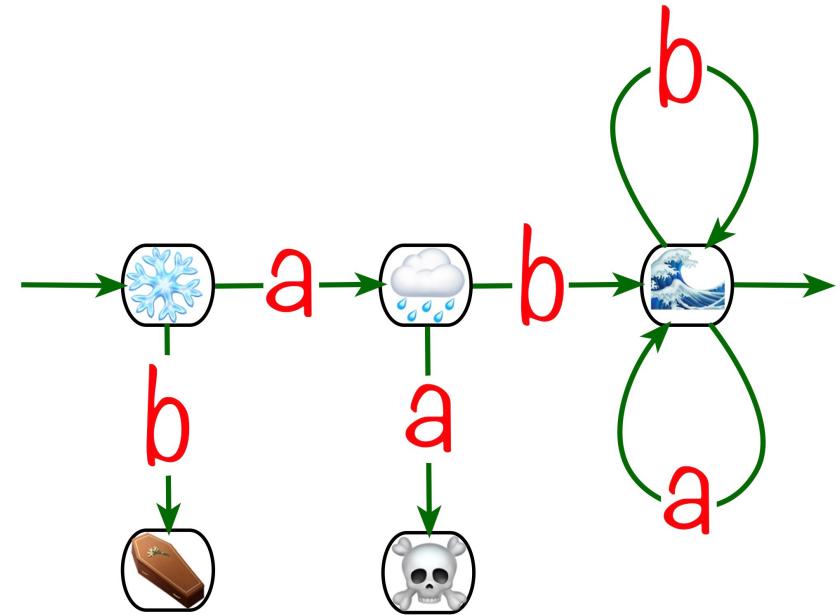
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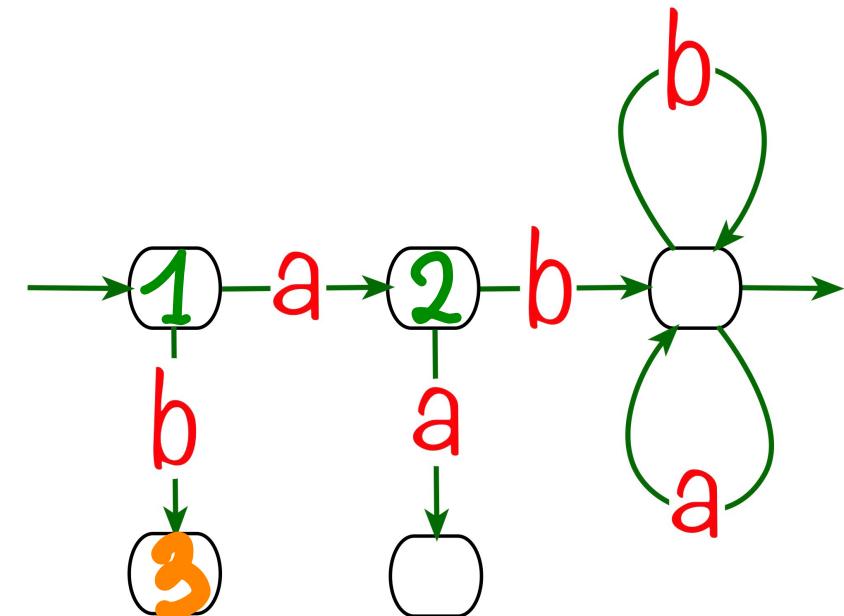
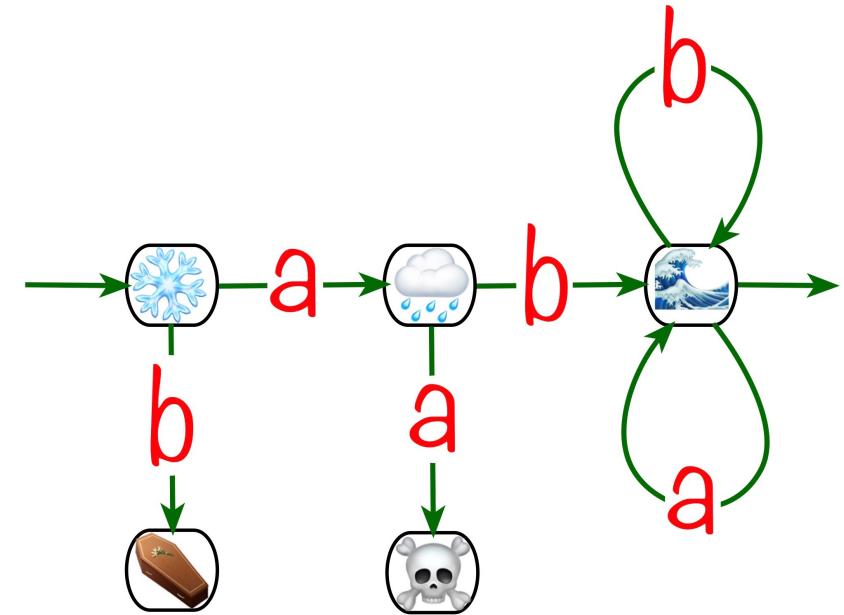


