

# **KHOVANOV HOMOLOGY OF SEMIADEQUATE LINKS VIA PRESIMPLICIAL SETS**

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**AT**  
**UNIVERSITAT DE BARCELONA**

[JOINT WORK WITH JÓZEF H. PRZYTICKI]

**WINTERBRAIDS VIII**  
**FEBRUARY 6 - LUMINY (FRANCE)**

# Why to study Khovanov homology?

(Khovanov, 2000)

- It is a bigraded homology  $H^{i,j}(L)$ .

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$$H^{0,-3}(5_1) = \mathbb{Z}$$

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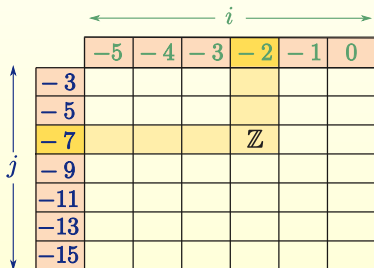
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$$H^{i,j}(5_1) = 0 \text{ otherwise}$$

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	-5	-4	-3	-2	-1	0
-3						
-5						
-7				$\mathbb{Z}$		
-9						
-11						
-13						
-15						



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	-5					$\mathbb{Z}$
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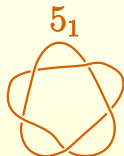
- It is a bigraded homology  $H^{i,j}(L)$ .
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- It **categorifies Jones polynomial**

$$V(L) = \sum_{i,j} q^j (-1)^i \operatorname{rk}(H^{i,j}(L)).$$

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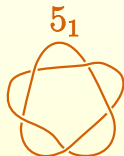
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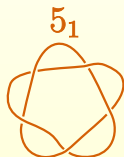
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	-5	-4	-3	-2	-1	0
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-5						Q
-7				Q		
-9						
-11		Q	Q			
-13						
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	-7	-6	-5	-4	-3	-2	-1	0
-1							Q	Q
-3								Q
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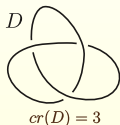
- $s$ -invariant  $\left\{ \begin{array}{l} \text{provides a lower bound on the slice genus of a knot.} \\ \text{gives a combinatorial proof of the (topological)} \\ \text{Milnor conjecture: } g_s(T(p, q)) = \frac{1}{2}(p-1)(q-1). \end{array} \right.$   
(Rasmussen, 2003)

# Why one should look for new approaches?

Although conceptually simple, the original definition of Khovanov homology becomes impractical.

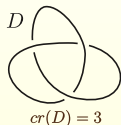
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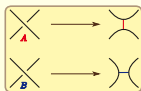


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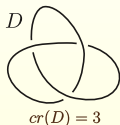


Kauffman states

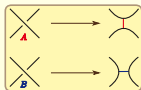


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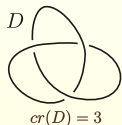


$$\# \{ \text{Kauffman states} \} = 2^{cr(D)} = 2^3.$$

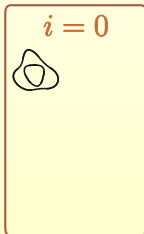
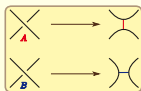


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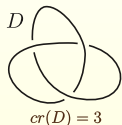


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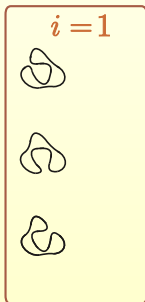
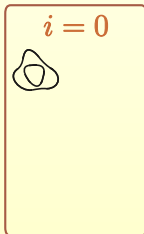
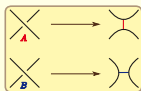


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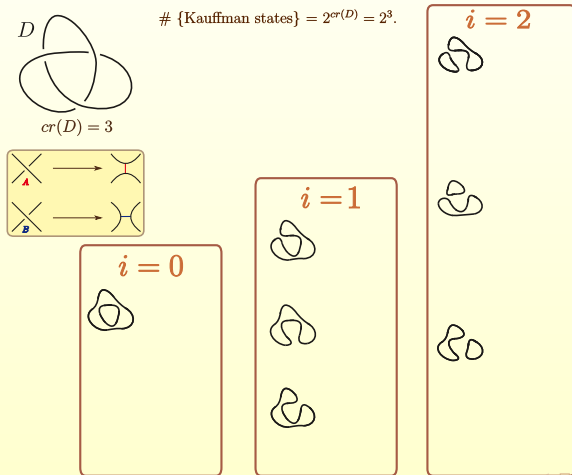
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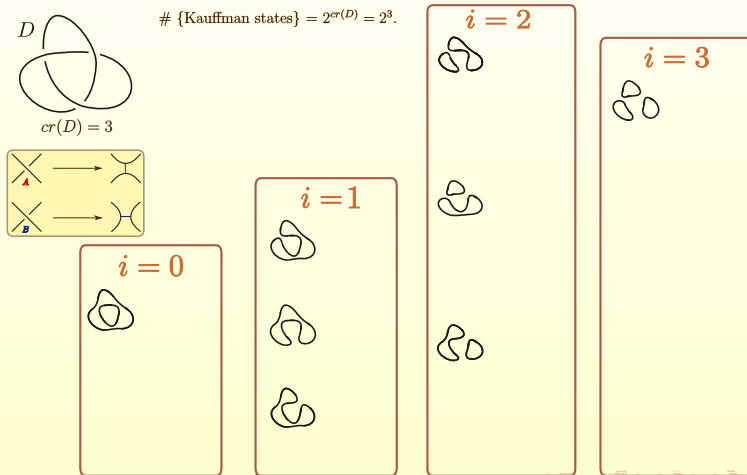
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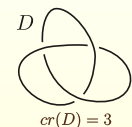
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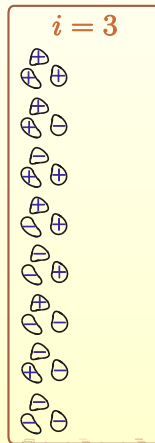
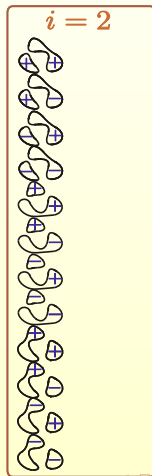
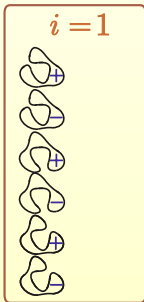
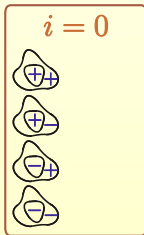
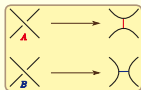
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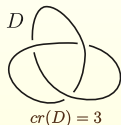
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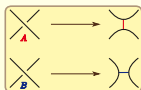
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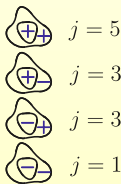


