

Uniform generation of infinite concurrent runs

The case of trace monoids

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Contents

- 1 Introduction: Heaps of pieces and trace monoids
- 2 Simulating Bernoulli distributions
- 3 Step-by-step simulation and pyramids
- 4 Conclusion

Heaps of pieces and trace monoids

Heap of pieces

- Pieces:



Trace monoid

- Alphabet:

$$\Sigma = \{a, b, c, d\}$$

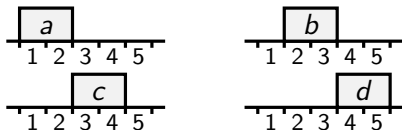
Heaps of pieces and trace monoids

Heap of pieces

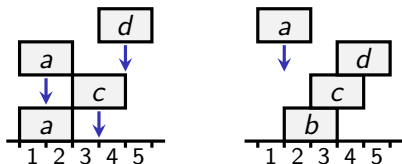
- Pieces:



- Horizontal layout:



- Vertical heaps:



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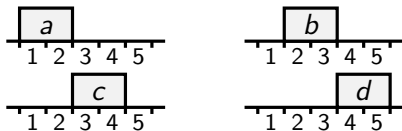
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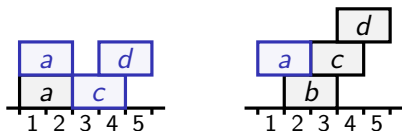
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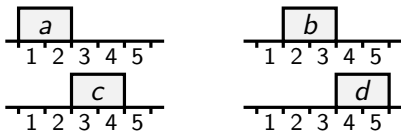
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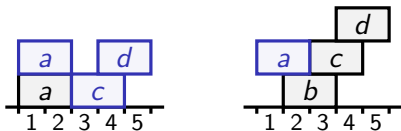
- Pieces:



- Horizontal layout:



- Vertical heaps:



Trace monoid

- Alphabet:

$$\Sigma = \{a, b, c, d\}$$

- Dependence relation:

$$\mathbf{D} = \{\{a, b\}, \{b, c\}, \{c, d\}\}$$

- Trace monoid:

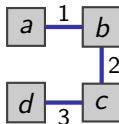
$$\mathcal{M} = \langle a, b, c, d \mid ac=ca, ad=da, bd=db \rangle^+$$



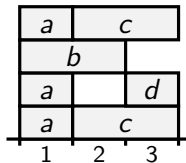
Dependency graph

Heaps of pieces and dependency graph

Dependency graph

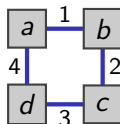
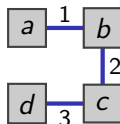


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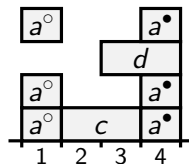
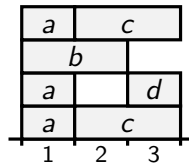


Heaps of pieces and dependency graph

Dependency graph

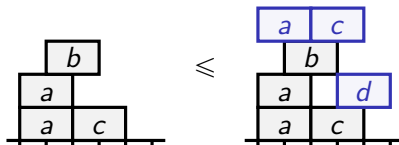


Heap of (disconnected) pieces



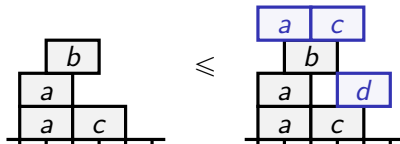
Heaps of pieces and left divisibility

Heap of pieces



Heaps of pieces and left divisibility

Heap of pieces

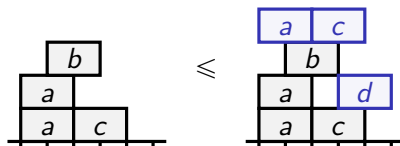


Can you pick one **infinite** heap **uniformly at random**?

- 1 Define your preferred notion of heap length
- 2 Study uniform distributions on traces of length k : what if $k \rightarrow \infty$?

Heaps of pieces and left divisibility

Heap of pieces



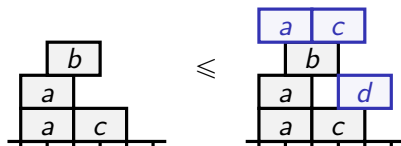
Can you pick one **infinite** heap **uniformly at random**?

