



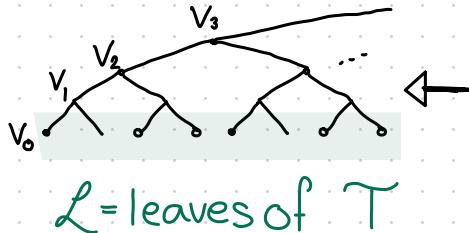
# Convergence of recursive equations via numerical analysis

Louigi Addario-Berry, Luc Devroye, Celine Kerriou, Rivka Maclaine Mitchell

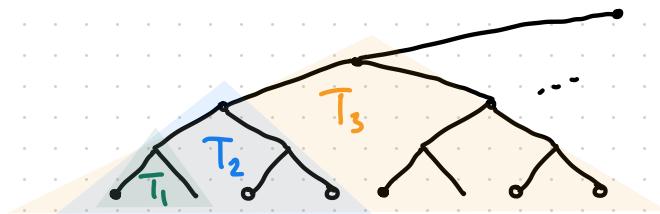
Analysis of Algorithms, CIRM, Luminy, France

June 24, 2019

$T$  = infinite binary canopy tree



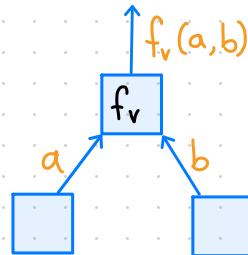
- one-way infinite path  $v_0, v_1, v_2, \dots$
- node  $v_n$  is the root of a complete binary tree of depth  $n$



$T_n$  = subtree rooted at  $v_n$   
 $L_n$  = leaves of  $T_n$

Functions on  $T$ :

- input from children
- combination function at nodes
- output to parents



Choose functions  $(f_v, v \in T_n \setminus L_n)$ ; this turns  $T_n$  into a function,

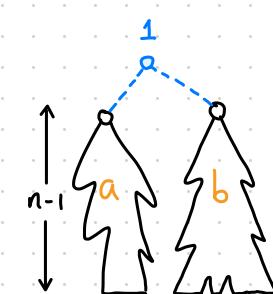
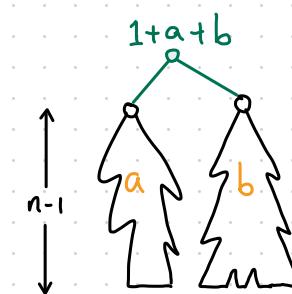
$$x = (x_v, v \in L_n) \xrightarrow{\quad} T_n(x) \leftarrow \text{output at root } v_n \text{, on input } x.$$

Either or both of  $x$  and  $(f_v, v \in T_n \setminus L_n)$  can be random

## Examples

①

$$f_v = \begin{cases} (a, b) \mapsto 1+a+b & \text{with prob. } p \\ (a, b) \mapsto 1 & \text{with prob. } 1-p \end{cases}$$

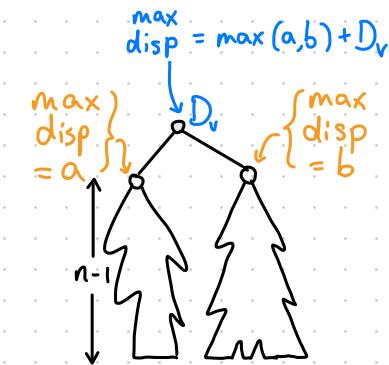


Then  $T_n(\vec{1}) \stackrel{d}{=} \# \text{ nodes at level } \leq n \text{ in a Galton-Watson tree}$

with offspring dist.  $\begin{cases} 2 & \text{with probability } p \\ 0 & \text{with probability } 1-p \end{cases}$

② Let  $(D_v, v \in T)$  be IID with law  $\mu$ , let  $f_v(a, b) = \max(a, b) + D_v$

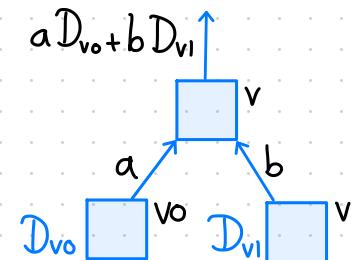
Then  $T_n(\vec{0}) \stackrel{d}{=} \text{maximum position in generation } n \text{ of a binary branching random walk with displacement dist } \mu$   
 (displacements at vertices)



③ Let  $(D_v, v \in T)$  be IID with law  $\mu$ , let  $f_v(a, b) = aD_{v_0} + bD_{v_1}$ .

This is a smoothing transform; fixed points studied by Durrett & Liggett (1983), many others.

