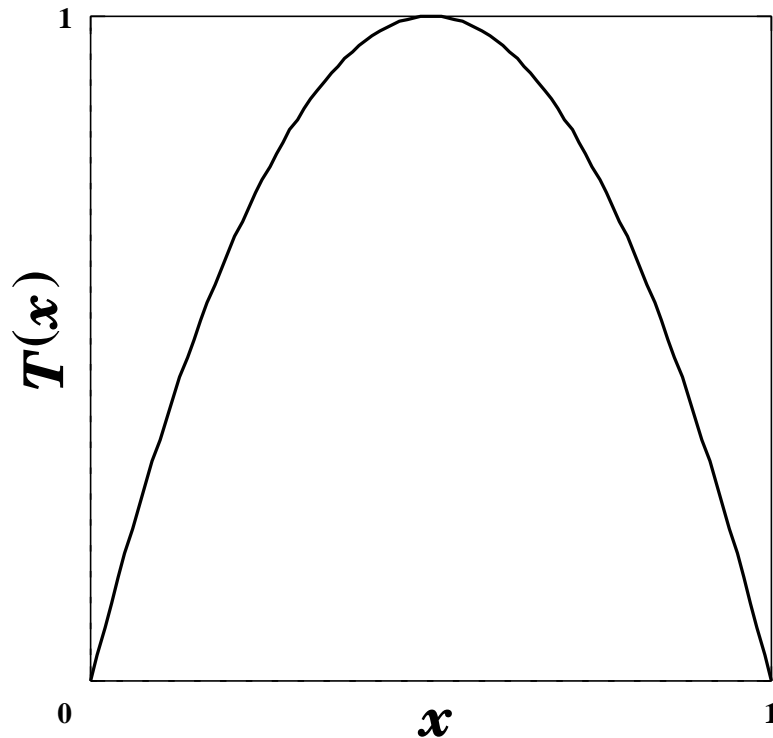


**The topological entropy and correlational  
properties of the discretized Markov  
 $\beta$ -transformations and their applications**

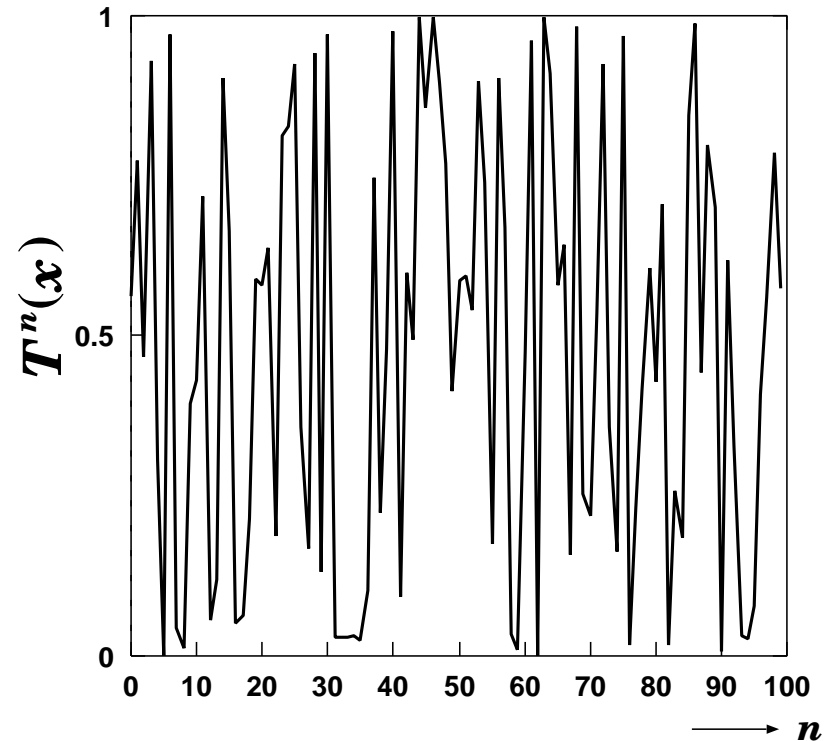
**Hiroshi FUJISAKI**  
**Kanazawa University**

**Jan. 31, 2017**

- Ulam and von Neumann (Bull. Math. Soc., 1947)  
Logistic map:  $T(x) = 4x(1 - x)$ ,  $0 \leq x \leq 1$



$\Rightarrow$



Given an initial value  $x = T^0(x)$ ,

$$T^n(x) = T(T^{n-1}(x)) \quad \text{for } n = 1, 2, \dots.$$

The sequence  $(T^n(x))_{n=0}^{\infty}$  is a good candidate for the pseudo-random numbers.