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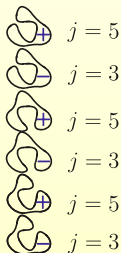
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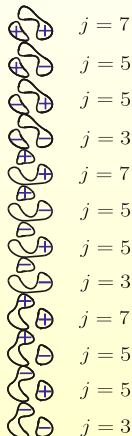
$i = 0$



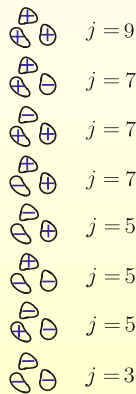
$i = 1$



$i = 2$

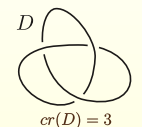


$i = 3$



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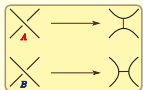
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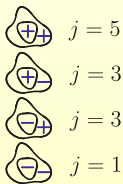
||

Generators

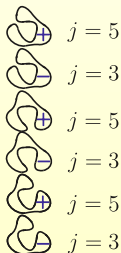
$$C^{i,j}(D) = \bigoplus_{\substack{s(a)=i \\ s(b)=j}} \mathbb{Z} s$$



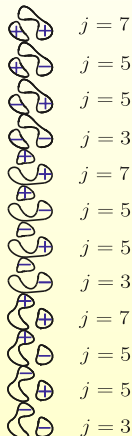
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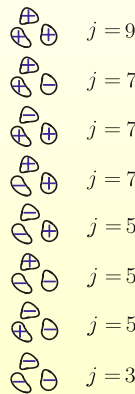
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We look for **new approaches**  
to Khovanov homology.

# Almost extreme Khovanov complex

$$\begin{array}{ccccccc}
 & \vdots & & \vdots & & \vdots & \\
 \cdots & \longrightarrow & C_{i+2,j+4}(D) & \longrightarrow & C_{i,j+4}(D) & \longrightarrow & C_{i-2,j+4}(D) \longrightarrow \cdots \\
 \cdots & \longrightarrow & C_{i+2,j}(D) & \longrightarrow & C_{i,j}(D) & \longrightarrow & C_{i-2,j}(D) \longrightarrow \cdots \\
 \cdots & \longrightarrow & C_{i+2,j-4}(D) & \longrightarrow & C_{i,j-4}(D) & \longrightarrow & C_{i-2,j-4}(D) \longrightarrow \cdots \\
 & \vdots & & \vdots & & \vdots & 
 \end{array}$$



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 & \vdots & & \vdots & & \vdots & 
 \end{array}$$

$$j_{\max}(D) = \max_{s \text{ state of } D} \{j(s)\}$$

# Almost extreme Khovanov complex

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Extreme Khovanov

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$$D \longrightarrow I_D \text{ simplicial complex} / H_i(I_D) = H^{i,j_{\max}}(D)$$

# Extreme Khovanov homology

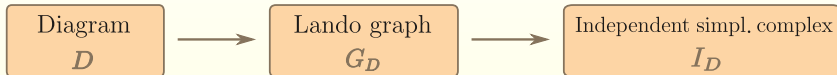
Diagram

$D$

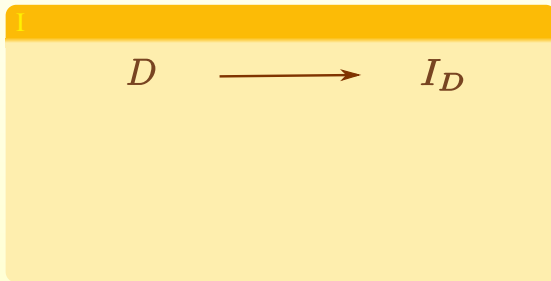
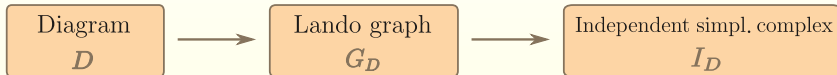
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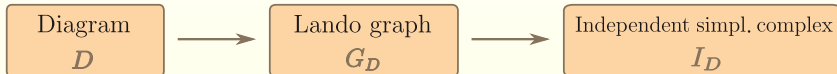
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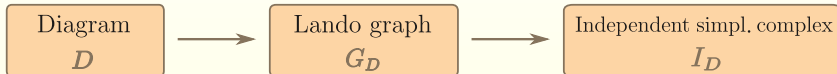


**THEOREM** [González-Meneses, Manchón, S.]

$$D \longrightarrow I_D$$
$$\{C_{i,j_{max}}(D), d_i\} \sim_h \{C_i(I_D), \partial_i\}$$



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$$H_{i,j_{\max}}(D) \underset{\text{some shifting}}{=} \tilde{H}_i(I_D)$$

# Almost extreme Khovanov complex

$$\cdots \longrightarrow C_{i+2,j_{\max}}(D) \longrightarrow C_{i,j_{\max}}(D) \longrightarrow C_{i-2,j_{\max}}(D) \longrightarrow \cdots$$