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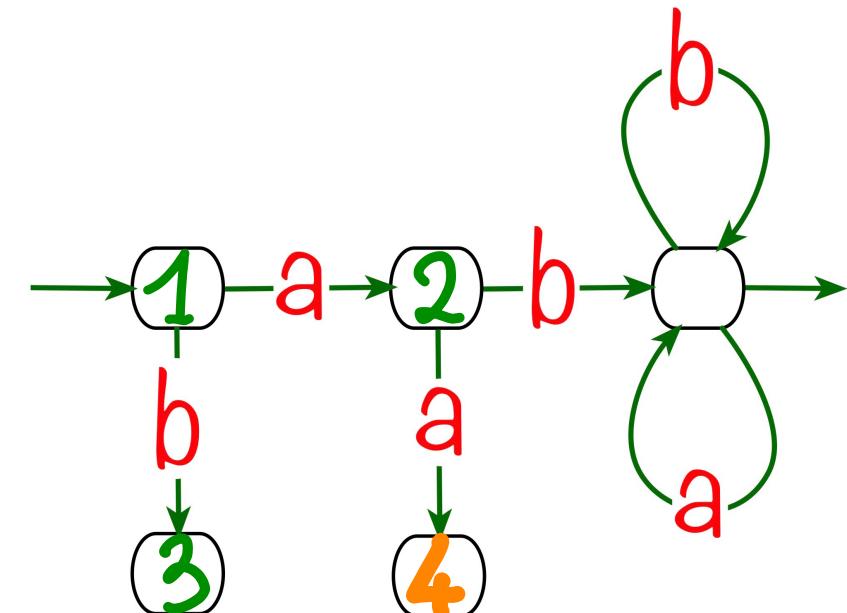
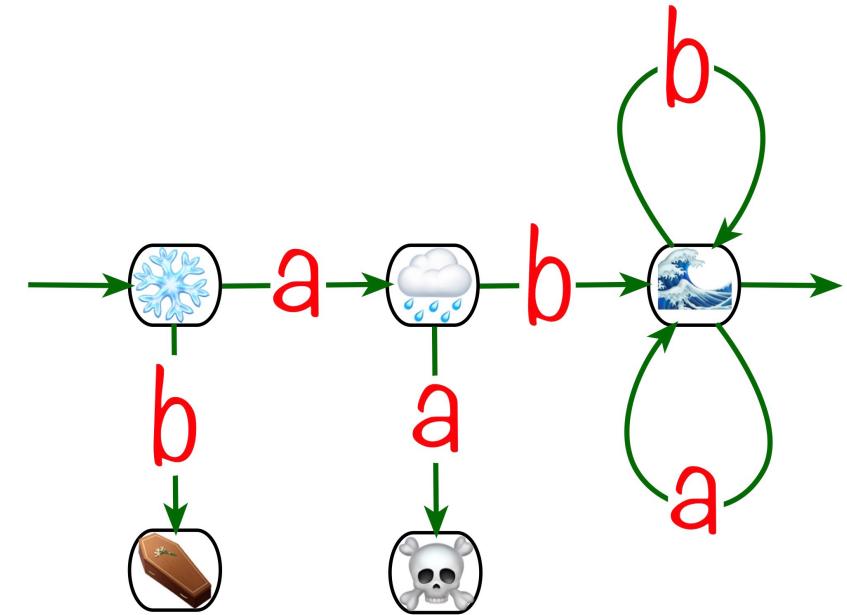
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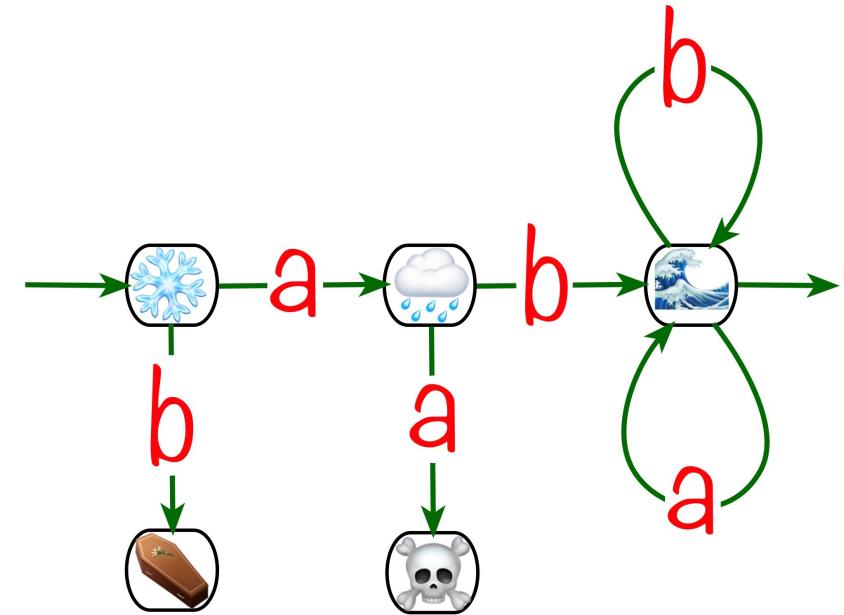
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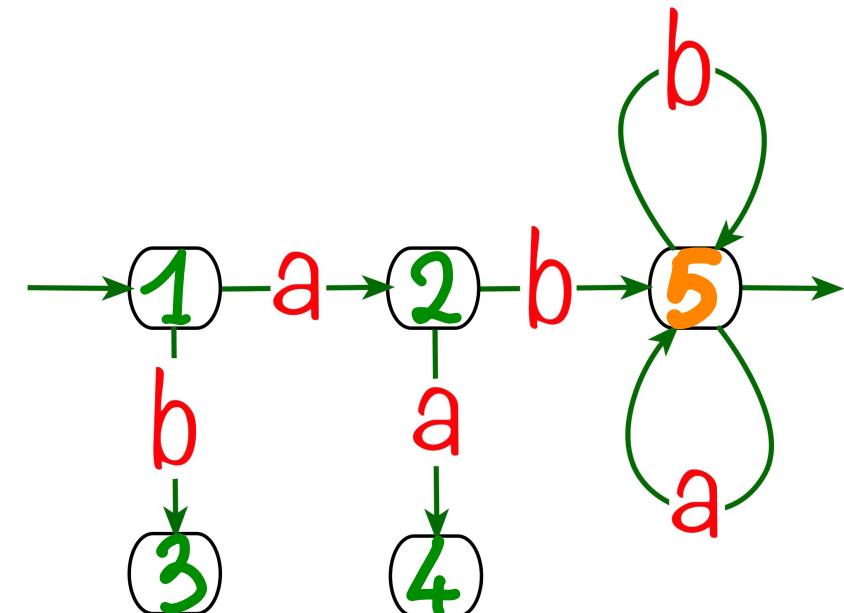
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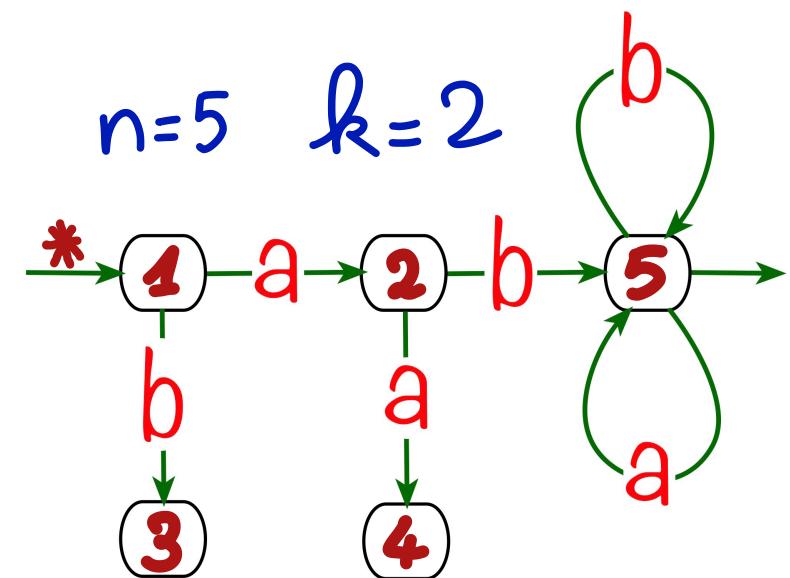
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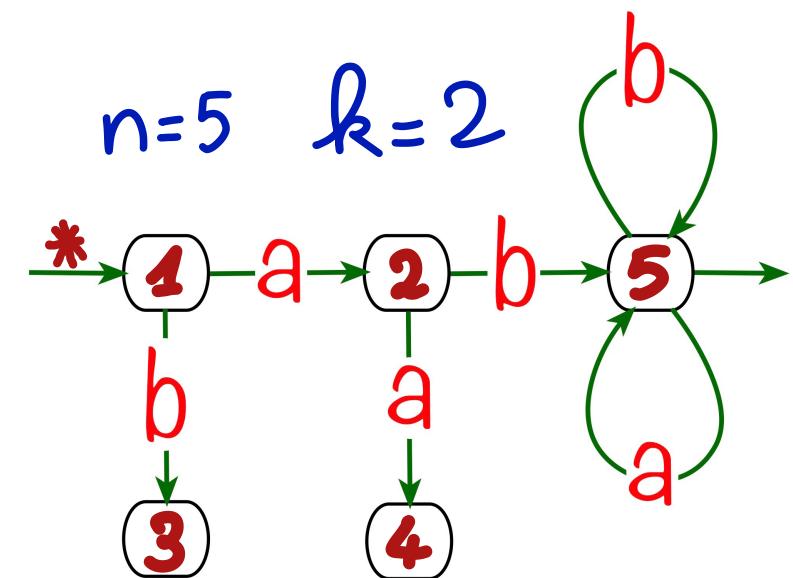