Ulam and von Neumann's idea requires handling real numbers for practice. On the contrary, computers can only deal with floating point numbers. Hence we need ergodic theory for a transformation from a finite set onto itself to understand the behaviour of the iterates of one-dimensional transformations implemented in computers.

No way is known to give a good theoretical model that tells us characteristics of the execution time for floating point numbers. ( D. Knuth, *The Art of Computer Programming*, vol. 2, 3rd ed., Addison-Wesley, '97)

# Discretized Bernoulli Transformations

## ○ Cryptosystems

- Permutation Cipher Based on Discretized Unimodal Bernoulli Transformations (N. Masuda and K. Aihara, *Trans. of IEICE*, '99 (in Japanese))

## ○ Spreading Seq.s for SSMA **Communication Systems**

- Maximal-Period Sequences Based on Discretized Bernoulli Transformations (A. Tsuneda, Y. Kuga, and T. Inoue, *IEICE Trans. on Funda.*, 2002)

A Generalization of *de Bruijn Sequences*

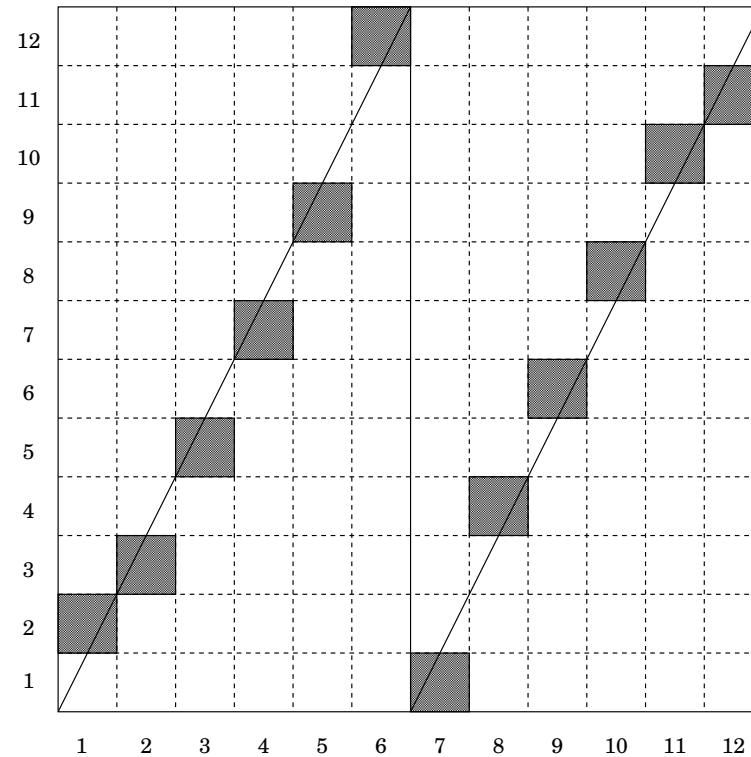
## Markov Partition

We use  $|E|$  to denote the cardinality of a set  $E$ .

**Definition 1** Let  $T : [0, 1) \rightarrow [0, 1)$ . Let  $\mathcal{P}$  be a partition of  $[0, 1)$  given by the point  $0 = a_0 < a_1 < \cdots < a_{|\mathcal{P}|} = 1$ . For  $i = 1, \dots, |\mathcal{P}|$ , let  $I_i = (a_{i-1}, a_i)$  and denote the restriction of  $T$  to  $I_i$  by  $T|_{I_i}$ . If  $T|_{I_i}$  is a homeomorphism from  $I_i$  onto the interior of some connected union of the closures of intervals of  $\mathcal{P}$ , then  $T$  is said to be Markov. The partition  $\mathcal{P} = \{I_i\}_{i=1}^{|\mathcal{P}|}$  is referred to as a *Markov partition* with respect to  $T$ .