Extreme Khovanov

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Extreme Khovanov

$$\cdots \longrightarrow C_{i+2,j_{\text{almost}}}(D) \longrightarrow C_{i,j_{\text{almost}}}(D) \longrightarrow C_{i-2,j_{\text{almost}}}(D) \longrightarrow \cdots$$

Almost-extreme Khovanov

$$\vdots$$

$$\vdots$$

$$\vdots$$

$$\cdots \longrightarrow C_{i+2,j+4}(D) \longrightarrow C_{i,j+4}(D) \longrightarrow C_{i-2,j+4}(D) \longrightarrow \cdots$$

$$\cdots \longrightarrow C_{i+2,j}(D) \longrightarrow C_{i,j}(D) \longrightarrow C_{i-2,j}(D) \longrightarrow \cdots$$

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$$\vdots$$

$$\vdots$$

$$\vdots$$

$$j_{\max}(D) = \max_{s \text{ state of } D} \{j(s)\}$$

$$j_{\text{almost}}(D) = j_{\max}(D) - 4$$

$$D \longrightarrow I_D \text{ simplicial complex} \quad / \quad H_i(I_D) = H^{i,j_{\max}}(D)$$

# Almost extreme Khovanov complex

$$\cdots \longrightarrow C_{i+2,j_{\max}}(D) \longrightarrow C_{i,j_{\max}}(D) \longrightarrow C_{i-2,j_{\max}}(D) \longrightarrow \cdots$$

Extreme Khovanov

$$\cdots \longrightarrow C_{i+2,j_{\max}}(D) \longrightarrow C_{i,j_{\max}}(D) \longrightarrow C_{i-2,j_{\max}}(D) \longrightarrow \cdots$$

Almost-extreme Khovanov

$$\vdots$$

$$\vdots$$

$$\vdots$$

$$\cdots \longrightarrow C_{i+2,j+4}(D) \longrightarrow C_{i,j+4}(D) \longrightarrow C_{i-2,j+4}(D) \longrightarrow \cdots$$

$$\cdots \longrightarrow C_{i+2,j}(D) \longrightarrow C_{i,j}(D) \longrightarrow C_{i-2,j}(D) \longrightarrow \cdots$$

$$\cdots \longrightarrow C_{i+2,j-4}(D) \longrightarrow C_{i,j-4}(D) \longrightarrow C_{i-2,j-4}(D) \longrightarrow \cdots$$

$$\vdots$$

$$\vdots$$

$$\vdots$$

$$j_{\max}(D) = \max_{s \text{ state of } D} \{j(s)\}$$

$$j_{\max}(D) = j_{\max}(D) - 4$$

$$D \longrightarrow I_D \text{ simplicial complex} / H_i(I_D) = H^{i,j_{\max}}(D)$$

Is it possible to find something similar for  $H^{i,j_{\max}}(D)$ ?

# Almost extreme Khovanov complex

$$\cdots \longrightarrow C_{i+2,j_{\max}}(D) \longrightarrow C_{i,j_{\max}}(D) \longrightarrow C_{i-2,j_{\max}}(D) \longrightarrow \cdots$$

Extreme Khovanov

$$\cdots \longrightarrow C_{i+2,j_{\max}}(D) \longrightarrow C_{i,j_{\max}}(D) \longrightarrow C_{i-2,j_{\max}}(D) \longrightarrow \cdots$$

Almost-extreme Khovanov

$$\vdots$$

$$\vdots$$

$$\vdots$$

$$\cdots \longrightarrow C_{i+2,j+4}(D) \longrightarrow C_{i,j+4}(D) \longrightarrow C_{i-2,j+4}(D) \longrightarrow \cdots$$

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$$j_{\max}(D) = j_{\max}(D) - 4$$

$$D \longrightarrow I_D \text{ simplicial complex} / H_i(I_D) = H^{i,j_{\max}}(D)$$

Is it possible to find something similar for  $H^{i,j_{\max}}(D)$ ? “YES”

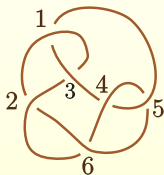
# Semiadequate links

$D$



# Semiadequate links

$D$

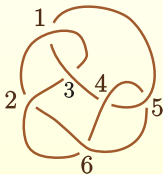


Kauffman state:

$$s : cr(D) \rightarrow \{A, B\}$$

# Semiadequate links

$D$



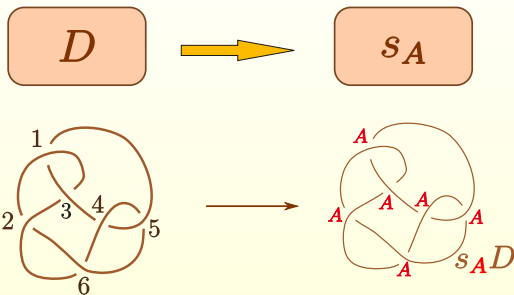
$s_A :=$  state where all crossings are marked with  $A$ -label.

Kauffman state:

$$s : cr(D) \rightarrow \{A, B\}$$



# Semiadequate links

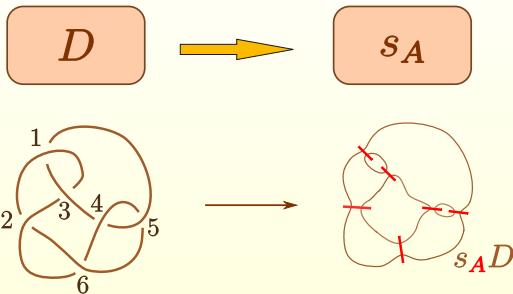


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Kauffman state:

$$s : cr(D) \rightarrow \{A, B\}$$

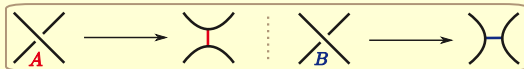
# Semiadequate links



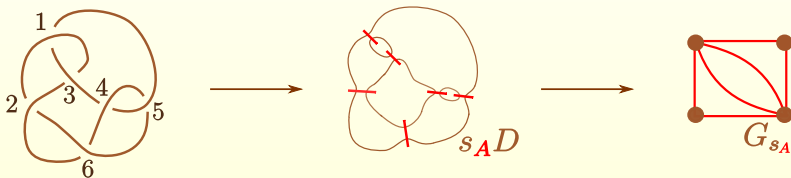
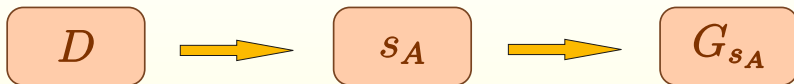
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Kauffman state:

$s : cr(D) \rightarrow \{A, B\}$



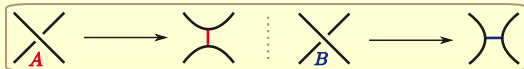
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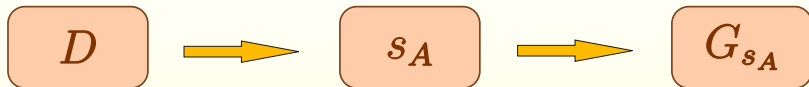
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Kauffman state:

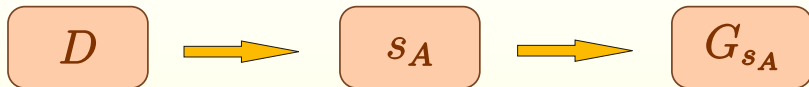
$s : cr(D) \rightarrow \{A, B\}$



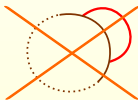
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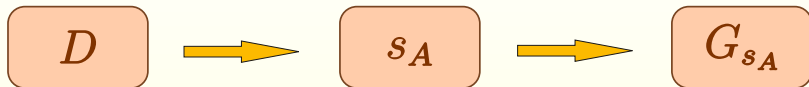
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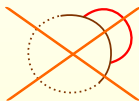
$D$  is  $A$ -adequate if



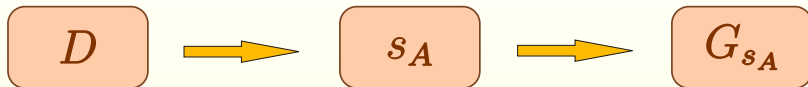
# Semiadequate links



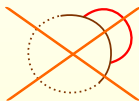
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# Semiadequate links

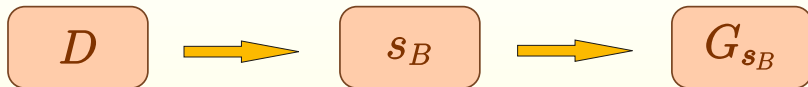


$D$  is  $A$ -adequate if

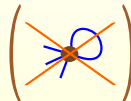
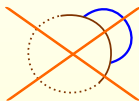


$D$   $A$ -adequate  $\longrightarrow$   $L$   $A$ -adequate

# Semiadequate links



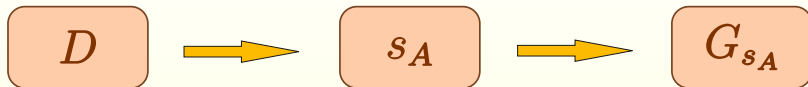
$D$  is  $B$ -adequate if



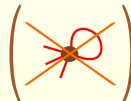
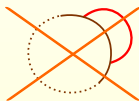
$D$   $B$ -adequate  $\longrightarrow$   $L$   $B$ -adequate



# Semiadequate links

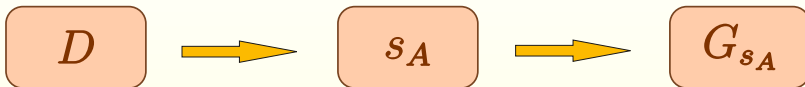


$D$  is  $A$ -adequate if

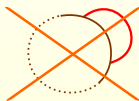


$D$   $A$ -adequate  $\longrightarrow$   $L$   $A$ -adequate

# Semiadequate links



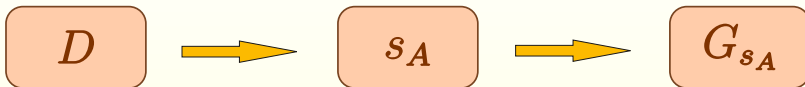
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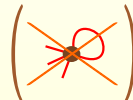
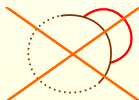
$D$   $A$ -adequate  $\longrightarrow L$   $A$ -adequate

Semiadequate  $:= A$ -adequate or  $B$ -adequate

# Semiadequate links



$D$  is  $A$ -adequate if

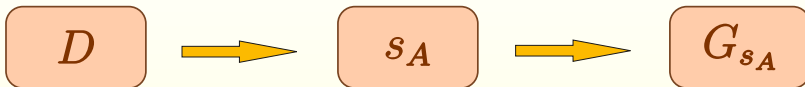


$D$   $A$ -adequate  $\longrightarrow$   $L$   $A$ -adequate

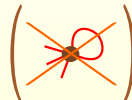
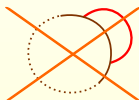
Semiadequate  $:=$   $A$ -adequate or  $B$ -adequate

Adequate  $:=$   $A$ -adequate and  $B$ -adequate

# Semiadequate links



$D$  is  $A$ -adequate if



$D$   $A$ -adequate  $\longrightarrow L$   $A$ -adequate

Semiadequate  $:= A$ -adequate or  $B$ -adequate

Adequate  $:= A$ -adequate and  $B$ -adequate

## (Partial) presimplicial sets

$$\mathcal{X} = (X_n, d_i) \left\{ \begin{array}{l} \{X_n\}_{n \geq 0} \\ d_{i,n} = d_i: \quad X_n \rightarrow X_{n-1} \quad 0 \leq i \leq n \end{array} \right.$$