- ① Heap length = #pieces in the heap
- ② **Weak convergence** of distributions:

$$\mu_k \rightarrow \nu \Leftrightarrow (\forall x \in \mathcal{M}^+, \mathbb{P}_{\mu_k}[x \leq \xi] \rightarrow \mathbb{P}_{\nu}[x \leq \xi])$$

Heaps of pieces and left divisibility

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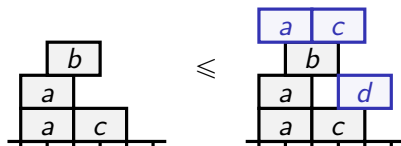
$$\mu_k \rightarrow \nu \Leftrightarrow (\forall x \in \mathcal{M}^+, \mathbb{P}_{\mu_k}[x \leq \xi] \rightarrow \mathbb{P}_{\nu}[x \leq \xi])$$

Theorem (Abbes & Mairesse 2015)

If μ_k is the uniform measure on $\mathcal{M}_k = \{\zeta \in \mathcal{M} : |\zeta| = k\}$,
 ν **exists** and is the **critical Bernoulli** distribution of \mathcal{M} .

Heaps of pieces and left divisibility

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Theorem (Abbes & Mairesse 2015) — not constructive!

If μ_k is the uniform measure on $\mathcal{M}_k = \{\zeta \in \mathcal{M} : |\zeta| = k\}$,
 ν **exists** and is the **critical Bernoulli** distribution of \mathcal{M} .

Bernoulli distributions

Definition

A probability measure μ on \mathcal{M} is:

- **Bernoulli** if $\forall x \in \mathcal{M}, \forall \sigma \in \Sigma, \mathbb{P}_\mu[x \sigma \leq \zeta \mid x \leq \zeta] = \mathbb{P}_\mu[\sigma \leq \zeta]$
- **uniform Bernoulli** of parameter p if, furthermore, $\mathbb{P}_\mu[\sigma \leq \zeta] = p$

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$$\text{Uniform Bernoulli} \Rightarrow \mathbb{P}_\mu[x \leq \zeta] = p^{|x|} \Rightarrow \mathbb{P}_\mu[x = \zeta] \propto p^{|x|}.$$

Bernoulli distributions

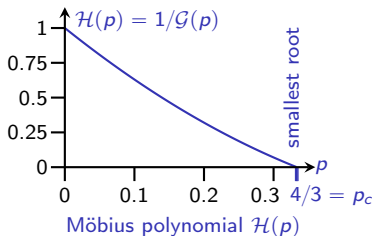
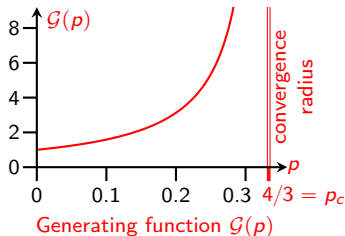
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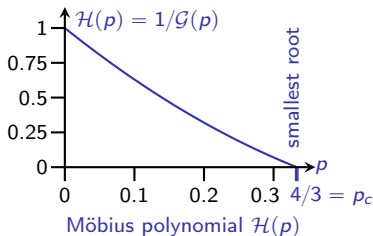
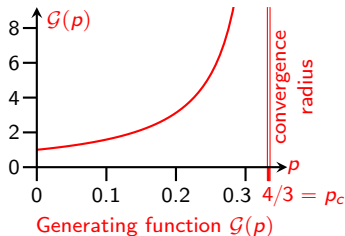
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- **critical Bernoulli** if $p = p_c$ (requires **infinite** heaps)

Which are the **possible values** of the parameter p ?

$$0 \leq p \leq p_c$$

$$\text{Uniform Bernoulli} \Rightarrow \mathbb{P}_\mu[x \leq \zeta] = p^{|x|} \Rightarrow \mathbb{P}_\mu[x = \zeta] \propto p^{|x|} \stackrel{?}{=} p^{|x|} \mathcal{H}(p).$$