An example of discretized dyadic transformations ( $2m = 12$ ):

$$\sigma = \begin{pmatrix} I_1 & I_2 & I_3 & I_4 & I_5 & I_6 & I_7 & I_8 & I_9 & I_{10} & I_{11} & I_{12} \\ I_2 & I_3 & I_5 & I_7 & I_9 & I_{12} & I_1 & I_4 & I_6 & I_8 & I_{10} & I_{11} \end{pmatrix}.$$



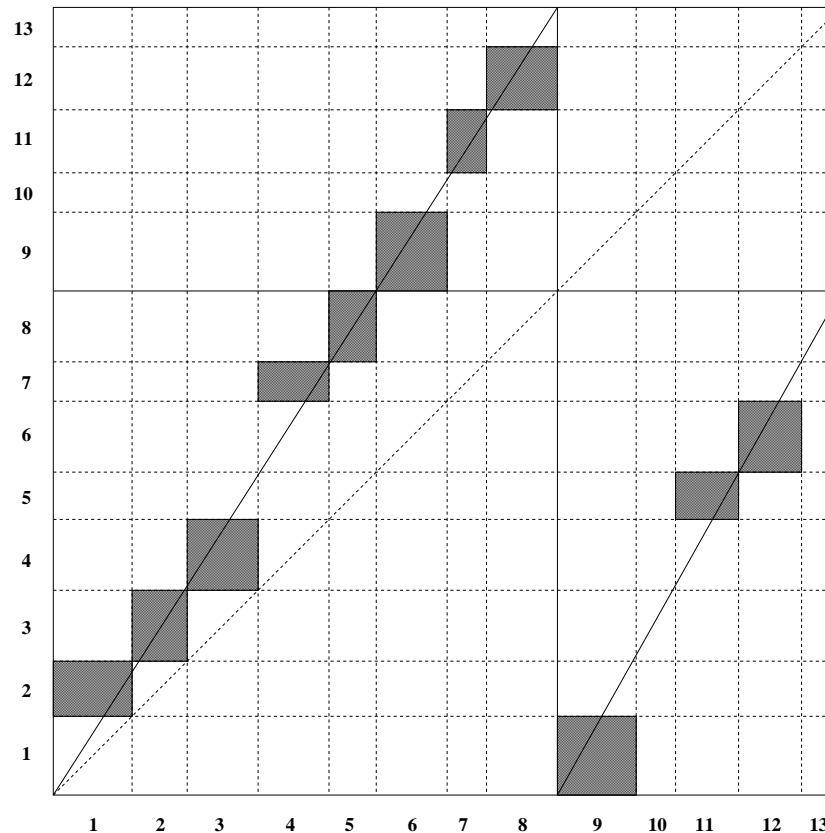
$\sigma$  determines a full-length sequence 000010111101.

If  $2m = 2^n$ , then the full-length sequence is called the [de Bruijn sequence](#).