# Relation automata - coupon collector



*	1a	1b	2a	2b	3a	3b	4a	4b	5a	5b
1	2	3	4	5	3	3	4	4	5	5

# Relation automata - coupon collector



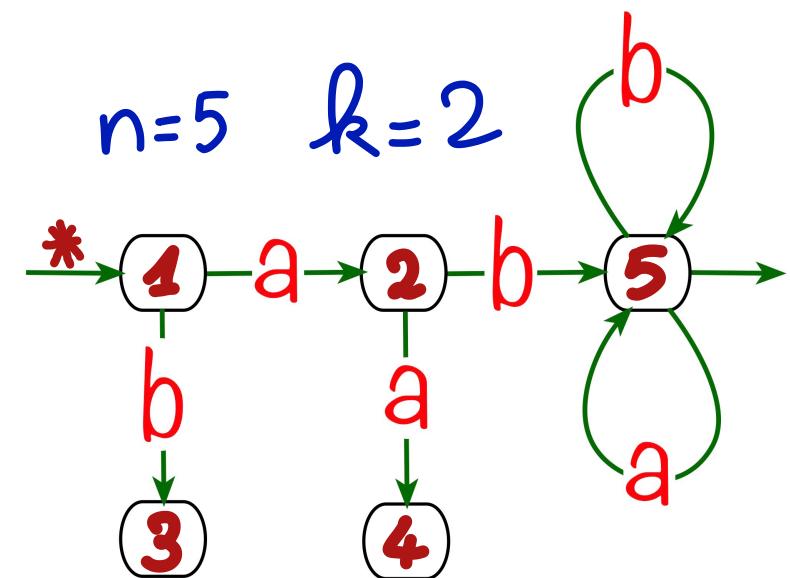
*	1a	1b	2a	2b	3a	3b	4a	4b	5a	5b
1	2	3	4	5	3	3	4	4	5	5

automata in  $\mathcal{Q}_{k_n}$



map from  $[1, 1+k_n]$  to  $[1, n]$ .

# Relation automata - coupon collector



*	1a	1b	2a	2b	3a	3b	4a	4b	5a	5b
1	2	3	4	5	3	3	4	4	5	5

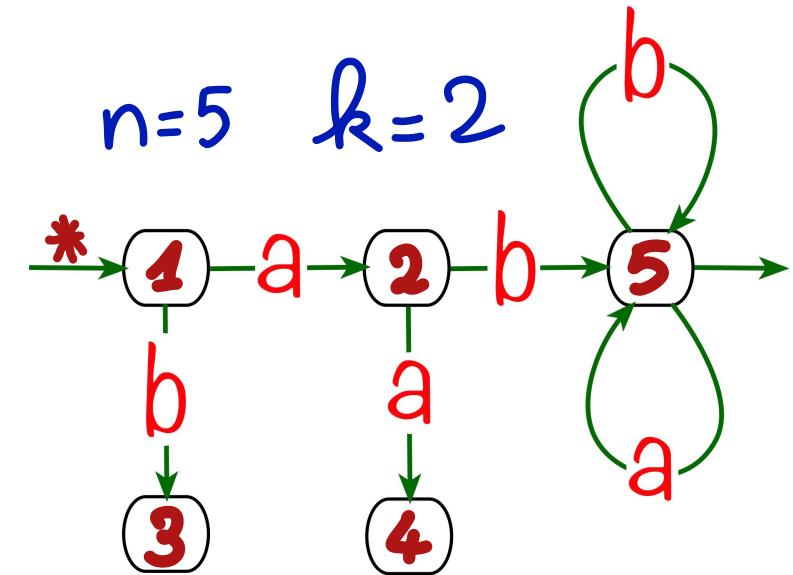
automata in  $\mathcal{Q}_{kn}$

accessible  
↓  
surjection! ~~map from  $[1, 1+kn]$  to  $[1, n]$ .~~

# Relation automata - coupon collector

automata  
 ??  
 impatient collector  
 $\lambda = k$   
 $\# \Omega_{k,n} = \binom{1+kn}{n} \times n!$   
 $\uparrow$

accessible  
 $\Downarrow$   
 surjection! map from  $[1, 1+kn]$  to  $[1, n]$ .



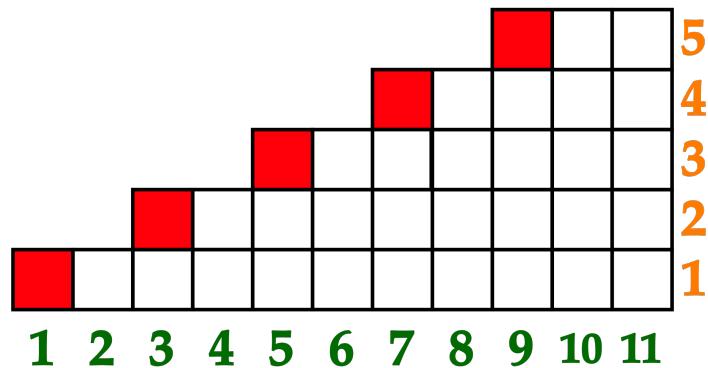
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1	2	3	4	5	3	3	4	4	5	5

automata in  $\Omega_{k,n}$

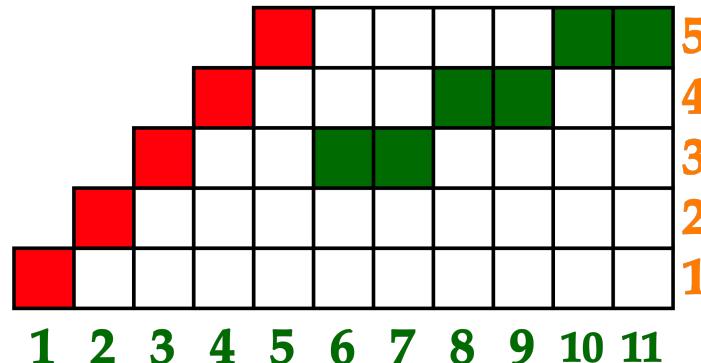


map from  $[1, 1+kn]$  to  $[1, n]$ .

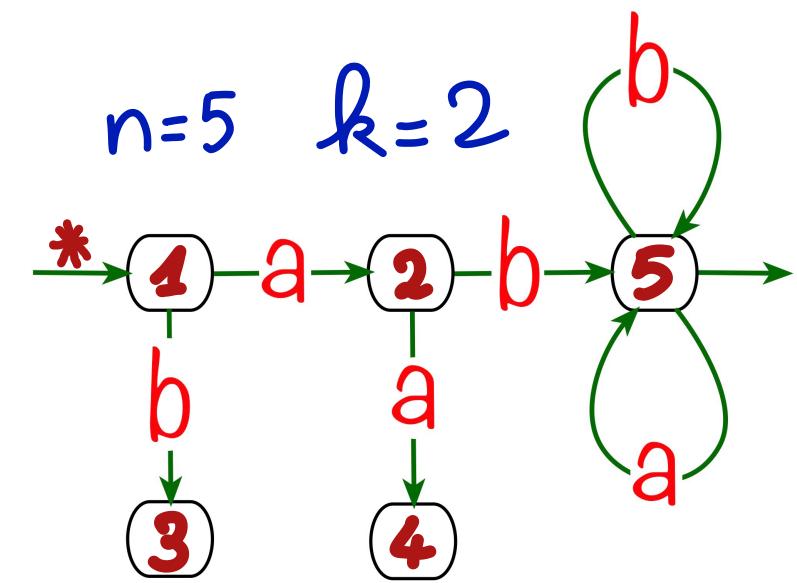
# Relation automata - coupon collector



automata  
??  
impatient collector  
 $\lambda = k$   
 $\#\Omega_{k,n} \stackrel{?}{=} \binom{1+kn}{n} \times n!$   
 $\uparrow$



accessible  
 $\Downarrow$   
 surjection!