Bernoulli distributions

Definition

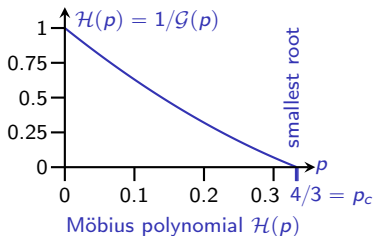
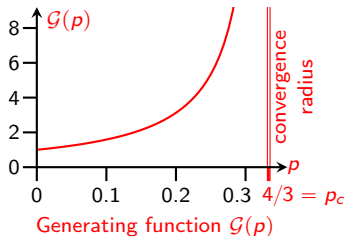
A probability measure μ on $\overline{\mathcal{M}}$ is:

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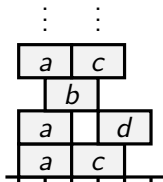
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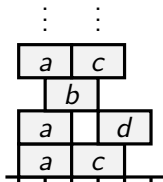
Infinite heaps

Heap of pieces

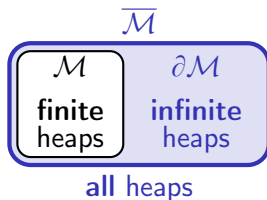


Infinite heaps

Heap of pieces



Sets of interest

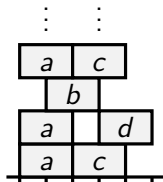


$\overline{\mathcal{M}}$ is the **compactification** of \mathcal{M} for the topology induced by $\uparrow \mathbf{x} = \{\zeta : \mathbf{x} \leq \zeta\}$

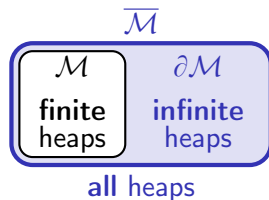
Reminder: $\mu_k \rightarrow \nu \Leftrightarrow (\forall x \in \mathcal{M}^+, \mathbb{P}_{\mu_k}[x \leq \xi] \rightarrow \mathbb{P}_{\nu}[x \leq \xi])$

Infinite heaps

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Fact #1 (Abbes & Mairesse 2015)

The limit ν is a distribution **on the set** $\overline{\mathcal{M}}$ with **support** $\partial\mathcal{M}$. ⚠

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Simulating the limit ν

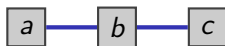
Idea #1: Pick $\zeta_k \sim \mu_k$, pick a piece x **wisely** and set $\zeta_{k+1} = \zeta_k \cdot x$

Simulating the limit ν

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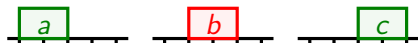
Problem: In general, ζ_{k+1} **cannot** be distributed according to μ_{k+1}

Example in the monoid $\langle a, b, c \mid ac = ca \rangle^+$



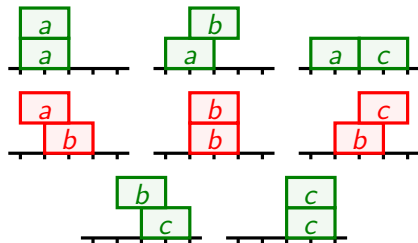
Dependency graph

$$\mu_1(\uparrow a \cup \uparrow c) = 2/3$$



>

$$\mu_2(\uparrow a \cup \uparrow c) = 5/8$$



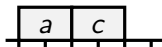
Simulating the limit ν

Idea #2: Simulate $\zeta \sim \nu$, **floor by floor**



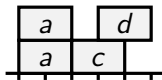
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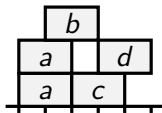
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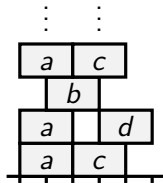
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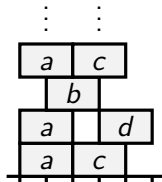
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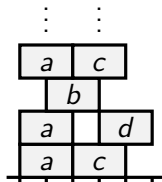


Fact #2 (Abbes & Mairesse 2015)

This approach works because ν is **Bernoulli**!

Simulating the limit ν

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Fact #2 (Abbes & Mairesse 2015)

This approach works because ν is **Bernoulli**!