# Discretized Golden Mean Transformations

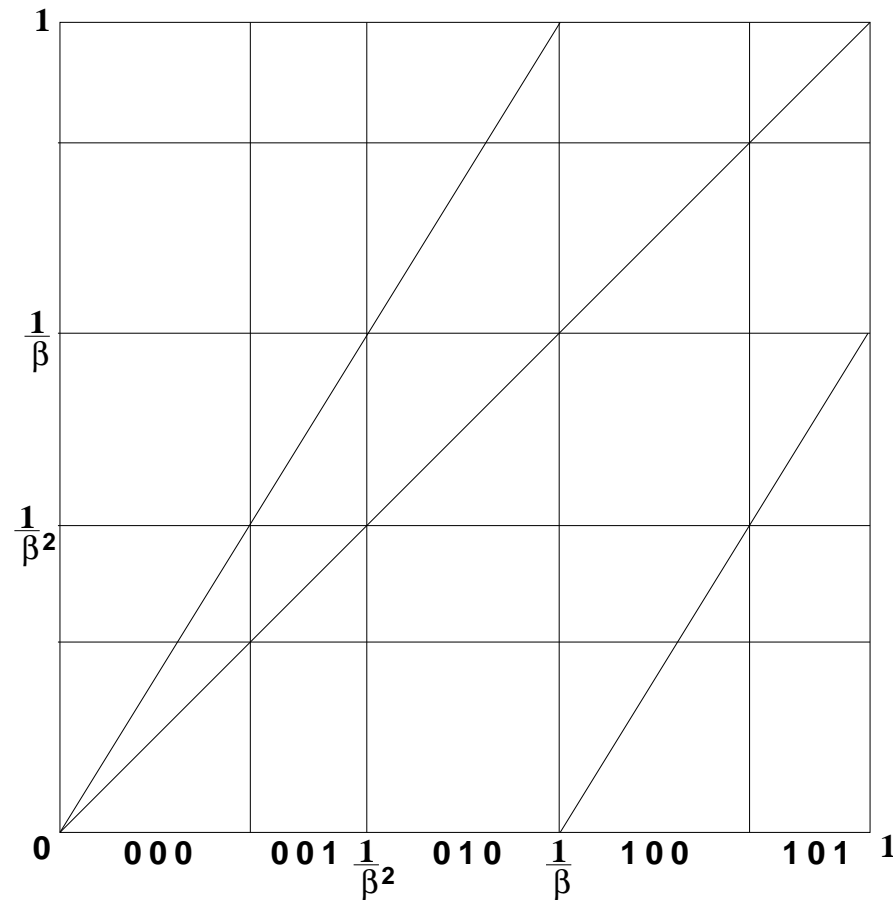
$$\sigma = \begin{pmatrix} I_1 & I_2 & I_3 & I_4 & I_5 & I_6 & I_7 & I_8 & I_9 & I_{11} & I_{12} \\ I_2 & I_3 & I_4 & I_7 & I_8 & I_9 & I_{11} & I_{12} & I_1 & I_5 & I_6 \end{pmatrix}.$$

Note that  $I_{10}$  and  $I_{13}$  are excluded from the Markov partition.



$\sigma$  determines a full-length sequence 00000100101. [F. Enomoto and S. Ito, Workshop Number Theory and Ergodic Theory, 2004)]

# Graph Representation of the Markov Transformation

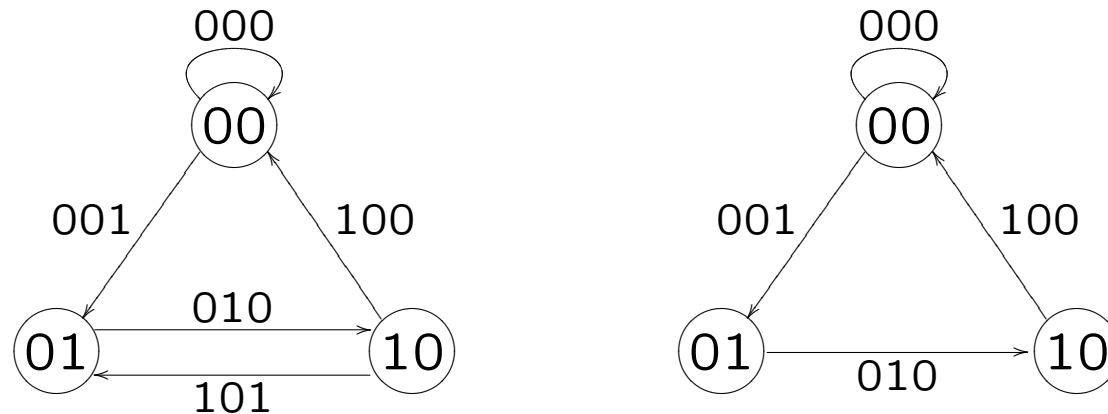


For an irreducible aperiodic Markov transformation  $T$ , given a [Markov partition](#)  $\mathcal{P}$  with respect to  $T$ , corresponding each subinterval  $I \in \mathcal{P}$  to one [edge](#)  $e(I)$ , we obtain the set  $\mathcal{A}$  of edges.



