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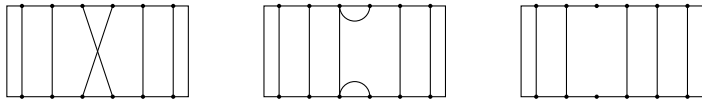
## The Partition algebra and the Kronecker coefficients II:

### The Littlewood–Richardson rule

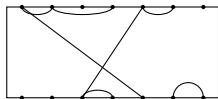
## Section 1

The partition algebra and the stable Kronecker coefficients

The partition algebra,  $P_s(n)$ , is the  $\mathbb{C}$ -algebra generated by elements of the form

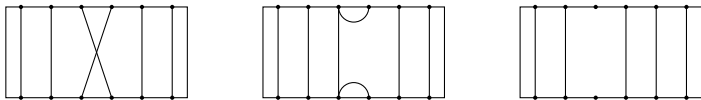


with product given by concatenation of diagrams (deformed by the parameter  $n$ ). For example,

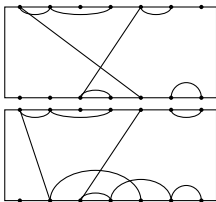


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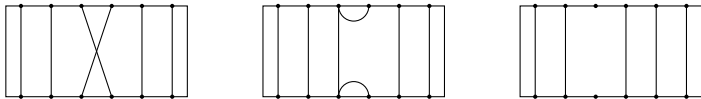


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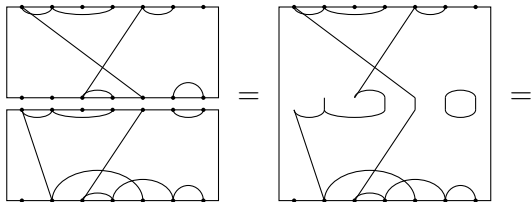


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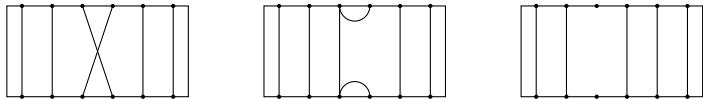


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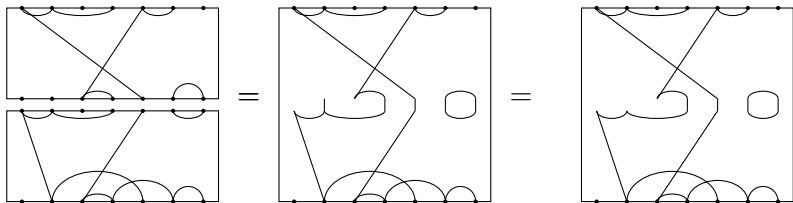


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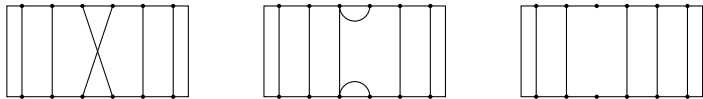


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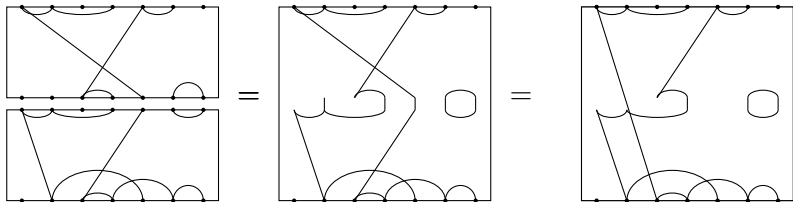


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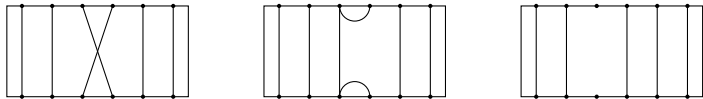


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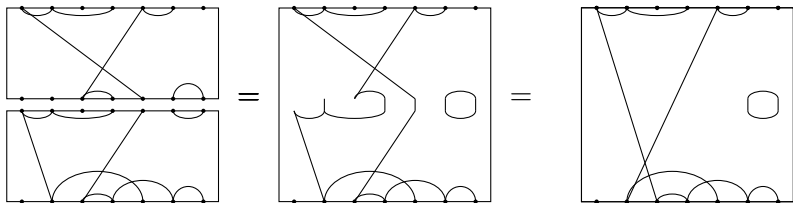