In fact, all these equations have been studied from the perspective of fixed-point equations (sometimes wish to introduce a rescaling or shift).



# Examples without a fixed-point theory

(4) Derrida-Retaux model / "Parking on trees". Here  $f_v(a,b) = \max(a+b-1, 0)$

Question: Large- $n$  behaviour of  $T_n(X)$  where  $X = (X_v, v \in \mathcal{L}_n)$  IID with some law  $\mu$

(answer of course depends on  $\mu$ )

[Refs: Hu, Mallein, Pain, 1811.08749v2 ; Hu, Shi, 1705.03792 ; Goldschmidt, Przykucki, 1610.08786]

There is exactly one model where this can be analyzed so far.

(5) Random hierarchical lattice.

$$f_v = \begin{cases} (a,b) \mapsto a+b & \text{with prob. } p \\ (a,b) \mapsto \frac{ab}{a+b} & \text{with prob. } 1-p \end{cases}$$

[Ref: Hambly - Jordan 2004.  $p > \frac{1}{2} \rightarrow T_n(\bar{1})$  grows exponentially;  $p < \frac{1}{2} \rightarrow T_n(\bar{1})$  decays exp.]

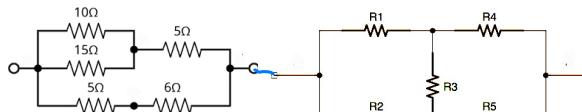
(6) Pemantle's Min-Plus tree

$$f_v = \begin{cases} (a,b) \mapsto a+b & \text{with prob. } p \\ (a,b) \mapsto \min(a,b) & \text{with prob. } 1-p \end{cases}$$

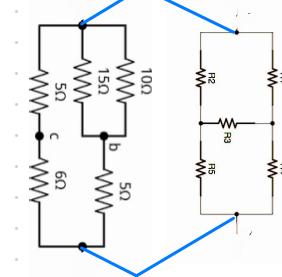
[Ref: Auffinger - Cable: 1709.07849]

(Open question: universality: what happens for other inputs?)

Series connection Resistance  $\rightarrow a+b$



Parallel connection  
Resistance  $\rightarrow \frac{ab}{a+b}$



Theorem (A-C)  $\frac{\log T_n(\bar{0})}{(\pi^2 n/3)^{1/2}} \xrightarrow{d} \text{Beta}(2,1)$

# New model Hipster random walk

Fix  $(D_v, v \in \mathcal{L})$  i.i.d. Let  $f_v$  be defined by

$$(a, b) \xrightarrow{f_v} a + D_v \mathbf{1}_{a=b} \quad \text{with prob. } \frac{1}{2}$$

$$(a, b) \xrightarrow{f_v} b + D_v \mathbf{1}_{a=b} \quad \text{with prob. } \frac{1}{2}$$

**Idea** Think of time as running up the tree

1 One of  $v_0, v_1$  is hipper than the other (chosen randomly)

2 If another particle shows up, hipper child takes off.

We will study • symmetric simple hipster random walk

SSHRW

• totally asymmetric lazy simple hipster random walk  $\rightarrow$  TALSHRW

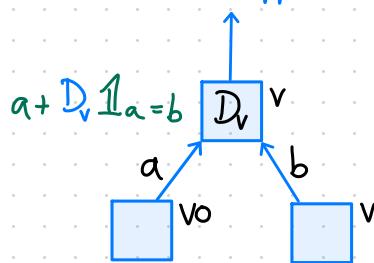
## Theorem

For SSHRW,  $\frac{T_n(\vec{0})}{(36n)^{\frac{1}{3}}} \xrightarrow{d} \text{Beta}(2,2) - \frac{1}{2}$ .

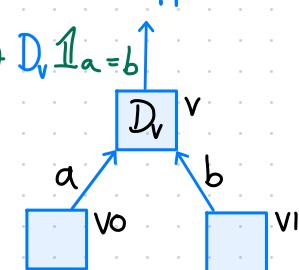
For TALSHRW

$\frac{T(\vec{0})}{(4(1-p)n)^{\frac{1}{2}}} \xrightarrow{d} \text{Beta}(2,1)$

$v_0$  is hipper



$v_1$  is hipper



**Note** Result for TALSHRW very similar to that of Auffinger-Cable.

Recall Auffinger-Cable:

$$f_v = \begin{cases} (a,b) \mapsto a+b & \text{with prob. } p \\ (a,b) \mapsto \min(a,b) & \text{with prob. } 1-p \end{cases}$$

Theorem (A-C)  $\frac{\log T_n(\bar{\alpha})}{(\pi^2 n / 3)^{1/2}} \xrightarrow{d} \text{Beta}(2,1)$

**Intuition** Suppose  $T_n(\bar{\alpha})$  is growing on a (stretched) exponential scale.

Write  $L, R$  for values at children of root of  $T_n$ .