## Eulerian Subgraph Spanning $G$

A directed graph  $H = (\mathcal{W}, \mathcal{B})$  is said to be a subgraph of the directed graph  $G = (\mathcal{V}, \mathcal{A})$  if  $\mathcal{W} \subset \mathcal{V}$  and  $\mathcal{B} \subset \mathcal{A}$ . In this case we write  $H \subset G$ . The directed graph  $H$  is called a spanning subgraph of  $G$  if  $\mathcal{W} = \mathcal{V}$ . Furthermore, if  $H$  is Eulerian, it is called Eulerian subgraph spanning  $G$ . We are interested in the spanning Eulerian subgraph of  $G$  with maximal number of edges.



The spanning **Eulerian** subgraph with **maximal** number of edges

Full-length sequences based on the discretized Markov transformation are exactly Eulerian circuits in  $H$ , whose length is given by  $|\mathcal{B}|$ .

## Preliminaries

Let  $\Sigma$  be a finite alphabet. The full  $\Sigma$ -shift is denoted by

$$\Sigma^{\mathbb{Z}} = \{x = (x_i)_{i \in \mathbb{Z}} : \forall i \in \mathbb{Z}, x_i \in \Sigma\}$$

which is endowed with the product topology arising from the discrete topology on  $\Sigma$ . The shift transformation  $\sigma : \Sigma^{\mathbb{Z}} \rightarrow \Sigma^{\mathbb{Z}}$  is defined by

$$\sigma((x_i)_{i \in \mathbb{Z}}) = (x_{i+1})_{i \in \mathbb{Z}}.$$

The closed shift-invariant subsets of  $\Sigma^{\mathbb{Z}}$  are called subshifts. For a subshift  $X$ , we use  $\sigma_X$  to denote the shift transformation on  $X$ , which is the restriction to  $X$  of  $\sigma$  on  $\Sigma^{\mathbb{Z}}$ . For simplicity, we shall write  $\sigma : X \rightarrow X$  rather than  $\sigma_X$ .

We call elements  $u = u_1u_2\cdots u_n \in \Sigma^n$  blocks over  $\Sigma$  of length  $n$  ( $n \geq 1$ ). We use  $\Sigma^*$  to denote the collection of all blocks over  $\Sigma$  and the empty block  $\epsilon$ . For a subshift  $X$ , we use  $\mathcal{L}_n(X)$  to denote the collection of all  $n$ -blocks appearing in points in  $X$ . The language of  $X$  is the collection  $\mathcal{L}(X) = \bigcup_{n=0}^{\infty} \mathcal{L}_n(X)$ , where  $\mathcal{L}_0(X) = \{\epsilon\}$ .

**Definition 2** *The topological entropy of a subshift  $X$  is defined by*

$$h(X) = \lim_{n \rightarrow \infty} \frac{1}{n} \log |\mathcal{L}_n(X)|.$$

We use  $|E|$  to denote the cardinality of a set  $E$ .

## Higher Edge Graph

**Definition 3** Let  $G$  be a graph. For  $n \geq 2$  we define the  $n$ th higher edge graph  $G^{[n]}$  of  $G$  to have vertex set  $\mathcal{L}_{n-1}(X_{A_G})$  and to have edge set containing exactly one edge from  $e_1e_2 \cdots e_{n-1}$  to  $f_1f_2 \cdots f_{n-1}$  whenever  $e_2e_3 \cdots e_{n-1} = f_1f_2 \cdots f_{n-2}$  (or  $t(e_1) = i(f_1)$  if  $n = 2$ ), and none otherwise. The edge is named

$$e_1e_2e_3 \cdots e_{n-1}f_{n-1} = e_1f_1f_2 \cdots f_{n-1}.$$

For  $n = 1$  we set  $G^{[1]} = G$ .