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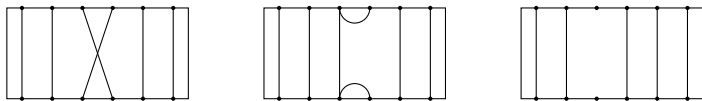
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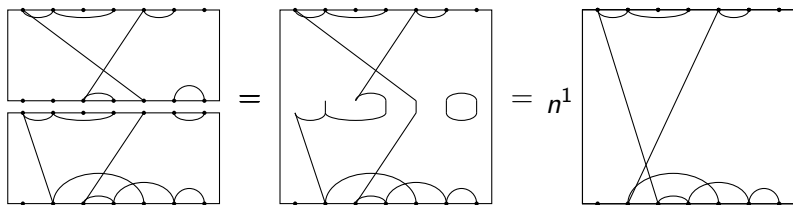
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# The branching graph of the tower of partition algebras



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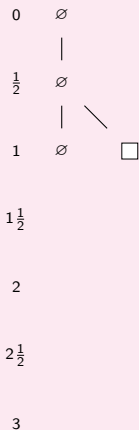
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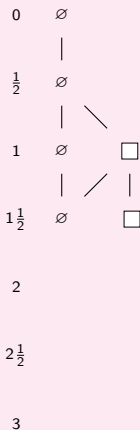
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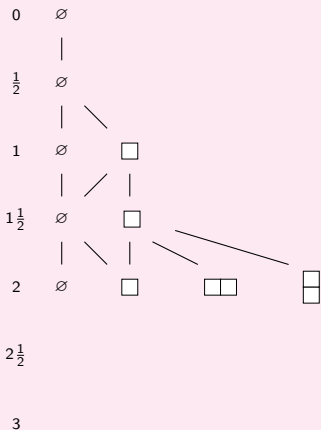
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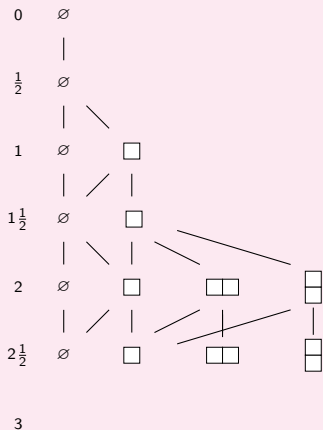
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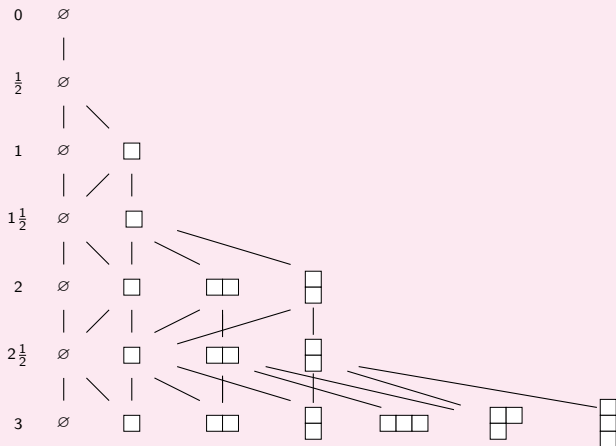
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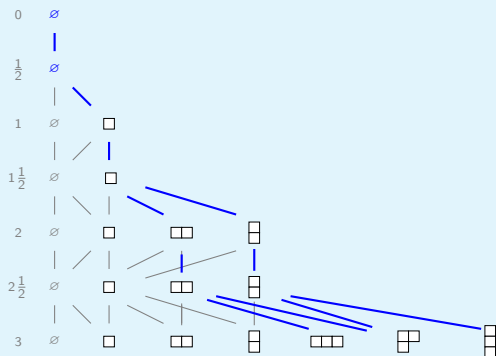
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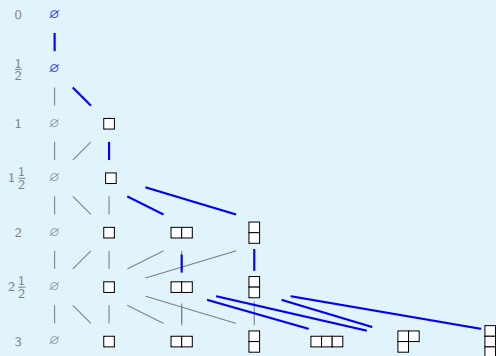
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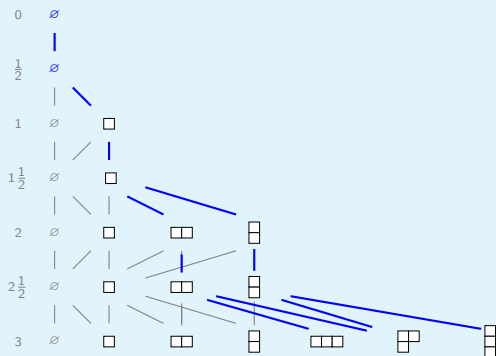
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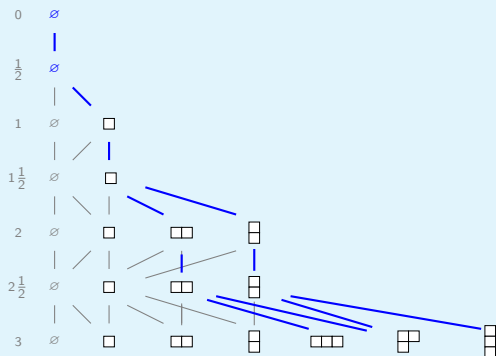
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The stable Kronecker coefficients equal dimensions of homomorphism spaces for path-theoretic  $P_s(n)$ -modules:

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- this is done via combinatorial resolutions.
- we do this in a language which is generalisable to the whole  $P_s(n)$  branching graph.

## The Littlewood–Richardson rule

Let  $\lambda \vdash r - s$ ,  $\mu \vdash s$  and  $\nu \vdash r$ . The multiplicities,

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- the Young tableau is semistandard;
- the  $\mu$ -reverse reading word of the Young tableau is a lattice permutation.

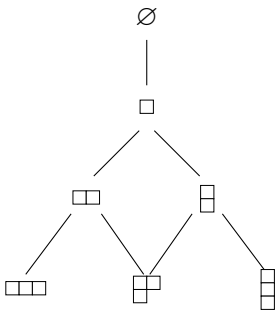
## This lecture

We unpack these terms and recast them in a manner which can be generalised to the wider Kronecker coefficients.

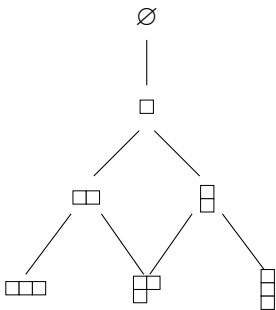
## Section 2

Standard tableaux and representations of  
symmetric groups

- The symmetric groups are controlled by the graph:

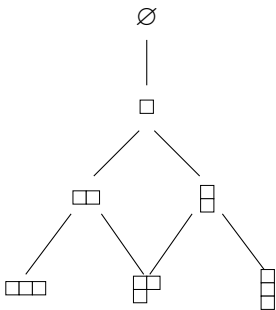


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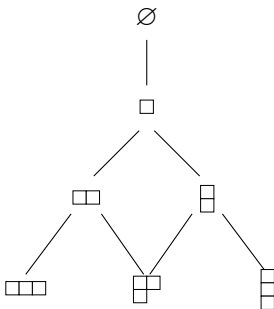


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- This gives a nice basis

$$\mathbb{C}\mathfrak{S}_s = \mathbb{C}\{m_{st} \mid s, t \in \text{Std}(\lambda), \lambda \vdash s\}$$

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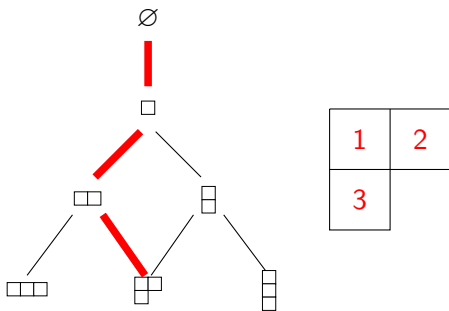
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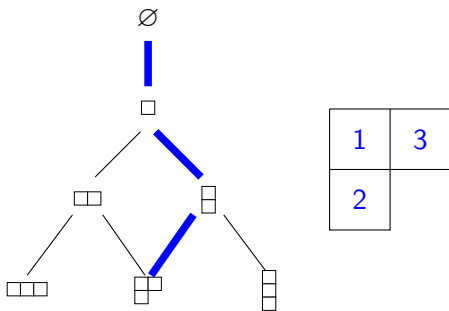
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## Definition

Each point  $\lambda \vdash s$  in the graph labels one left and one 2-sided ideal of the algebra. Fix any  $t^\lambda \in \text{Std}(\lambda)$ , we define

$$M^\lambda = \mathbb{C}\mathfrak{S}_s m_{t^\lambda t^\lambda} = \text{ind}_{\mathfrak{S}_{\lambda_1} \times \mathfrak{S}_{\lambda_2} \times \dots}^{\mathfrak{S}_n}(\mathbb{C})$$