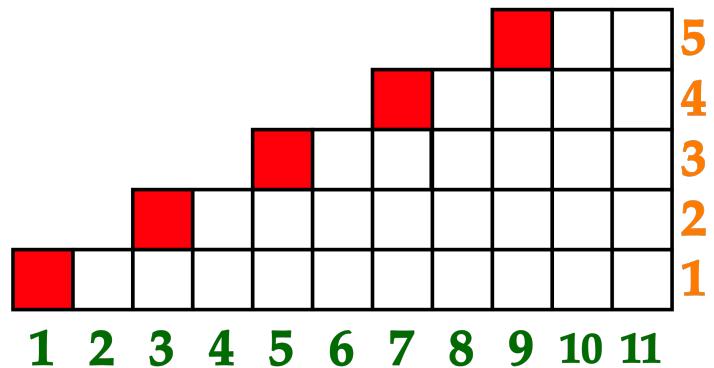
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1	2	3	4	5	3	3	4	4	5	5

automata in  $\Omega_{k,n}$

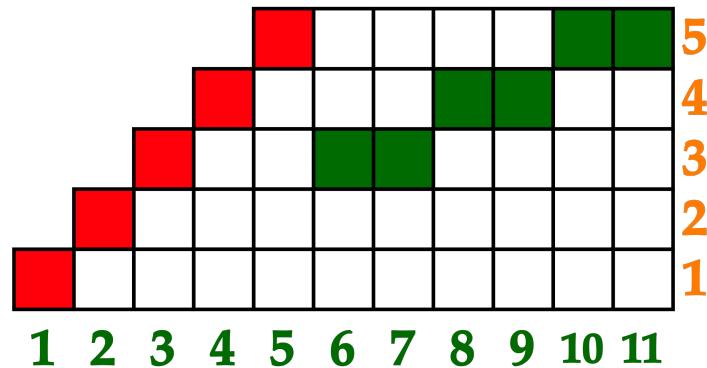
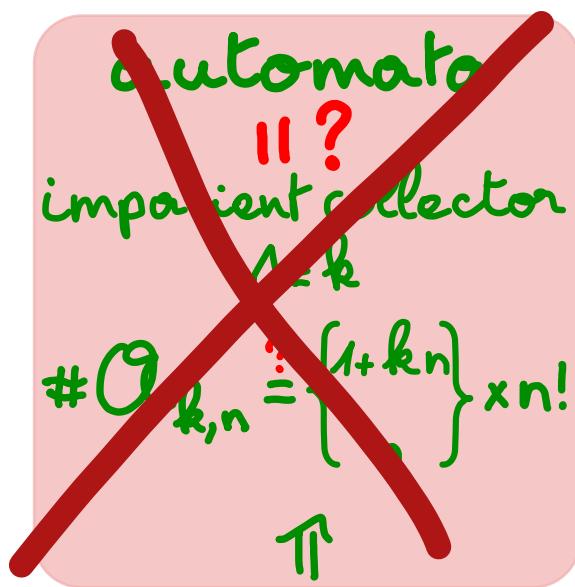


map from  $[1, 1+kn]$  to  $[1, n]$ .

# Relation automata - coupon collector

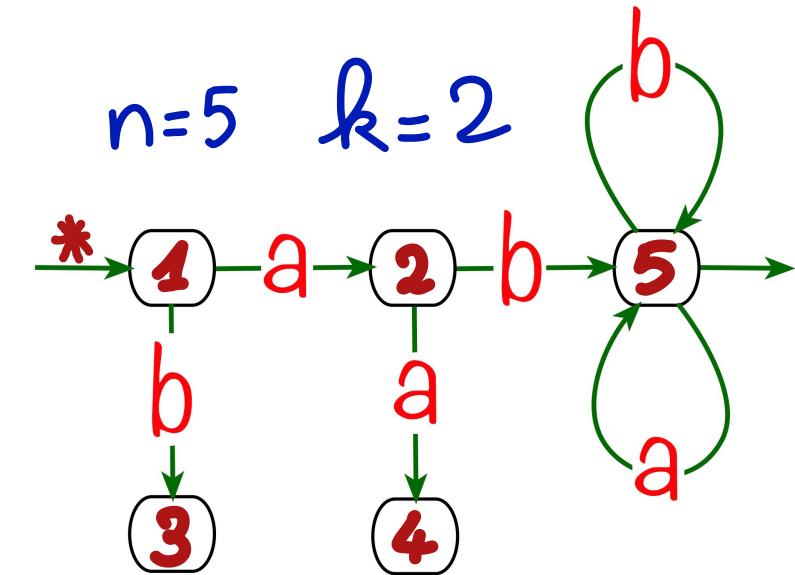


$$\#\Omega_{k,n} = \binom{1+kn}{n} \times n! \times C(k,n) \in [0, 1[$$



accessible  
 $\Downarrow$   
surjection !

~~map from  $[1, 1+kn]$  to  $[1, n]$ .~~



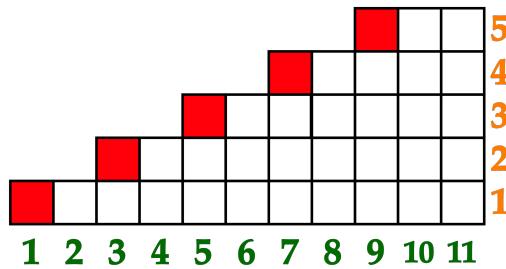
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automata in  $\Omega_{k,n}$



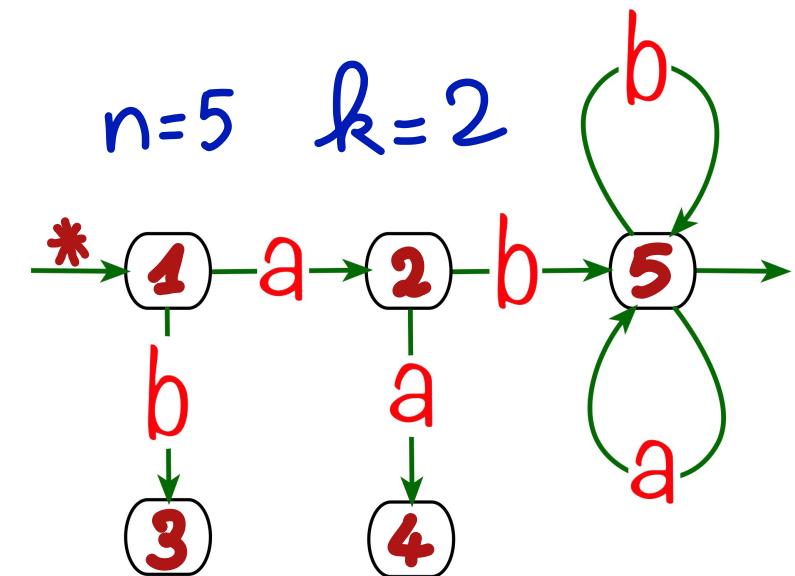
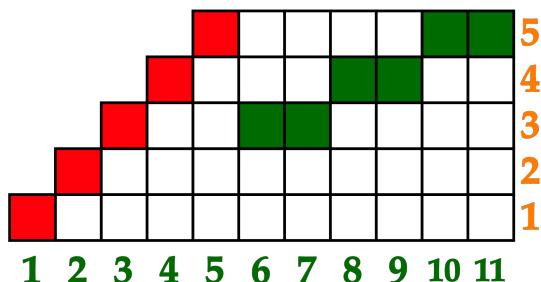
~~map from  $[1, 1+kn]$  to  $[1, n]$ .~~