Problems: Huge state space (exponential number of possible floors)
Adding one floor implies synchronising many pieces

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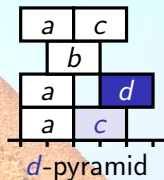
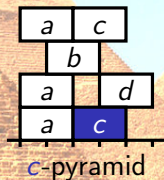
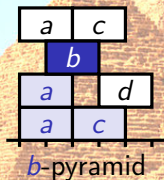
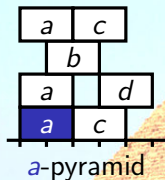
Simulating the limit ν piece by piece[⚠]

Idea #3: Decompose heaps recursively by using **pyramids**



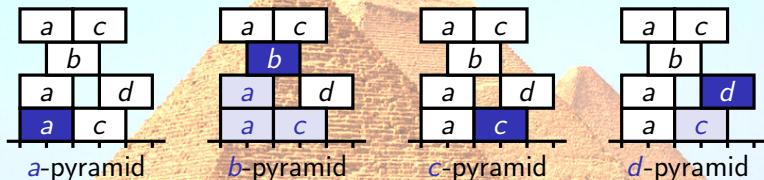
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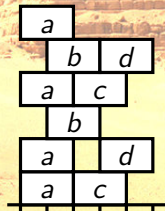


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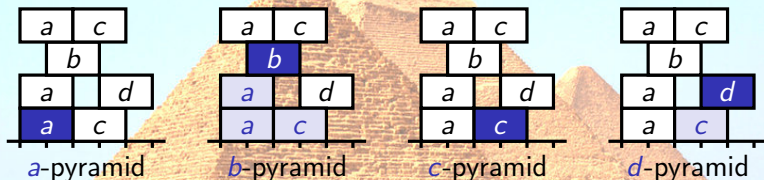


Example: Recursive decomposition using *b*-pyramids

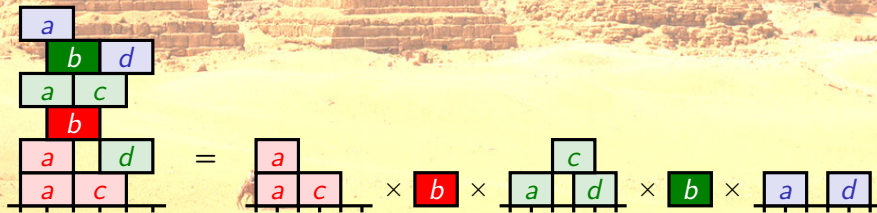


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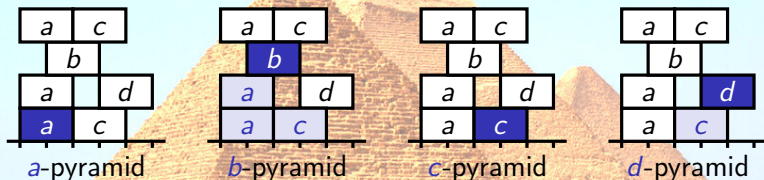


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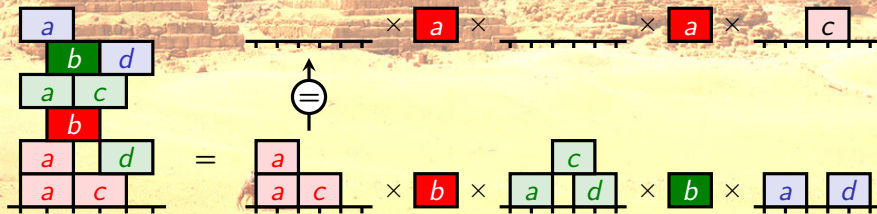


Simulating the limit ν piece by piece[⚠]

Idea #3: Decompose heaps recursively by using **pyramids**



Example: Recursive decomposition using *b*-pyramids, then *a*-pyramids...



Simulating the limit ν piece by piece[!]

Idea #3: Decompose heaps recursively by using **independent** pyramids

Theorem – continued (Abbes & J. 2017)

If $\mathbf{a} \in \Sigma$ and ν is **Bernoulli** on $\overline{\mathcal{M}}(\Sigma)$, then



The diagram shows the equation: **Heap** = **a-pyramid** \times **Heap** + **a-free heap**. The visual elements are: a battery (Heap), an equals sign, a pyramid (a-pyramid), a multiplication sign, another battery (Heap), a plus sign, and a third battery (a-free heap). Below each image is its label: Heap, a-pyramid, Heap, and a-free heap.

where right-hand side random variables are **independent**.