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We define the Specht module  $S(\lambda)$  to be the quotient

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$$S(\nu \setminus \lambda) = \mathbb{C}\{m_s \mid s \in \text{Std}(\nu \setminus \lambda)\}$$

## Action on skew modules

Given  $t \in \text{Std}_s(\nu \setminus \lambda)$ , say we let  $t_{k \leftrightarrow k+1}$  denote the tableau with  $k$  and  $k+1$  swapped. We have that

$$s_k(t) = \begin{cases} t_{k \leftrightarrow k+1} & \text{if } k \leftrightarrow k+1 \text{ is a standard tableau} \\ t & \text{if } k \text{ \& } k+1 \text{ are in same row} \\ -t + \sum_{s \triangleright t} a_s s & \text{if } k \text{ \& } k+1 \text{ are in same column} \end{cases}$$

## Example

$$S((2, 2) \setminus (1)) = \left\{ \begin{array}{|c|c|} \hline & 1 \\ \hline 2 & 3 \\ \hline \end{array}, \begin{array}{|c|c|} \hline & 2 \\ \hline 1 & 3 \\ \hline \end{array} \right\}$$

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$$s_1 = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \quad s_2 = \begin{pmatrix} 1 & -1 \\ 0 & -1 \end{pmatrix}$$

## Section 3

# Semistandard Young tableaux and homomorphisms

## Question

Can we give a basis of

$$\text{Hom}_{\mathbb{C}\mathfrak{S}_s}(M^\mu, \Delta_s(\nu \setminus \lambda))$$

and then count how many of these homomorphisms factor through the projection

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## Question

- The homomorphisms are indexed by and constructed from SEMISTANDARD tableaux.
- Those which factor through  $\pi$  are counted by the semistandard tableaux satisfying the LATTICE PERMUTATION condition.

A semistandard tableau of shape  $\nu \setminus \lambda$  and weight  $\mu$  is a filling of the boxes of  $\nu \setminus \lambda$  with the entries

$$\underbrace{1, \dots, 1}_{\mu_1}, \underbrace{2, \dots, 2}_{\mu_2}, \dots, \underbrace{\ell, \dots, \ell}_{\mu_\ell}$$

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What really **IS** a semistandard tableaux?

We think of it as an orbit under the  $\mathfrak{S}_{\mu_1} \times \mathfrak{S}_{\mu_2} \times \dots$  permutation action on  $\text{Std}(\nu \setminus \lambda)$ .

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So mapping the generator of  $M(\mu)$  to the sum over these four tableaux is a  $\mathfrak{S}_5$ -homomorphism by Frobenius reciprocity.

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- We define a tableau,  $T$ , of weight  $\mu$  to be an **equivalence class of standard tableaux** under  $\stackrel{\mu}{\sim}$ .
- $T$  is **semistandard** if  $t_{k \leftrightarrow k+1} \in \text{Std}(\nu \setminus \lambda)$  for all  $t \in T$  and all  $k_1, \dots, k_{d-1} \notin \{\mu_1, \mu_1 + \mu_2, \dots\}$ .

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$$\varphi \left( \begin{array}{|c|c|} \hline 1 & 2 \\ \hline 3 & 4 \\ \hline 5 & \\ \hline \end{array} \right) = \begin{array}{|c|c|c|} \hline & & 1 \\ \hline & 2 & 3 \\ \hline 4 & 5 & \\ \hline \end{array} + \begin{array}{|c|c|c|} \hline & & 2 \\ \hline & 1 & 3 \\ \hline 4 & 5 & \\ \hline \end{array} + \begin{array}{|c|c|c|} \hline & & 1 \\ \hline & 2 & 4 \\ \hline 3 & 5 & \\ \hline \end{array} + \begin{array}{|c|c|c|} \hline & & 2 \\ \hline & 1 & 4 \\ \hline 3 & 5 & \\ \hline \end{array}$$