# Relation automata - coupon collector



$$\#\Omega_{k,n} = \binom{1+kn}{n} \times n! \times C(k,n) \in [0, 1[$$

some probability in the  
impatient collector model



*	1a	1B	2a	2B	3a	3B	4a	4B	5a	5B
1	2	3	4	5	3	3	4	4	5	5

automata ??  
impatient collector  
 $\# \Omega_{k,n} = \binom{1+kn}{n} \times n!$   
↑

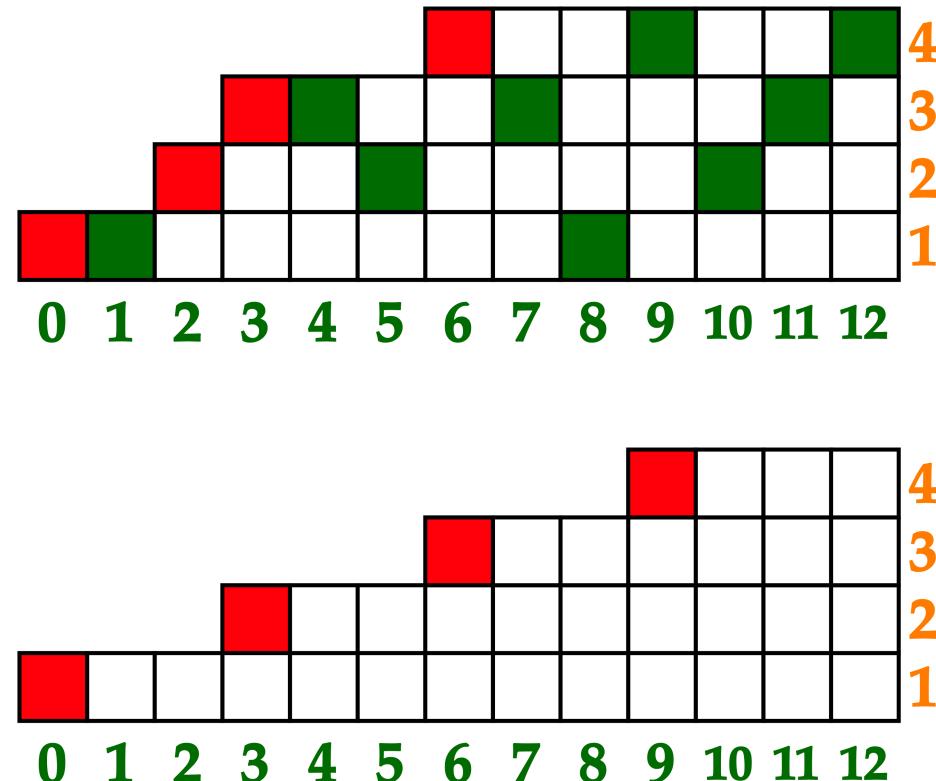
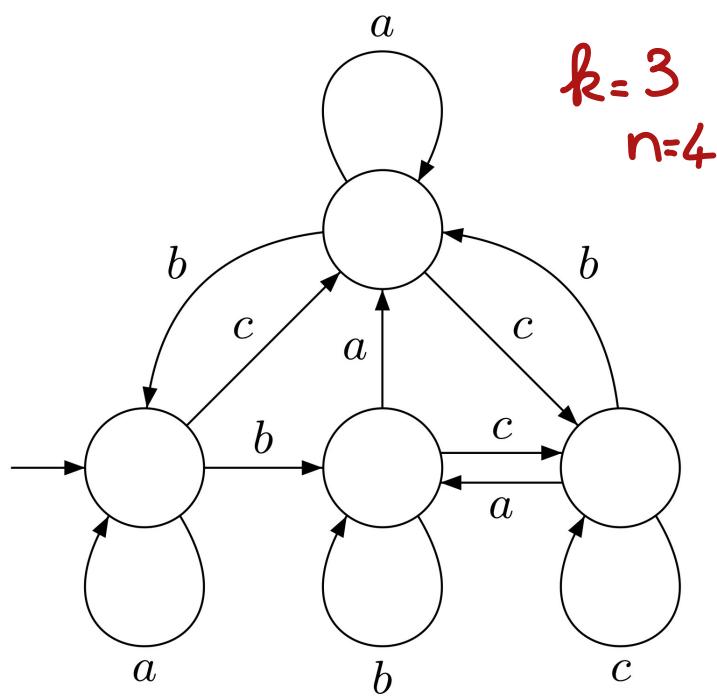
accessible  
↓  
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automata in  $\Omega_{k,n}$