Simulating the limit ν piece by piece[⚠]

Idea #3: Decompose heaps recursively by using **independent** pyramids

Theorem – continued (Abbes & J. 2017)

If $\mathbf{a} \in \Sigma$ and ν is **Bernoulli** on $\overline{\mathcal{M}}(\Sigma)$, then



The diagram illustrates the decomposition of a heap. On the left, a single battery is labeled "Heap". This is equal to a pyramid (labeled "a-pyramid") multiplied by a battery (labeled "Heap") plus another battery (labeled "a-free heap").

$$\text{Heap} = \text{a-pyramid} \times \text{Heap} + \text{a-free heap}$$

where right-hand side random variables are **independent**.



a-pyramid



The diagram shows a battery (representing a free heap) multiplied by a red circle containing the letter 'A'.

$$\text{a-free heap } \xi \text{ s.t. } \mathbf{S}(\xi) \subseteq \mathbf{D}(\mathbf{a})$$

a-free heap ξ s.t.

$$\mathbf{S}(\xi) \subseteq \mathbf{D}(\mathbf{a})$$

- $\mathbf{S}(\xi) = \{\sigma \in \Sigma : \xi \in \mathcal{M}^+ \cdot \sigma\}$
- $\mathbf{D}(\mathbf{a}) = \{\sigma \in \Sigma : \sigma \cdot \mathbf{a} \neq \mathbf{a} \cdot \sigma\}$

Simulating the limit ν piece by piece[⚠]

Idea #3: Decompose heaps recursively by using **independent** pyramids

Theorem – continued (Abbes & J. 2017)

If $X \subseteq \Sigma$, $\mathbf{a} \in \Sigma$ and ν is **Bernoulli** on $\overline{\mathcal{M}}(\Sigma)$, then

$$\begin{array}{ccccccc}
 \text{Heap } \xi \text{ s.t.} & = & \sum_{k=1}^{\infty} & \text{a-pyramid} & \times & \text{a-free heap } \xi \text{ s.t.} & + & \text{a-free heap } \xi \text{ s.t.} \\
 \mathbf{S}(\xi) \subseteq X & & & & & \mathbf{S}(\mathbf{a} \cdot \xi) \subseteq X & & \mathbf{S}(\xi) \subseteq X
 \end{array}$$

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a-pyramid

$$= \text{a-free heap } \xi \text{ s.t. } \mathbf{S}(\xi) \subseteq \mathbf{D}(\mathbf{a}) \times \textcircled{\mathbf{A}}$$

a-free heap ξ s.t.
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- $\mathbf{S}(\xi) = \{\sigma \in \Sigma : \xi \in \mathcal{M}^+ \cdot \sigma\}$
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Our algorithm: Generating heaps distributed according to ν

Remarks:

- if ν has support $\partial\mathcal{M}$, $\xi =$ infinite sequence of **a**-pyramids $(k \leftarrow \infty)$
- **a**-pyramid = (**a**-free sequence ξ with $\mathbf{S}(\xi) \subseteq \mathbf{D}(\mathbf{a})$) $\cdot \mathbf{a}$
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Running time/piece: nq or n **Read-only memory usage:** n or 2^n

where $\mathcal{M}' = \mathcal{M}(\Sigma \setminus \{a\})$ and $q = \mathcal{G}_{\mathcal{M}'}(p_c) \leq 1/p_c^n$ ($q = n^{\Theta(n)}$ is possible)

Contents

- 1 Introduction: Heaps of pieces and trace monoids
- 2 Simulating Bernoulli distributions
- 3 Step-by-step simulation and pyramids
- 4 Conclusion**

A distributed simulation algorithm

Algorithm based on:

- precomputing and storing **Möbius polynomials** of sub-monoids
- decomposing heaps into **independent** pyramids/heaps in sub-monoids
- outputting pieces **one by one** with little synchronisation

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Very efficient on graphs with:

- few cycles (small storage/high efficiency for a mix between variants)
- small tree-width (no preprocessing/storage)

Thank
you