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## Theorem

For  $T \in \text{SStd}_s(\nu \setminus \lambda, \mu)$ , define

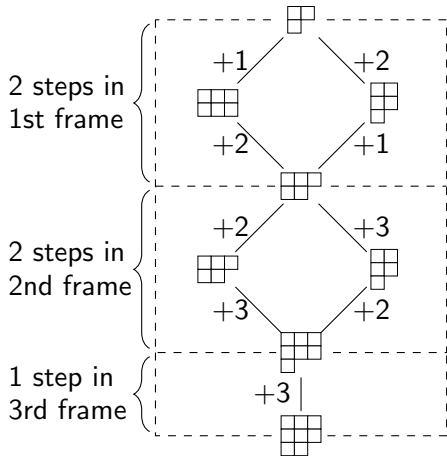
$$\varphi_T(\mathbf{t}^\mu) = \sum_{\varphi(\mathbf{t})=T} \mathbf{t}.$$

We have that

$$\{\varphi_T \mid T \in \text{SStd}_s(\nu \setminus \lambda, \mu)\}$$

is a  $\mathbb{Z}$ -basis for  $\text{Hom}_{\mathfrak{S}_s}(M(\mu), \Delta_s(\nu \setminus \lambda))$ .

How can we picture this in a way that generalises to  $P_s(n)$ ?



$$\left[ \begin{array}{|c|c|c|} \hline & & 1 \\ \hline & 1 & 2 \\ \hline 2 & 3 & \\ \hline \end{array} \right]_{\mu} = \left\{ \begin{array}{|c|c|c|} \hline & & 1 \\ \hline & 2 & 3 \\ \hline 4 & 5 & \\ \hline \end{array} , \begin{array}{|c|c|c|} \hline & & 2 \\ \hline & 1 & 3 \\ \hline 4 & 5 & \\ \hline \end{array} , \begin{array}{|c|c|c|} \hline & & 1 \\ \hline & 2 & 4 \\ \hline 3 & 5 & \\ \hline \end{array} , \begin{array}{|c|c|c|} \hline & & 2 \\ \hline & 1 & 4 \\ \hline 3 & 5 & \\ \hline \end{array} \right\}$$

## Section 4

### The lattice permutation condition

## Reverse reading word

Given a semistandard tableau,  $S$ , the reverse reading word  $\text{read}(S)$ , is given by reading the entries of  $S$  from **right-to-left** beginning with the first row and continuing in order down the rows.



## Reverse reading word

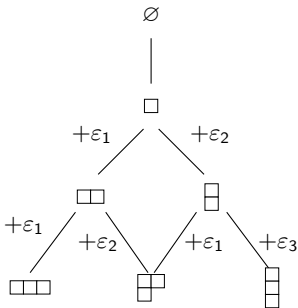
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$$\text{read} \left( \begin{array}{|c|c|c|} \hline & & 1 \\ \hline & 1 & 2 \\ \hline 2 & 3 & \\ \hline \end{array} \right) = 1, 2, 1, 3, 2$$

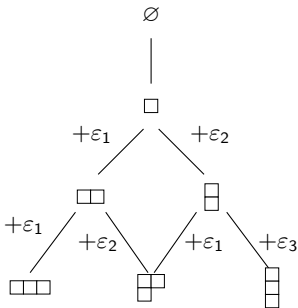
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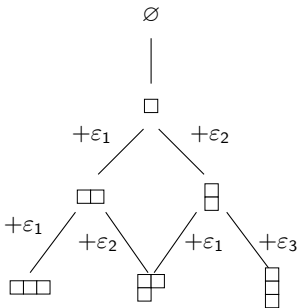


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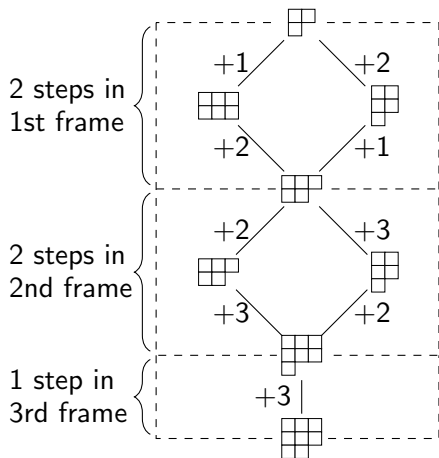


- The dominance ordering on steps is  $+\varepsilon_1 < +\varepsilon_2 < +\varepsilon_3 < \dots$

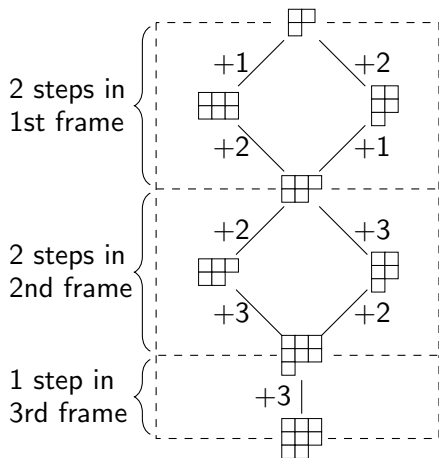
- What meaning does the ordering defining  $\text{read}(S)$  have?
- This comes from the dominance ordering on  $\mathbb{CG}_s$ .
- We decorate the edges of the branching graph with the corresponding row in which a node is added.