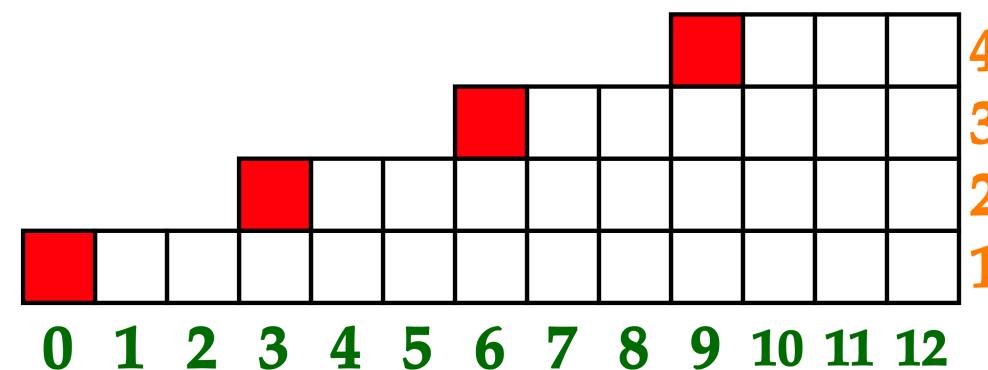
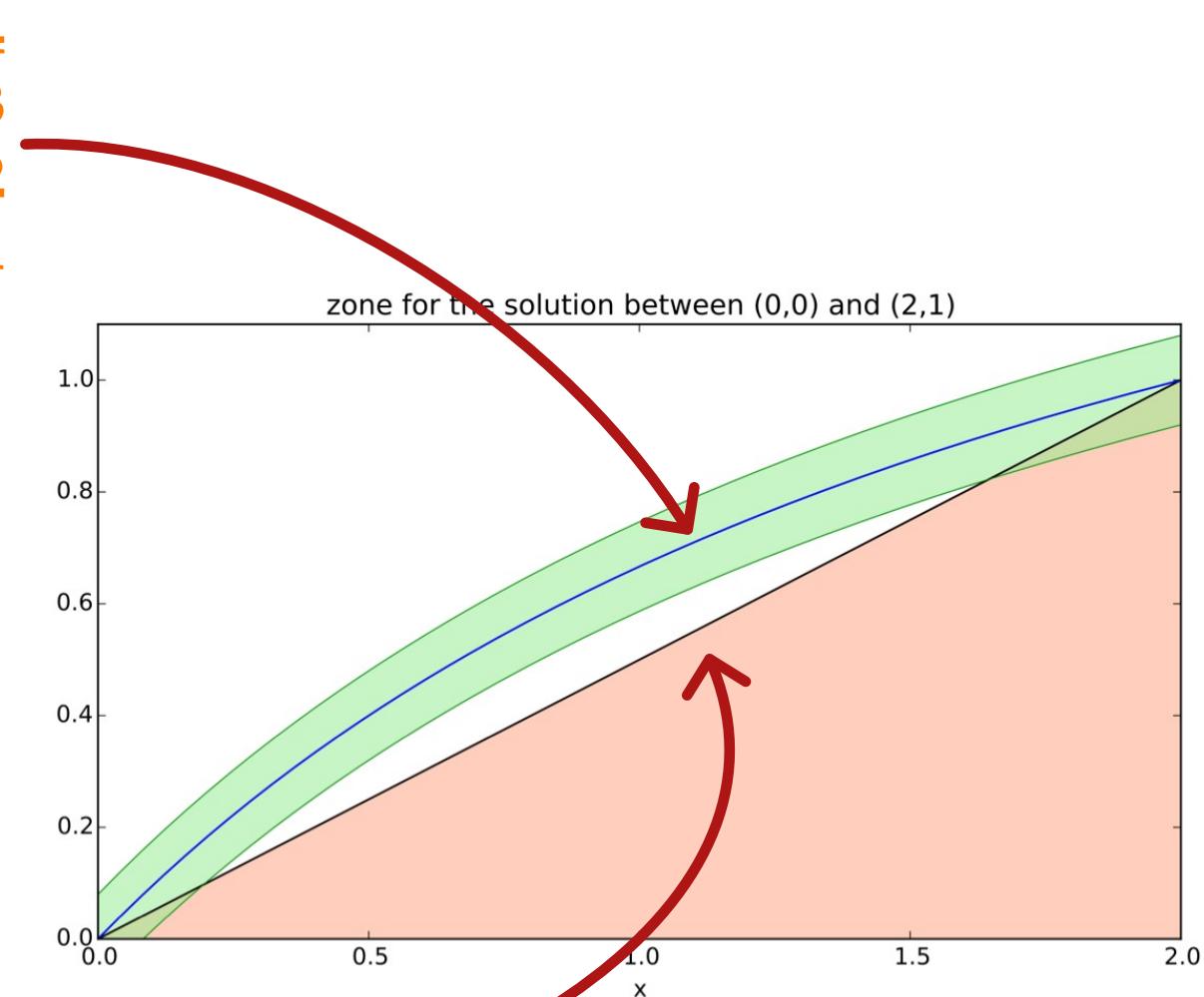
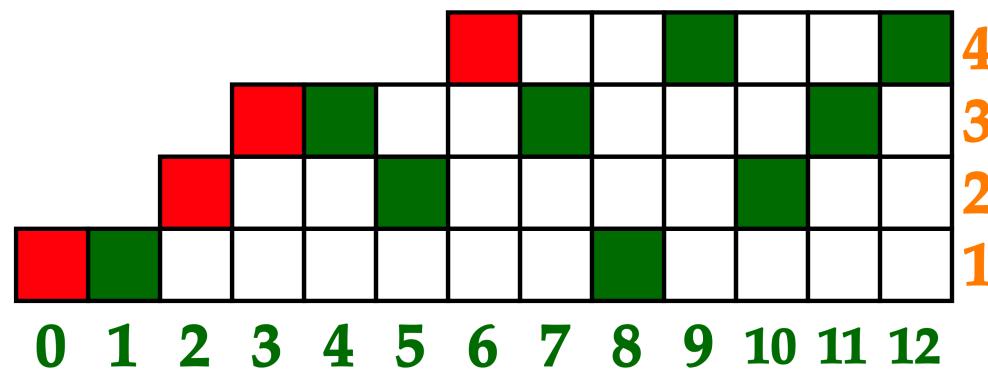


# Koršunov 1978

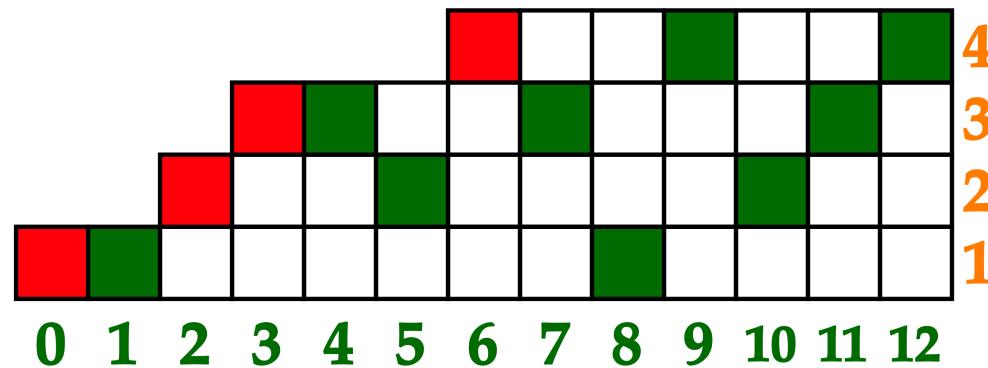
$$\lim_n \frac{\#\alpha_{k,n}}{\binom{1+kn}{n} \times n!} = C(k) \in ]0, 1[$$