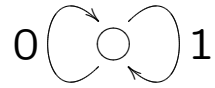
# Discretized Dyadic Transformations

Let  $T : [0, 1] \rightarrow [0, 1]$  be the dyadic transformation:

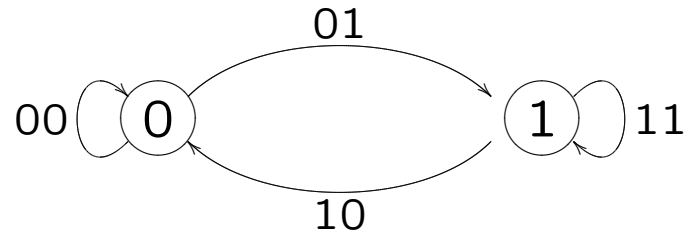
$$T(x) = 2x \pmod{1}, \quad x \in [0, 1].$$

If we take a Markov partition of  $[0, 1]$  given by the point  $0 < 1/2 < 1$ , then we obtain the graph  $G$  representing the dyadic transformation.

$$0 \left( \begin{array}{c} \circlearrowleft \\ \circlearrowright \end{array} \right) 1 \quad G = G^{[1]} = H_1.$$



$$G = G^{[1]} = H_1.$$



$$G^{[2]} = H_2.$$

For each  $n (\geq 1)$ , we obtain  $G^{[n]} = (\{0, 1\}^{n-1}, \{0, 1\}^n)$ . Since  $G^{[n]}$  is Eulerian, we have  $H_n = G^{[n]}$  for each  $n$ . The Eulerian circuits in  $G^{[n]}$  are called the **de Bruijn sequences** of length  $2^n$  because of the following theorem. For the same reason,  $G^{[n]}$  is called the **de Bruijn graph**.

**Theorem 1 (de Bruijn, 1946, Flye Sainte-Marie, 1894)** *For each positive integer  $n$ , there are exactly  $2^{2^{n-1}-n}$  Eulerian circuits in  $G^{[n]}$ .*

# The Topological Entropy of Discretized Markov Transformation

Let  $G$  be the graph representing the Markov transformation. Then we obtain a sequence  $(G^{[n]})_{n=1}^{\infty}$  of higher edge graphs of  $G$ . For each  $n \geq 1$ , we use  $H_n = (\mathcal{L}_{n-1}(X_{A_G}), \mathcal{B}_n)$  to denote the Eulerian subgraph spanning  $G^{[n]}$  with maximal number of edges, each of which leads to a discretized Markov transformation  $\hat{T}_n$ .

We use  $\nu_n$  to denote the number of the full-length sequence in  $H_n$ . Recall that the length is given by  $|\mathcal{B}_n|$ .

**Definition 4** The topological entropy of the discretized Markov transformation  $\mathcal{T} = (\hat{T}_n)_{n=1}^{\infty}$  of  $T$  is defined by

$$h_{\mathcal{T}} = \lim_{n \rightarrow \infty} \frac{1}{|\mathcal{B}_n|} \log \nu_n.$$



**Example 1** *The topological entropy of the discretized dyadic transformation  $\mathcal{T}$  is given by*

$$h_{\mathcal{T}} = \frac{1}{2} \log 2.$$

**Remark 1** *Since it is also shown in [de Bruijn, 1946] and [Flye Sainte-Marie, 1894] that, for each  $n (\geq 1)$ , there are exactly*

$$\{(k-1)!\}^{k^{n-1}} k^{k^{n-1}-n}$$

*Eulerian circuits of length  $k^n$  in*

$$G^{[n]} = (\{0, 1, \dots, k-1\}^{n-1}, \{0, 1, \dots, k-1\}^n),$$

*the topological entropy of the discretized  $k$ -adic transformation is given by*