- The dominance ordering on steps is  $+\varepsilon_1 < +\varepsilon_2 < +\varepsilon_3 < \dots$
- We read steps in  $S$  according to the dominance ordering and then we refine this by recording the frames in which these steps occur in decreasing fashion.

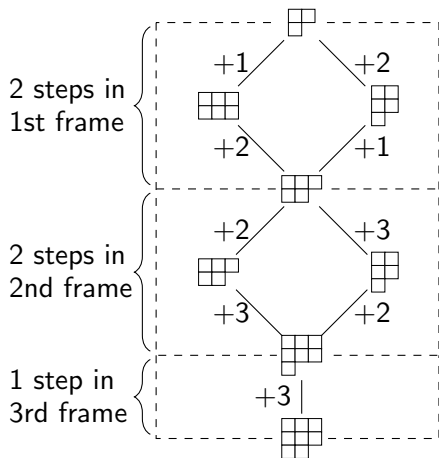


$$\left( \begin{array}{c|cc|cc} +\varepsilon_1 & +\varepsilon_2 & +\varepsilon_2 & +\varepsilon_3 & +\varepsilon_3 \\ \hline \end{array} \right)$$

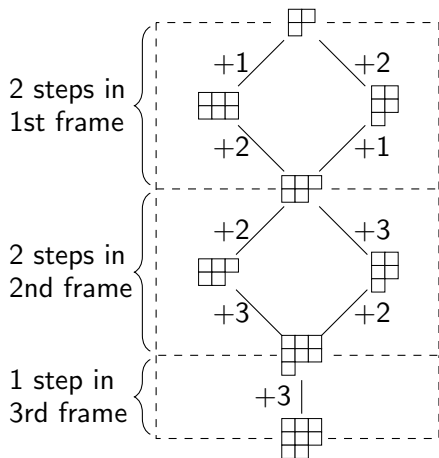


$$\left( \begin{array}{c|cc|cc} +\varepsilon_1 & +\varepsilon_2 & +\varepsilon_2 & +\varepsilon_3 & +\varepsilon_3 \\ 1 & & & & \end{array} \right)$$

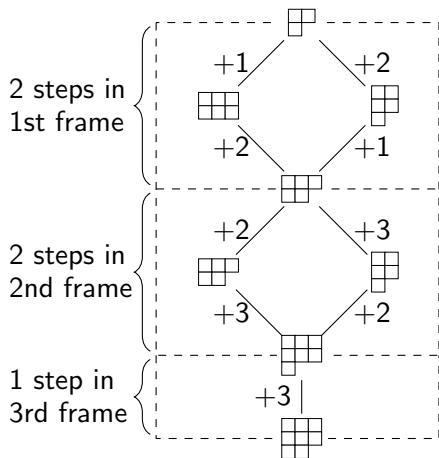




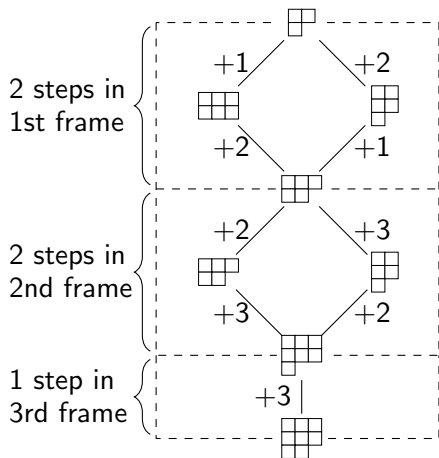
$$\left( \begin{array}{c|cc|cc} +\varepsilon_1 & +\varepsilon_2 & +\varepsilon_2 & +\varepsilon_3 & +\varepsilon_3 \\ 1 & & & & \end{array} \right)$$