# Limit profile & forbidden zone

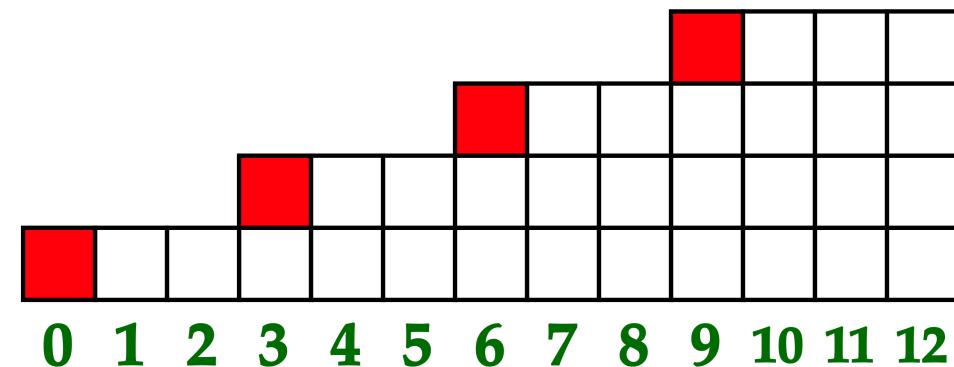


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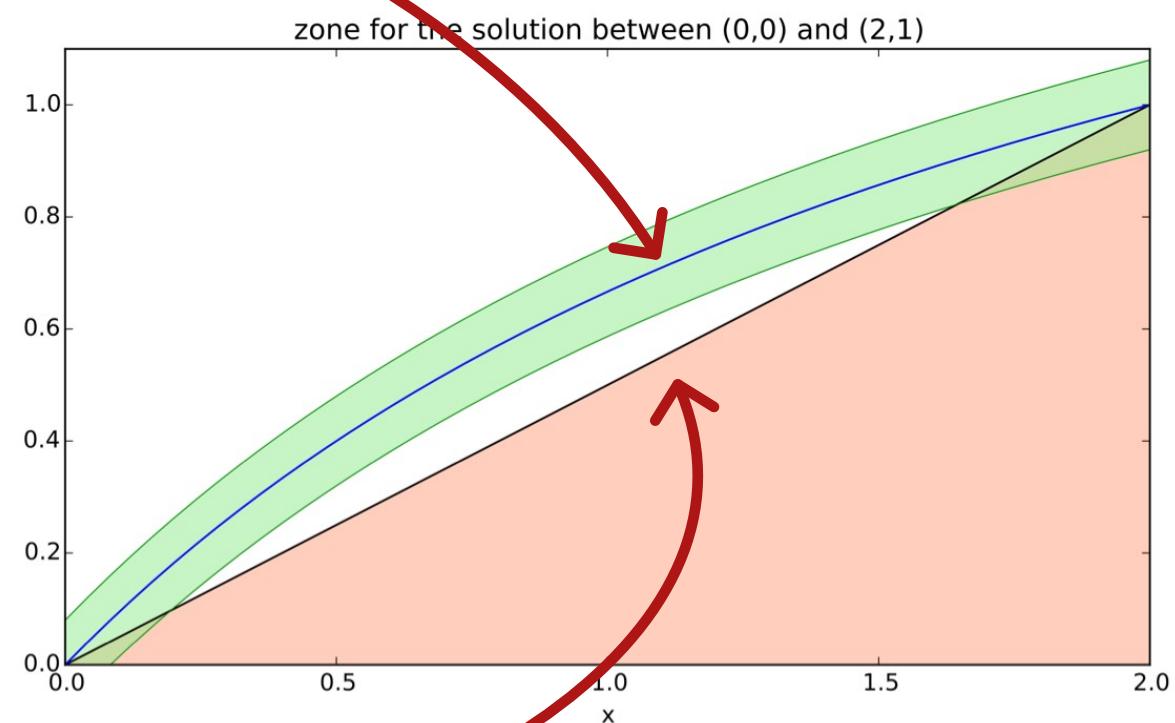


4  
3  
2  
1

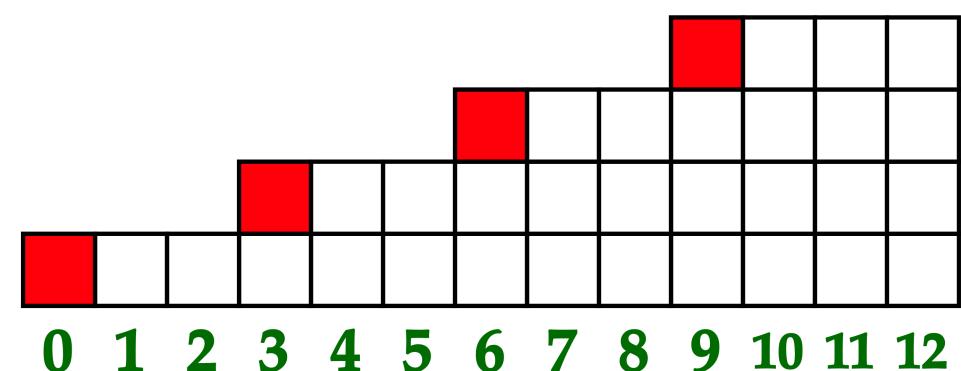
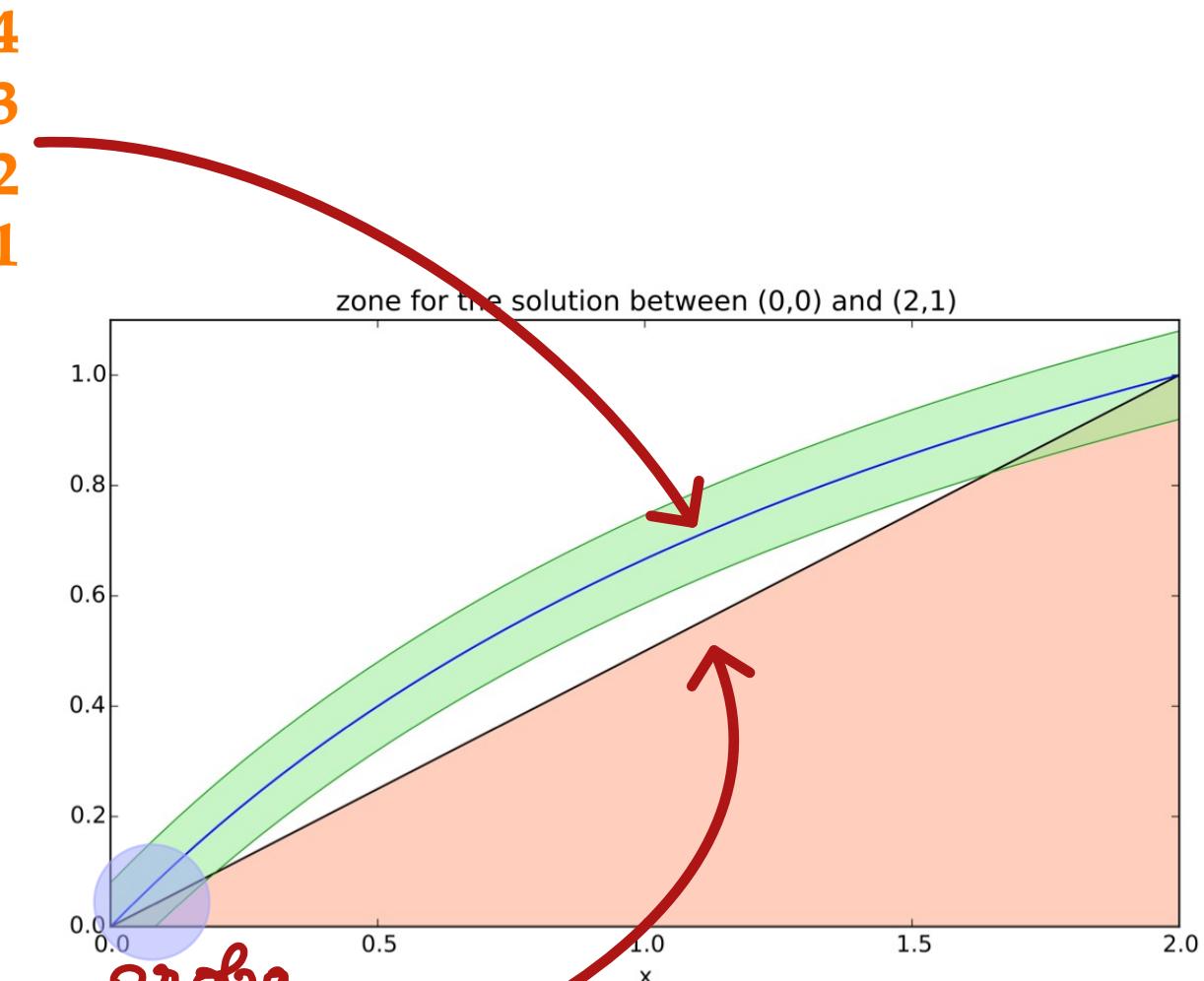
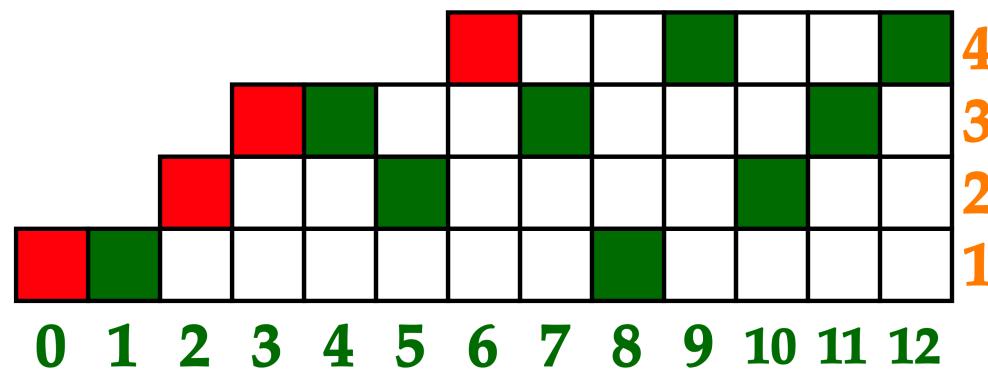
$$\mathbb{P}_{N,n} \left( \sup_{[a,1+\Lambda]} |\zeta_n - \zeta(\Lambda, \cdot)| \geq Cn^{-1/3} \right) \leq n^{1/3} e^{-\ln^2 n/2}$$