If  $|\log L - \log R|$  small then  $\begin{cases} L+R \approx 2L \\ \min(L,R) \approx L \end{cases}$   $\begin{cases} \log(L+R) \approx \log(L) + 1 \\ \min(\log L, \log R) \approx \log L \end{cases}$  This is the common value plus a  $\{\bar{0}, \bar{1}\}$ -valued increment

If  $|\log L - \log R|$  big then  $\begin{cases} L+R \approx \max(L, R) \\ \min(L, R) = \min(L, R) \end{cases}$   $\begin{cases} \log(L+R) \approx \max(\log L, \log R) \\ \log(\min(L, R)) = \min(\log L, \log R) \end{cases}$  This is just  $\log(\text{value of a random child})$

Similar intuition should work for the hierarchical lattice:

$$f_v = \begin{cases} (a,b) \mapsto a+b & \text{with prob. } p \\ (a,b) \mapsto \frac{ab}{a+b} & \text{with prob. } 1-p \end{cases}$$

Intuition: Suppose  $\tilde{T}_n(\bar{o})$  is growing on a (stretched) exponential scale.

Write  $L, R$  for values at children of root.

If  $|\log L - \log R|$  small then  $\left\{ \begin{array}{l} L+R \approx 2L \quad \log(L+R) \approx \log(L) + 1 \\ \frac{LR}{L+R} \approx L/2 \quad \log\left(\frac{LR}{L+R}\right) \approx \log(L) - 1 \end{array} \right\}$  This is the common value plus a  $\{-1, 1\}$ -valued increment

If  $|\log L - \log R|$  big then  $\left\{ \begin{array}{l} L+R \approx \max(L, R) \quad \log(L+R) \approx \max(\log L, \log R) \\ \frac{LR}{L+R} \approx \min(L, R) \quad \log\left(\frac{LR}{L+R}\right) = \min(\log L, \log R) \end{array} \right\}$  This is just  $\log(\text{value of a random child})$

Motivates the following conjecture: in the random hierarchical lattice with  $p=\frac{1}{2}$ ,  $\exists c > 0$  s.t.

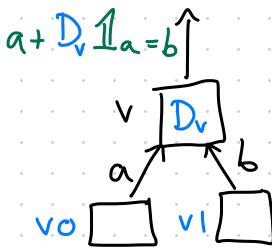
$$\frac{\log \tilde{T}_n(\bar{o})}{(c n)^3} - \frac{1}{2} \xrightarrow{d} \text{Beta}(2,2)$$

Theorem (Totally asymmetric lazy)  $\frac{T(\vec{0})}{(2n)^{\frac{1}{2}} \alpha} \xrightarrow{d} \text{Beta}(2,1)$

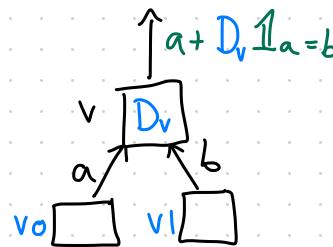
Proof Idea

Original dynamics:

$v_0$  is hipper



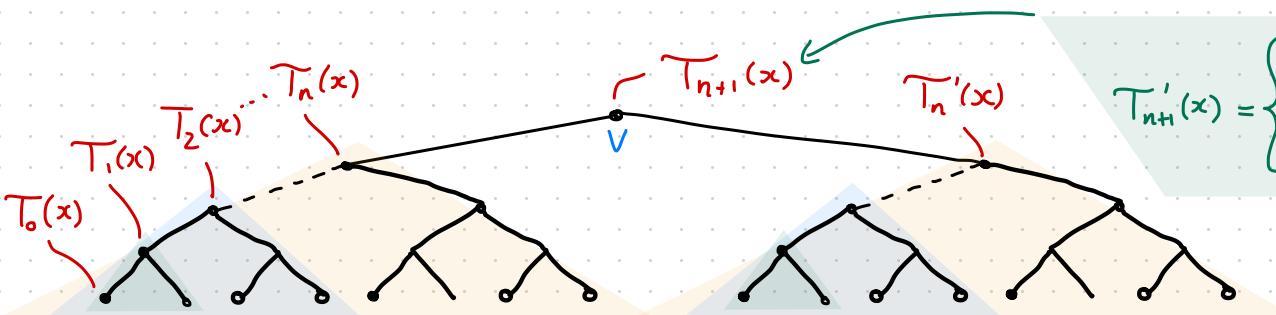
$v_1$  is hipper



$$D_v \sim \text{Bernoulli}\left(\frac{1}{2}\right)$$

By symmetry, can assume left child is always chosen.