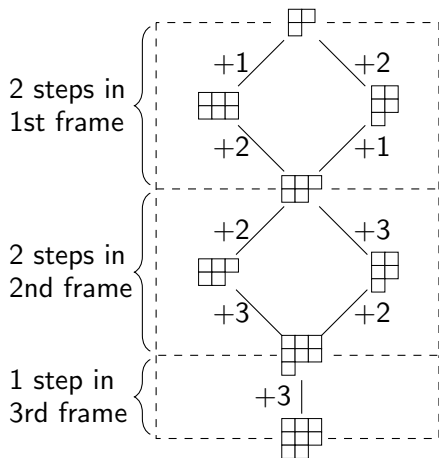
$$\left( \begin{array}{c|cc|cc} +\varepsilon_1 & +\varepsilon_2 & +\varepsilon_2 & +\varepsilon_3 & +\varepsilon_3 \\ 1 & 2 & & & \end{array} \right)$$



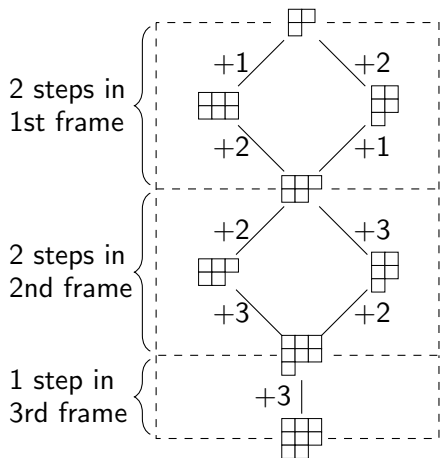
$$\left( \begin{array}{c|cc|cc} +\varepsilon_1 & +\varepsilon_2 & +\varepsilon_2 & +\varepsilon_3 & +\varepsilon_3 \\ \hline 1 & 2 & 1 & & \end{array} \right)$$



$$\left( \begin{array}{c|cc|cc} +\varepsilon_1 & +\varepsilon_2 & +\varepsilon_2 & +\varepsilon_3 & +\varepsilon_3 \\ \hline 1 & 2 & 1 & 3 & 3 \end{array} \right)$$



$$\left( \begin{array}{c|cc|cc} +\varepsilon_1 & +\varepsilon_2 & +\varepsilon_2 & +\varepsilon_3 & +\varepsilon_3 \\ \hline 1 & 2 & 1 & 3 & 2 \end{array} \right)$$



read  $\left( \begin{array}{ccc} & & 1 \\ & 1 & 2 \\ 2 & 3 & \end{array} \right) = 1, 2, 1, 3, 2$

$$\left( \begin{array}{c|cc|cc} +\varepsilon_1 & +\varepsilon_2 & +\varepsilon_2 & +\varepsilon_3 & +\varepsilon_3 \\ \hline 1 & 2 & 1 & 3 & 2 \end{array} \right)$$

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Now that we have defined the reading word of a tableau, we are ready to define the quality (good/bad) of each term in  $\text{read}(S)$  as follows.

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A sequence of positive integers is called a lattice permutation if every term in the sequence is good.

		1
	1	2
2	3	

$$\left( \begin{array}{c|cc|cc} +\varepsilon_1 & & & & \\ \hline 1 & & & & \\ \hline +\varepsilon_2 & +\varepsilon_2 & & +\varepsilon_3 & +\varepsilon_3 \\ \hline 2 & 1 & & 3 & 2 \\ \hline \end{array} \right)$$

		1
	1	2
2	3	

$$\left( \begin{array}{c|c|c|c|c} +\varepsilon_1 & +\varepsilon_2 & +\varepsilon_2 & +\varepsilon_3 & +\varepsilon_3 \\ \hline 1 & 2 & 1 & 3 & 2 \end{array} \right)$$

		1
	1	3
2	2	

		1
	1	2
2	3	

$$\left( \begin{array}{c|cc|cc} +\varepsilon_1 & +\varepsilon_2 & +\varepsilon_2 & +\varepsilon_3 & +\varepsilon_3 \\ 1 & 2 & 1 & 3 & 2 \end{array} \right)$$

		1
	1	3
2	2	

$$\left( \begin{array}{c|cc|cc} +\varepsilon_1 & +\varepsilon_2 & +\varepsilon_2 & +\varepsilon_3 & +\varepsilon_3 \\ 1 & 3 & 1 & 2 & 2 \end{array} \right)$$

		1
	1	2
2	3	

$$\left( \begin{array}{c|cc|cc} +\varepsilon_1 & +\varepsilon_2 & +\varepsilon_2 & +\varepsilon_3 & +\varepsilon_3 \\ \hline 1 & 2 & 1 & 3 & 2 \end{array} \right)$$

		1
	1	3
2	2	

$$\left( \begin{array}{c|cc|cc} +\varepsilon_1 & +\varepsilon_2 & +\varepsilon_2 & +\varepsilon_3 & +\varepsilon_3 \\ \hline 1 & \mathbf{3} & 1 & 2 & 2 \end{array} \right)$$