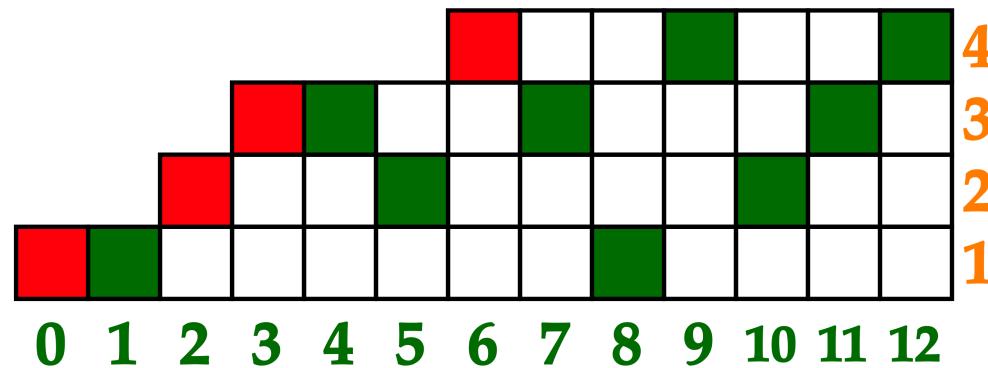
4  
3  
2  
1



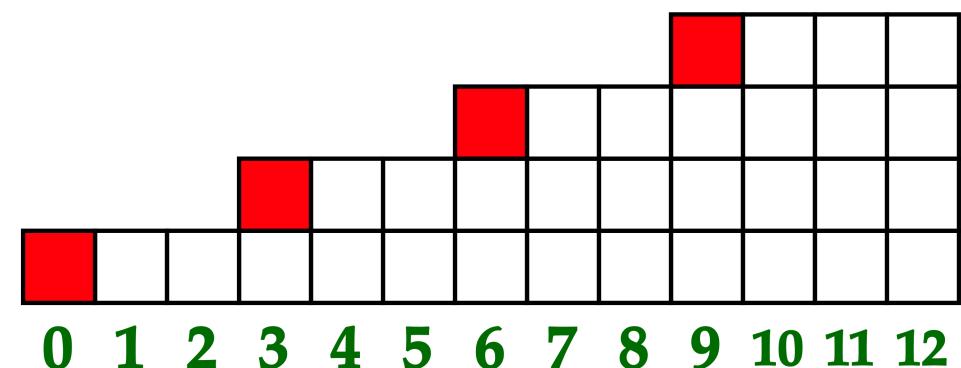
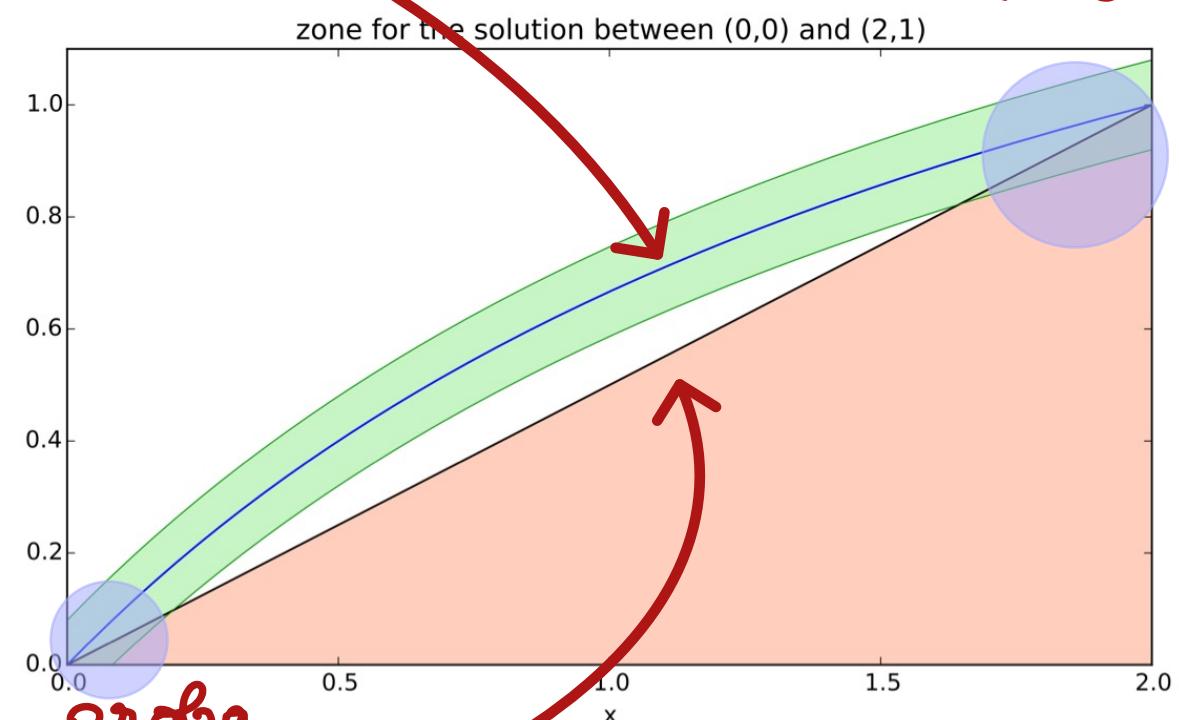
# Limit profile & forbidden zone



# Limit profile & forbidden zone



4  
3  
2  
1



4  
3  
2  
1

# Asymptotics for $\binom{m}{\ell}$

$$\binom{m}{\ell} \sim \psi(m, \ell) = \frac{m!(e^\xi - 1)^\ell}{\ell! \xi^m \sqrt{2\pi m \left(1 - \frac{m}{\ell} e^{-\xi}\right)}}.$$

Good 1961

## AN ASYMPTOTIC FORMULA FOR THE DIFFERENCES OF THE POWERS AT ZERO

By I. J. Good

*Admiralty Research Laboratory, Teddington, England*

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I. J. GOOD

solved by the above iterative method, thus  $\eta_{m+1} = (1 + \kappa) \log(1 + \eta_m)$ . Table II lists the values of  $\eta + 1 = (1 - \rho)^{-1}$ , where  $\rho$  is the solution of equation (15).

I am indebted to the Admiralty for permission to publish this paper, and for the use of the computer.

### REFERENCES

- [1] R. A. BUCKINGHAM, *Numerical Methods*, Pitman, London, 1957.
- [2] H. E. DANIELS, "Saddlepoint approximations in statistics," *Ann. Math. Stat.*, Vol. 25 (1954), pp. 631–650.

# Thanks !