For inputs  $x = (x_v, v \in \mathcal{L})$ , useful notation:  $T_n(x) := T_n((x_v, v \in \mathcal{L}_n))$



$$T'_{n+1}(x) = \begin{cases} T_n(x) & \text{if } T_n(x) \neq T'_n(x) \\ T_n(x) + D_v & \text{if } T_n(x) = T'_n(x) \end{cases}$$

Proof Idea

(Totally asymmetric case)

$$P_n(k)(1 - P_n(k)) \leftarrow \text{left child} = k, \text{right child} \neq k$$

$$\frac{1}{2} P_n(k-1)^2 \leftarrow \text{both} = k-1, \text{make a step}$$

$$\frac{1}{2} P_n(k)^2 \leftarrow \text{both} = k, \text{be lazy}$$

Let  $P_n(k) = P(T_n(\vec{\omega}) = k)$

Then  $P_{n+1}(k) = P_n(k)(1 - P_n(k)) + \frac{1}{2} P_n(k-1)^2 + \frac{1}{2} P_n(k)^2$

Rearranging gives  $P_{n+1}(k) - P_n(k) = \frac{1}{2} (P_n(k)^2 - P_n(k-1)^2)$

This is a discretization of the inviscid Burgers' equation  $\frac{\partial}{\partial t} u(x,t) = -\frac{1}{2} \frac{\partial}{\partial x} (u(x,t)^2)$

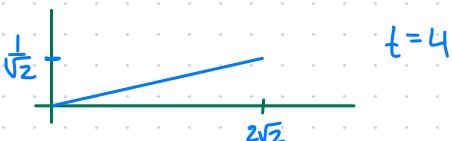
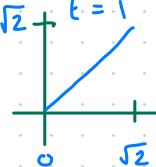
So we are trying to solve the (measure-valued) initial-value problem

$$u_t = -\frac{1}{2} (u^2)_x = -u u_x$$

$$\begin{cases} u_t = -u u_x, & t \geq 0, x \in \mathbb{R} \\ u_0(x) = \delta_0(x) = \mathbf{1}_{[x=0]} & (\text{Dirac mass at } 0) \end{cases}$$

Ignoring space-time points of discontinuity, this is solved\* by  $u: \mathbb{R} \times [0, \infty) \rightarrow \mathbb{R}$  given by

$$u(x,t) = \begin{cases} \frac{x}{t}, & 0 \leq x < \sqrt{2t} \\ 0, & \text{otherwise} \end{cases}$$



Note  $u(t,x)$  is always a prob. dist.: the density of a scaled Beta(2,1).

\*But solution is not unique!

First step Start Burgers' from a smooth initial condition of the form  $u_0(x) = \frac{x}{t_0} \mathbf{1}_{0 \leq x \leq \sqrt{2t_0}}$   
 Probabilistically what does this mean? (think of  $t_0$  as small).

$u_0$  is density of  $\sqrt{2t_0} \cdot B$  where  $B \sim \text{Beta}(2, 1)$

Fix  $M > 0$  and define  $u_j^0(M) = M \int_{j/M}^{(j+1)/M} u_0(x) dx$ . for  $j \geq 0$  s.t.  $\frac{j}{M} \leq \sqrt{2t_0}$ .

Then  $\sum_j u_j^0(M) = 1$ , so  $(u_j^0(M), j \geq 0)$  defines a probability distribution on  $\{0, 1, \dots, \lfloor M \cdot \sqrt{2t_0} \rfloor\}$

Let  $X^M = (X_v^M, v \in \mathbb{Z})$  be vector of IID's with  $P(X_v^M = j) = u_j^0(M)$  (discretization of  $u_0$  at mesh size  $\frac{1}{M}$ )

$T_n(X^M)$  is value of TALSHRW when initial distribution is  $\frac{1}{M}$ -mesh discretization of  $\sqrt{2t_0} \cdot B$ .

**Lemma** We have  $P(T_n(X^M) = j) = \frac{1}{M} \cdot u_j^n(M)$ , where  $(u_j^n(M))_{n \geq 0, j \geq 0}$  is defined by the recurrence  $M \cdot u_j^{n+1} = M \cdot u_j^n - \frac{1}{2} ((u_j^n)^2 - (u_{j-1}^n)^2)$ .

**Proof** Easy induction  $\square$

Second step Convergence of the fine-mesh approximation.