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### Geometric realization

## (Partial) presimplicial sets

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### Geometric realization:

$$\Delta^n = \left\{ (x_0, \dots, x_n) \in \mathbb{R}^{n+1} \mid \sum_{i=0}^n x_i = 1, x_i \geq 0 \right\}$$

$$d^{i,n} = d^i: \Delta^{n-1} \rightarrow \Delta^n, \quad d^i(x_0, \dots, x_{n-1}) = (x_0, \dots, x_{i-1}, 0, x_i, \dots, x_{n-1})$$



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$$|\mathcal{X}| = \bigsqcup_{n \geq 0} X_n \times \Delta^n / \sim$$

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“gluing instructions”

$$\hookrightarrow (a, d^i(t)) \sim (d_i(a), t), \quad a \in X_n, \quad t \in \Delta^{n-1}$$

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$$\mathcal{X} = (X_n, d_i) \left\{ \begin{array}{ll} \{X_n\}_{n \geq 0} & \\ d_{i,n} = d_i: \text{Dom}(X_n) \rightarrow X_{n-1} & 0 \leq i < j \leq n \\ d_i d_j = d_{j-1} d_i & \text{as long as both sides are defined} \end{array} \right.$$

## Geometric realization:

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“gluing instructions”

$$\hookrightarrow (a, d^i(t)) \sim (d_i(a), t), \quad \text{if } d_i(a) \neq 0$$

$$\begin{array}{l} a \in X_n \\ t \in \Delta^{n-1} \end{array}$$

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$$|\mathcal{X}| = b \cup \bigsqcup_{n \geq 0} X_n \times \Delta^n / \sim$$

“gluing instructions”

$$\begin{array}{ll} \rightarrow (a, d^i(t)) \sim (d_i(a), t), & \text{if } d_i(a) \neq 0 \\ \rightarrow (a, d^i(t)) \sim b, & \text{if } d_i(a) = 0 \end{array}$$

$$\begin{array}{l} a \in X_n \\ t \in \Delta^{n-1} \end{array}$$

# From semiadequate link diagram to presimplicial set

$D$

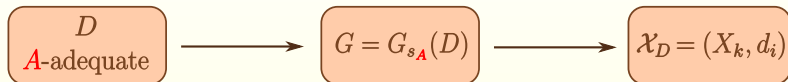
$A$ -adequate



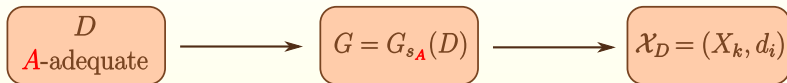
# From semiadequate link diagram to presimplicial set



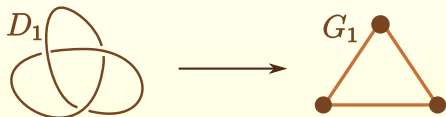
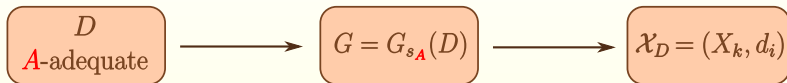
# From semiadequate link diagram to presimplicial set



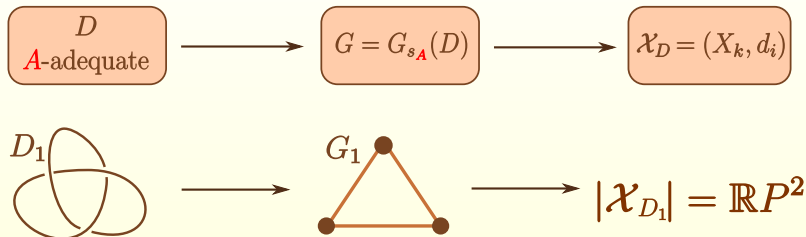
# From semiadequate link diagram to presimplicial set



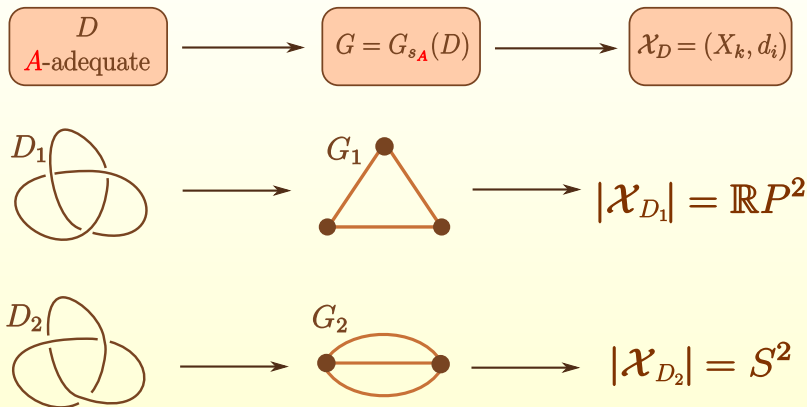
# From semiadequate link diagram to presimplicial set



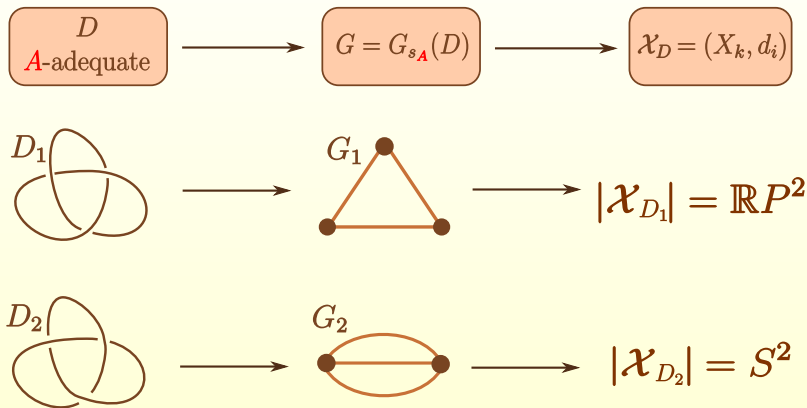
# From semiadequate link diagram to presimplicial set