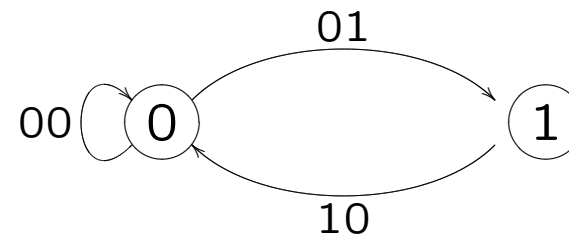
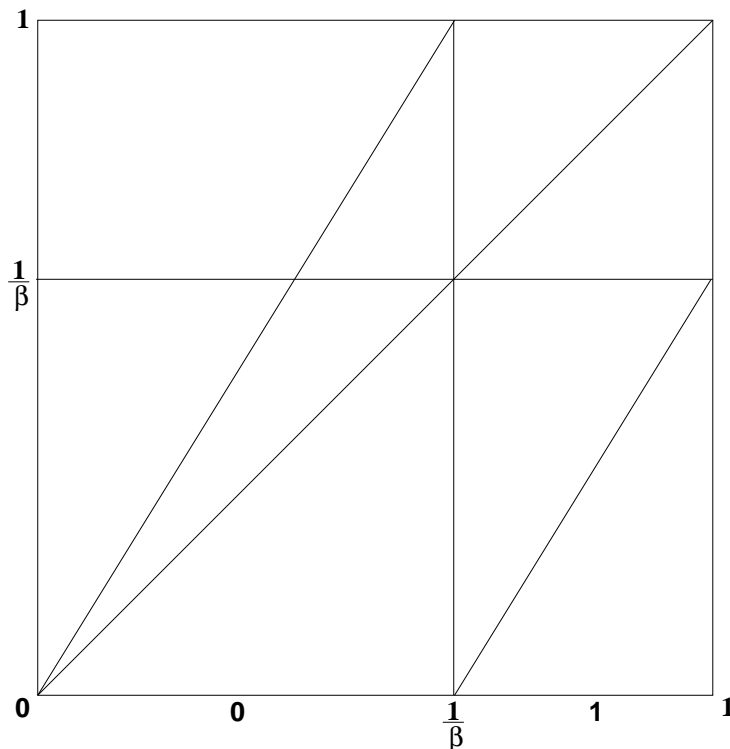
$$\frac{1}{k} \log(k!).$$

# Discretized Golden Mean Transformation

Let  $T : [0, 1] \rightarrow [0, 1]$  be the golden mean transformation:

$$T(x) = \beta x \pmod{1}, \quad x \in [0, 1],$$

where  $\beta$  is the golden mean number  $\frac{1+\sqrt{5}}{2}$ .



$$G^{[2]} = H_2.$$

In view of  $G^{[2]}$ , the set of forbidden blocks is given by  $\mathcal{F} = \{11\}$ .

For each  $n$  ( $\geq 2$ ), we obtain  $G^{[n]} = (\mathcal{L}_{n-1}(X_{\mathcal{F}}), \mathcal{L}_n(X_{\mathcal{F}}))$  and the **Eulerian subgraph**  $H_n = (\mathcal{L}_{n-1}(X_{\mathcal{F}}), \mathcal{B}_n)$  **spanning**  $G^{[n]}$  with **maximal** number of edges. Although  $G^{[2]}$  is Eulerian, which implies  $H_2 = G^{[2]}$ ,  $G^{[n]}$  is not always Eulerian for  $n$  ( $\geq 3$ ). In fact,  $H_3$  is a proper subgraph of  $G^{[3]}$ , in symbols  $H_3 \subsetneq G^{[3]}$ . We observed that  $H_n \subsetneq G^{[n]}$  for any  $n$  ( $\geq 3$ ).

Noting that the sequence  $(|\mathcal{B}_n|)_{n=2}^{\infty}$  is the Fibonacci numbers defined by the recurrence relation  $|\mathcal{B}_n| = |\mathcal{B}_{n-1}| + |\mathcal{B}_{n-2}|$  ( $\geq 4$ ) with  $|\mathcal{B}_2| = 3$  and  $|\mathcal{B}_3| = 4$ , we obtain

$$|\mathcal{B}_n| = \beta^n + \bar{\beta}^n \quad \text{for } n \geq 2, \quad (1)$$

where  $\bar{\beta} = \frac{1-\sqrt{5}}{2}$ .