The spatial mesh is  $\frac{1}{M}$ . We take a temporal mesh of  $\frac{1}{M^2}$ .

$$U_M(t, x) = U_{\lfloor t M^2 \rfloor}^{\lfloor t M^2 \rfloor}(M) = P(T_{\lfloor t M^2 \rfloor}(X^M) = \lfloor x M \rfloor) \quad \text{for } t, x \geq 0.$$

Call  $U_M$  a  $\frac{1}{M}$ -fine mesh approximation of Burgers' equation

Theorem (Evje & Karlsen, 2000)

From a bounded variation initial condition, the  $\frac{1}{M}$ -fine mesh approximation converges to the BV entropy solution  $u$  of Burgers' equation almost everywhere on  $\mathbb{R} \times [0, \infty)$ , and for any compact  $C \subset \mathbb{R} \times [0, \infty)$ ,  $\int_C |U^m(x, t) - u(x, t)| dx dt \rightarrow 0$ .

- BV  $\rightarrow$  Bounded variation
- BV entropy solution  $\rightarrow$  The correct solution of our problem  
(verifying this takes some work)

Conclusion  $U_M \rightarrow u$  defined by  $u(t, x) = \frac{x}{t+t_0} \mathbf{1}_{0 \leq x \leq \sqrt{2(t+t_0)}}$

# Implication for TALSHRW

**Corollary** For  $\varepsilon > 0$  small, if  $U = \text{Unif}[1-\varepsilon, 1+\varepsilon]$  is independent of  $X$ , then as  $M \rightarrow \infty$ ,

$$\frac{T_{\lfloor UM^2 \rfloor}(X^M)}{\sqrt{2(t_0+U)M}} \xrightarrow{d} \text{Beta}(2,1).$$

**Proof:** For any compact  $C \subset \mathbb{R} \times [0, \infty)$ ,

$$\iint_C \left| P(T_{\lfloor tM^2 \rfloor}(X^M) = \lfloor xM \rfloor) - \frac{x}{t+t_0} \mathbf{1}_{0 \leq x \leq \sqrt{2(t+t_0)}} \right| dx dt \rightarrow 0$$

Taking  $C = \{(x,t) : |t-t_0| \leq \varepsilon, 0 \leq x \leq \alpha \sqrt{2(t+t_0)}\}$ , this yields by the triangle inequality that

$$\int_{[-\varepsilon, \varepsilon]} \left| P(T_{\lfloor tM^2 \rfloor}(X^M) \leq \alpha \sqrt{2(t+t_0)} M) - \int_0^{\alpha} \frac{x}{t+t_0} dx \right| \cdot \frac{1}{2\varepsilon} dt \rightarrow 0 \quad \text{as } M \rightarrow \infty$$

(There are "discretization errors" coming from the floors, but it's easy to see these tend to 0 as  $M \rightarrow \infty$ .)

Since  $U$  has density  $\frac{1}{2\varepsilon} \mathbf{1}_{|t-t_0| \leq \varepsilon}$ , the result follows.  $\square$

Last step stochastic domination.

### Proposition

If  $x = (x_v, v \in \mathcal{L})$  and  $y = (y_v, v \in \mathcal{L})$  are such that  $x_v \in \mathbb{Z}$ ,  $y_v \in \mathbb{Z}$  and  $x_v \leq y_v$  for all  $v \in \mathcal{L}$ , then  $T_n(x) \leq_{st} T_n(y)$  for all  $n \geq 1$ .

Proof: A straightforward induction.  $\blacksquare$

**Corollary 1** For all  $n, M \in \mathbb{N}$ ,  $T_n(X^M) - \lfloor \sqrt{2t_0M} \rfloor \leq_{st} T_n(\vec{0}) \leq_{st} T_n(\vec{0})$ .  $\blacksquare$

Allows us to compare all-0 input to random input with  $o(M)$  error (recall  $t_0 > 0$  is fixed but arbitrarily small).

**Corollary 2** For all  $M \in \mathbb{N}$ ,  $T_{(1-\varepsilon)M^2}(X^M) \leq_{st} T_{UM^2}(X^M) \leq_{st} T_{(1+\varepsilon)M^2}(X^M)$

Allows us to compare fixed time near  $M^2$  to random time  $UM^2$ .

Since  $\frac{1}{\sqrt{2(t_0+U)}} M T_{UM^2}(X_M) \xrightarrow{d} \text{Beta}(2,1)$ , corollaries yield that  $\frac{T_n(X^M)}{\sqrt{2M}} \xrightarrow{d} \text{Beta}(2,1)$ .



# CIRM

## HIPSTERS