# From semiadequate link diagram to presimplicial set



# From semiadequate link diagram to presimplicial set



(Construction of  $\mathcal{X}_D$  detailed in the poster)

# Almost extreme Khovanov homology of semiadequate links

I

$D$   $A$ -adequate



# Almost extreme Khovanov homology of semiadequate links

I

$$D \text{ } \textcolor{red}{A}\text{-adequate} \longrightarrow \mathcal{X}_D$$

# Almost extreme Khovanov homology of semiadequate links

## THEOREM

$$D \text{ \textcolor{red}{A}-adequate} \longrightarrow \mathcal{X}_D$$

$$\{C_{i,j_{\text{almax}}}(D), d_i\} \sim_h |\mathcal{X}_D|$$

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$$D \text{ \textcolor{red}{A}-adequate} \longrightarrow \mathcal{X}_D$$

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$$H_{i,j_{\text{almax}}}(D) \underset{\substack{\text{some} \\ \text{shifting}}}{=} \tilde{H}_i(|\mathcal{X}_D|)$$

# Almost extreme Khovanov homology of semiadequate links

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$$H_{i,j_{\text{almax}}}(D) \underset{\text{some shifting}}{=} \tilde{H}_i(|\mathcal{X}_D|)$$

$$D_1 \quad |\mathcal{X}_{D_1}| = S^2$$

$$D_2 \quad |\mathcal{X}_{D_2}| = \mathbb{R}P^2$$


# Almost extreme Khovanov homology of semiadequate links

## THEOREM


$$D \text{ \textcolor{red}{A}-adequate} \longrightarrow \mathcal{X}_D$$

$$\{C_{i,j_{\text{almax}}}(D), d_i\} \sim_h |\mathcal{X}_D|$$

$$H_{i,j_{\text{almax}}}(D) \underset{\text{some shifting}}{=} \tilde{H}_i(|\mathcal{X}_D|)$$

$D_1$    $|\mathcal{X}_{D_1}| = S^2$

		$i$			
		-3	-1	1	3
$j$	7				$\mathbb{Z}$
	3				$\mathbb{Z}$
	-1		$\mathbb{Z}$		
	-5	$\mathbb{Z}/2$			
	-9	$\mathbb{Z}$			

$D_2$    $|\mathcal{X}_{D_2}| = \mathbb{R}P^2$

		$i$			
		-3	-1	1	3
$j$	9				$\mathbb{Z}$
	5			$\mathbb{Z}/2$	
	1			$\mathbb{Z}$	
	-3	$\mathbb{Z}$			
	-7	$\mathbb{Z}$			

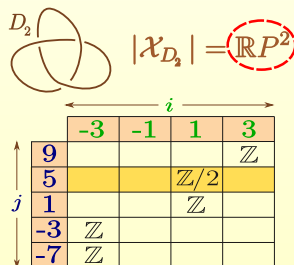
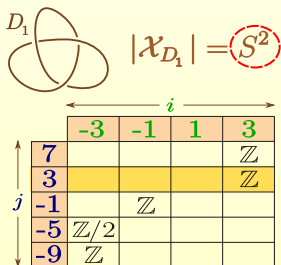
# Almost extreme Khovanov homology of semiadequate links

## THEOREM

$$D \text{ \textcolor{red}{A}-adequate} \longrightarrow \mathcal{X}_D$$

$$\{C_{i,j_{\text{almax}}}(D), d_i\} \sim_h |\mathcal{X}_D|$$

$$H_{i,j_{\text{almax}}}(D) \underset{\text{some shifting}}{=} \tilde{H}_i(|\mathcal{X}_D|)$$



# Almost-extreme Khovanov homology II

## THEOREM

$$D \text{ \textcolor{red}{A}-adequate} \longrightarrow \mathcal{X}_D$$

$$\{C_{i,j_{\text{almax}}}(D), d_i\} \sim_h |\mathcal{X}_D|$$

$$H_{i,j_{\text{almax}}}(D) \underset{\text{some shifting}}{=} \tilde{H}_i(|\mathcal{X}_D|)$$

# Almost-extreme Khovanov homology II

## THEOREM

$$D \text{ \textcolor{red}{A}-adequate} \longrightarrow \mathcal{X}_D$$

$$\{C_{i,j_{\text{almax}}}(D), d_i\} \sim_h |\mathcal{X}_D|$$

$$H_{i,j_{\text{almax}}}(D) \underset{\substack{\text{some} \\ \text{shifting}}}{=} \tilde{H}_i(|\mathcal{X}_D|)$$

I

$$D \text{ \textcolor{red}{A}-adequate} / G = G_{s_{\textcolor{red}{A}}} \text{ simple} \longrightarrow \mathcal{X}_D$$



# Almost-extreme Khovanov homology II

## THEOREM

$$D \text{ \textcolor{red}{A}-adequate} \longrightarrow \mathcal{X}_D$$

$$\{C_{i,j_{\text{almax}}}(D), d_i\} \sim_h |\mathcal{X}_D|$$

$$H_{i,j_{\text{almax}}}(D) \underset{\substack{\text{some} \\ \text{shifting}}}{=} \tilde{H}_i(|\mathcal{X}_D|)$$

I

$$D \text{ \textcolor{red}{A}-adequate} / \begin{array}{l} c \text{ crossings} \\ G = G_{s_{\textcolor{red}{A}}} \text{ simple} \\ p_1(G) \text{ cyclomatic number} \end{array} \longrightarrow \mathcal{X}_D$$

# Almost-extreme Khovanov homology II

## THEOREM

$$D \text{ \textcolor{red}{A}-adequate} \longrightarrow \mathcal{X}_D$$

$$\{C_{i,j_{\text{almax}}}(D), d_i\} \sim_h |\mathcal{X}_D|$$

$$H_{i,j_{\text{almax}}}(D) \underset{\text{some shifting}}{=} \tilde{H}_i(|\mathcal{X}_D|)$$