		1
	2	2
1	3	

		1
	1	2
2	3	

$$\left( \begin{array}{c|cc|cc} +\varepsilon_1 & +\varepsilon_2 & +\varepsilon_2 & +\varepsilon_3 & +\varepsilon_3 \\ 1 & 2 & 1 & 3 & 2 \end{array} \right)$$

		1
	1	3
2	2	

$$\left( \begin{array}{c|cc|cc} +\varepsilon_1 & +\varepsilon_2 & +\varepsilon_2 & +\varepsilon_3 & +\varepsilon_3 \\ 1 & 3 & 1 & 2 & 2 \end{array} \right)$$

		1
	2	2
1	3	

$$\left( \begin{array}{c|cc|cc} +\varepsilon_1 & +\varepsilon_2 & +\varepsilon_2 & +\varepsilon_3 & +\varepsilon_3 \\ 1 & 2 & 2 & 3 & 1 \end{array} \right)$$

		1
	1	2
2	3	

$$\left( \begin{array}{c|cc|cc} +\varepsilon_1 & +\varepsilon_2 & +\varepsilon_2 & +\varepsilon_3 & +\varepsilon_3 \\ 1 & 2 & 1 & 3 & 2 \end{array} \right)$$

		1
	1	3
2	2	

$$\left( \begin{array}{c|cc|cc} +\varepsilon_1 & +\varepsilon_2 & +\varepsilon_2 & +\varepsilon_3 & +\varepsilon_3 \\ 1 & 3 & 1 & 2 & 2 \end{array} \right)$$

		1
	2	2
1	3	

$$\left( \begin{array}{c|cc|cc} +\varepsilon_1 & +\varepsilon_2 & +\varepsilon_2 & +\varepsilon_3 & +\varepsilon_3 \\ 1 & 2 & 2 & 3 & 1 \end{array} \right)$$

		2
	1	3
1	2	



		1
	1	2
2	3	

$$\left( \begin{array}{c|cc|cc} +\varepsilon_1 & +\varepsilon_2 & +\varepsilon_2 & +\varepsilon_3 & +\varepsilon_3 \\ 1 & 2 & 1 & 3 & 2 \end{array} \right)$$

		1
	1	3
2	2	

$$\left( \begin{array}{c|cc|cc} +\varepsilon_1 & +\varepsilon_2 & +\varepsilon_2 & +\varepsilon_3 & +\varepsilon_3 \\ 1 & 3 & 1 & 2 & 2 \end{array} \right)$$

		1
	2	2
1	3	

$$\left( \begin{array}{c|cc|cc} +\varepsilon_1 & +\varepsilon_2 & +\varepsilon_2 & +\varepsilon_3 & +\varepsilon_3 \\ 1 & 2 & 2 & 3 & 1 \end{array} \right)$$

		2
	1	3
1	2	

$$\left( \begin{array}{c|cc|cc} +\varepsilon_1 & +\varepsilon_2 & +\varepsilon_2 & +\varepsilon_3 & +\varepsilon_3 \\ 2 & 3 & 1 & 2 & 1 \end{array} \right)$$

		1
	1	2
2	3	

$$\left( \begin{array}{c|cc|cc} +\varepsilon_1 & +\varepsilon_2 & +\varepsilon_2 & +\varepsilon_3 & +\varepsilon_3 \\ \hline 1 & 2 & 1 & 3 & 2 \end{array} \right)$$

		1
	1	3
2	2	

$$\left( \begin{array}{c|cc|cc} +\varepsilon_1 & +\varepsilon_2 & +\varepsilon_2 & +\varepsilon_3 & +\varepsilon_3 \\ \hline 1 & 3 & 1 & 2 & 2 \end{array} \right)$$

		1
	2	2
1	3	

$$\left( \begin{array}{c|cc|cc} +\varepsilon_1 & +\varepsilon_2 & +\varepsilon_2 & +\varepsilon_3 & +\varepsilon_3 \\ \hline 1 & 2 & 2 & 3 & 1 \end{array} \right)$$

		2
	1	3
1	2	

$$\left( \begin{array}{c|cc|cc} +\varepsilon_1 & +\varepsilon_2 & +\varepsilon_2 & +\varepsilon_3 & +\varepsilon_3 \\ \hline 2 & 3 & 1 & 2 & 1 \end{array} \right)$$

And so

$$c((3^2, 2), (2, 1), (2^2, 1)) = 1.$$