I

$$\begin{array}{ccc} D \text{ \textcolor{red}{A}-adequate} & / & G = G_{s_{\textcolor{red}{A}}} \text{ simple} \longrightarrow \mathcal{X}_D \\ \text{\textcolor{brown}{c} crossings} & & \text{\textcolor{green}{p}_1(G) cyclomatic number} \end{array}$$

$$|\mathcal{X}_D| \sim_h \begin{cases} \bigvee_{\text{\textcolor{green}{p}_1(G)}} S^{\textcolor{brown}{c}-1} & \text{if } G \text{ is bipartite.} \end{cases}$$

# Almost-extreme Khovanov homology II

## THEOREM

$$D \text{ \textcolor{red}{A}-adequate} \longrightarrow \mathcal{X}_D$$

$$\{C_{i,j_{\text{almax}}}(D), d_i\} \sim_h |\mathcal{X}_D|$$

$$H_{i,j_{\text{almax}}}(D) \stackrel{\text{some shifting}}{=} \tilde{H}_i(|\mathcal{X}_D|)$$

## THEOREM

$$D \text{ \textcolor{red}{A}-adequate} / \begin{matrix} c \text{ crossings} \\ G = G_{s_{\textcolor{red}{A}}} \text{ simple} \\ p_1(G) \text{ cyclomatic number} \end{matrix} \longrightarrow \mathcal{X}_D$$

$$|\mathcal{X}_D| \sim_h \begin{cases} \bigvee_{p_1(G)} S^{\textcolor{red}{c}-1} & \text{if } G \text{ is bipartite.} \\ \bigvee_{p_1(G)-1} S^{\textcolor{red}{c}-1} \vee \Sigma^{\textcolor{red}{c}-3} \mathbb{R}P^2 & \text{otherwise.} \end{cases}$$

# Almost-extreme Khovanov homology III

## COROLLARY

$D$   $A$ -adequate /  $G = G_{s_A}$  simple  
 $c$  crossings  $p_1(G)$  cyclomatic number

# Almost-extreme Khovanov homology III

## COROLLARY

$D$   $A$ -adequate /  $G = G_{s_A}$  simple  
 $c$  crossings  $p_1(G)$  cyclomatic number

$$H_{i, \underbrace{c+2|s_A|-4}_{j_{almax}}}(D) \simeq \begin{cases} \mathbb{Z}^{p_1} & \text{if } G \text{ is bipartite.} \\ \mathbb{Z}^{p_1-1} \oplus \mathbb{Z}_2 & \text{otherwise.} \end{cases}$$

# Almost-extreme Khovanov homology III

## COROLLARY

$D$   $\textcolor{red}{A}$ -adequate /  $G = G_{s\textcolor{red}{A}}$  simple  
 $c$  crossings  $p_1(G)$  cyclomatic number

$$H_{i, \underbrace{c+2|s\textcolor{red}{A}|-4}_{j_{almax}}}(D) \simeq \begin{cases} \mathbb{Z}^{\textcolor{green}{p}_1} & \text{if } G \text{ is bipartite.} \\ \mathbb{Z}^{\textcolor{green}{p}_1-1} \oplus \mathbb{Z}_2 & \text{otherwise.} \end{cases}$$



10<sub>44</sub>

# Almost-extreme Khovanov homology III

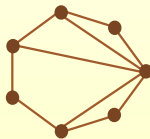
## COROLLARY

$D$   $\textcolor{red}{A}$ -adequate /  $G = G_{s\textcolor{red}{A}}$  simple  
 $c$  crossings  $p_1(G)$  cyclomatic number

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$10_{44}$



$G$

# Almost-extreme Khovanov homology III

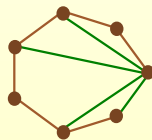
## COROLLARY

$D$   $\mathbf{A}$ -adequate /  $G = G_{s\mathbf{A}}$  simple  
 $c$  crossings  $p_1(G)$  cyclomatic number

$$H_{i, \underbrace{c+2|s\mathbf{A}|-4}_{j_{almax}}}(D) \simeq \begin{cases} \mathbb{Z}^{p_1} & \text{if } G \text{ is bipartite.} \\ \mathbb{Z}^{p_1-1} \oplus \mathbb{Z}_2 & \text{otherwise.} \end{cases}$$



$10_{44}$



$G$   $p_1 = 4$

$$H_{i,20}(10_{44}) = \mathbb{Z}^3 \oplus \mathbb{Z}_2$$







## Khovanov homology of semiadequate links via presimplicial sets.

Marithania Silveira - Barcelona Graduate School of Mathematics at Universitat de Barcelona - marithania@ub.es

Kieff H. Przytycki - The George Washington University - przytycki@gwu.edu

**Abstract:** Given a semiadequate diagram  $D$  representing a link  $L$ , we present an algorithm for constructing a presimplicial set such that its geometric realization is homotopy equivalent to the almost-extreme Khovanov complex of  $L$ . Moreover, we determine the homotopy type of the presimplicial set that one obtains when the link is strongly  $A$ -adequate.

### Presimplicial sets

A (partial) presimplicial set  $X = (X_n, \partial_n)$  consists of a sequence of sets  $(X_n)_{n \geq 0}$  together with a collection of partial maps  $\partial_n$  defined only on subsets of  $X_n$ ,  $\partial_n: X_n \rightarrow X_{n-1}$ ,  $\partial_n(x) = \partial_n(y)$  if and only if  $x \sim y$  and  $\partial_n(x) = \partial_n(y)$  for all  $x, y \in X_n$  such that  $x \sim y$ .

### Generative realizations

Let  $\mathcal{D}^n = \{x_1, \dots, x_n\} \subset \mathbb{R}^n$  with  $x_i = (x_{i1}, \dots, x_{in})$  be a "model simplex" and set  $\mathcal{D}^n = \mathcal{D}^n \times \mathcal{D}^n \rightarrow \mathcal{D}^n$ ,  $\mathcal{D}^n(x_1, \dots, x_n) = (x_{11}, \dots, x_{n1})$ . The generative realizations of  $(X, \partial_n)$  are the CW-complexes

$$|X| = \bigcup_{n \geq 0} |X_n| \times \mathcal{D}^n$$

where, for  $x \in X_n$  and  $t \in \mathcal{D}^{n-1}$ , the relation  $\sim$  is given by  $(x, t) \sim (x, t')$  if  $\partial_n(x) = t$  and  $\partial_n(x) = t'$ .

### Khovanov homology

For each of the  $\mathcal{D}^{n-1}$  Khovanov states associated to  $D$ , define  $\sigma(x) = \#A - \#B$ .

For the state  $x$ , consider each crossing as shown.

Reduce it by giving a sign to each of the cycles obtained, and define  $\sigma(x) = \#A - \#B$ .

For the enhanced state  $x$ , define  $\sigma(x) = \#A - \#B + \sum_{i=1}^n \sigma(x_i)$ .

The enhanced state  $x$  is adjacent to  $y$  if and only if  $\sigma(x) = \sigma(y) \pm 1$ .

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### Semiadequate links

Let  $D$  be a link diagram with  $n$  crossings.

A link diagram is semiadequate if it is the result of a crossing change of  $D$  with a letter  $A$  or  $B$ .

For each crossing following the figure on the left, we obtain  $A$  or  $B$ .

The state graph  $G_D$  associated to  $D$  is obtained from  $D$  by collapsing each circle to a vertex.

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### Almost-extreme Khovanov homology

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## Khovanov homology of semiadequate links via presimplicial sets.

Marithania Silveira - Barcelona Graduate School of Mathematics at Universitat de Barcelona - marithania@ub.es

Kieff H. Przytycki - The George Washington University - przytycki@gwu.edu

**Abstract:** Given a semiadequate diagram  $D$  representing a link  $L$ , we present an algorithm for constructing a presimplicial set such that its geometric realization is homotopy equivalent to the almost-extreme Khovanov complex of  $L$ . Moreover, we determine the homotopy type of the presimplicial set that one obtains when the link is strongly  $A$ -adequate.

### Presimplicial sets

A (partial) presimplicial set  $X = (X_n, \partial_n)$  consists of a sequence of sets  $(X_n)_{n \geq 0}$  together with a collection of partial maps  $\partial_n$  defined only on subsets of  $X_n$ ,  $\partial_n: X_n \rightarrow X_{n-1}$ ,  $\partial_n(x) = \partial_n(y)$  if and only if  $x \sim y$  and  $0 \leq i < j \leq n$ , so long as both sides of the equality are defined.

### Generative realizations

Let  $\mathcal{D}^n = \{x_1, \dots, x_n\} \subset \mathbb{R}^{n+1}$  with  $x_i = (x_i, 1)$  be a "model simplex" and set  $\mathcal{D}^n = \mathcal{D}^n \cup \partial \mathcal{D}^n \subset \mathbb{R}^{n+1}$ . The generators  $\partial_n$  of  $(X, \partial_n)$  are the  $\mathbb{R}^n$ -complexes

$$\partial_n(x) = \bigcup_{i=1}^n \partial_i \mathcal{D}^n$$

where, for  $x \in X_n$  and  $i \in \mathcal{D}^{n-1}$ , the relation  $\sim$  is given by  $(x, \partial_i(x)) \sim (y, \partial_i(y))$  if and only if  $(x, \partial_i(x)) \sim (y, \partial_i(y))$ .

### Khovanov homology

For each of the  $\mathcal{D}^{n-1}$  Khovanov states associated to  $\mathcal{D}$ , define  $\pi(x) = \partial_i(x) - \partial_j(x)$ .

For the state  $\pi$ , consider each crossing as shown.

Reduce by giving a sign to each of the cycles obtained, and define  $\pi(x) = \pi(y) - \pi(z)$ .

For the enhanced state  $\pi$ , define

$$\pi(x) = \pi(y) - \pi(z) + \pi(w)$$

The enhanced state  $\pi$  is adjacent to  $\pi'$  if and only if  $\pi(x) = \pi'(x) + \pi(y) - \pi(z) - \pi(w)$ .

The algebra of the crossing divides  $\pi$  and  $\pi'$  are equal.

For all numbers  $\pi$  and  $\pi'$  are adjacent except in one (sign) crossing  $\pi = \pi' + 1$ .

### Combinatorics

$\mathcal{C}^n(D)$  is free abelian group generated by the set of enhanced states  $\pi$  of  $D$  with  $\pi(x) = \pi(y) - \pi(z) - \pi(w)$ .

$$\partial_n: \mathcal{C}^n(D) \rightarrow \mathcal{C}^{n-1}(D)$$

$$R(x) = \sum_{i=1}^n (-1)^{i-1} \pi(x)$$

with  $\pi(x) = \pi(y) - \pi(z) - \pi(w)$ .

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### Semiadequate links

Let  $D$  be a link diagram with  $n$  crossings.

A (partial) state is the result of smoothing each crossing of  $D$  with a letter  $A$  or  $B$ .

Show each crossing following the figure on the left to obtain a state.

The state graph  $G_D$  associated to  $D$  is obtained from  $\mathcal{C}^0(D)$  by collapsing each circle to a vertex.

Write  $\pi_1, \pi_2, \dots$  for the states where all crossings are marked with  $A$  or  $B$ .

$D$  is  $A$ -adequate if the graph  $G_D = G_D(\pi_1, \pi_2, \dots)$  has no loops.

$D$  is  $B$ -adequate if the graph  $G_D = G_D(\pi_1, \pi_2, \dots)$  has no loops.

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Merci beaucoup.