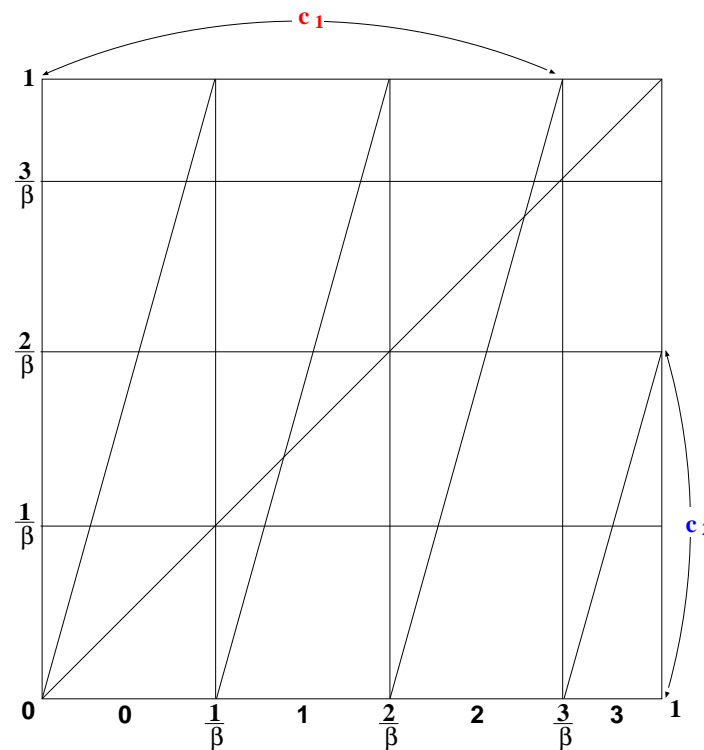
The topological entropy of the discretized golden mean transformation is given by

## Theorem 2

$$\lim_{n \rightarrow \infty} \frac{1}{|\mathcal{B}_n|} \log \nu_n = \frac{1}{\beta(\beta - \bar{\beta})} \log 2.$$

# A Class of Markov Transformations Associated with Greedy $\beta$ -Expansion

Now we are in the position to consider the discretized Markov  $\beta$ -transformations with the alphabet  $\Sigma = \{0, 1, \dots, k-1\}$  ( $k \geq 2$ ) and the set  $\mathcal{F} = \{(k-1)c, \dots, (k-1)(k-1)\}$  ( $1 \leq c \leq k-1$ ) of  $(k-c)$  forbidden blocks. Setting  $c_1 = k-1$  and  $c = c_2$ ,  $\beta$  is the positive solution of  $t^2 - c_1 t - c_2 = 0$ .



Generally we have

### Theorem 3

$$\lim_{n \rightarrow \infty} \frac{1}{|\mathcal{B}_n|} \log \nu_n = \frac{c_2}{\beta(\beta - \bar{\beta})} \log(k!) + \frac{\beta - c_2}{\beta(\beta - \bar{\beta})} \log\{(k-1)!\} + \frac{1}{\beta(\beta - \bar{\beta})} \log(c_2!).$$

## Variational Principle

**Definition 5** If  $\lambda$  is the *Perron-Frobenius eigenvalue* of  $A$  and  $(u_0, \dots, u_{|\Sigma|-1})$  is a strictly positive left eigenvector and  $(v_0, \dots, v_{|\Sigma|-1})$  is a strictly positive right eigenvector with  $\sum_{i=0}^{|\Sigma|-1} u_i v_i = 1$ , then we obtain a Markov measure given by a probability vector  $p = (p_0, \dots, p_{|\Sigma|-1})$  and a stochastic matrix  $P = (p_{i,j})$  where

$$p_i = u_i v_i \quad \text{and} \quad p_{i,j} = \frac{a_{i,j} v_j}{\lambda v_i}.$$

We call this measure the *Parry measure* for  $\sigma : X_A \rightarrow X_A$ .

We obtain

$$-\sum_{i,j} p_i p_{i,j} \log p_{i,j} = \log \lambda.$$

The left hand side is the *Shannon entropy* for a Markov chain given by  $(p, P)$  while the right hand side is the *topological entropy* of  $\sigma : X_A \rightarrow X_A$ .

## Correlational Properties of the de Bruijn Sequences

**Definition 6** *The cross-correlation function of time delay  $\ell$  for the sequences  $\mathbf{X} = (X_i)_{i=0}^{N-1}$  and  $\mathbf{Y} = (Y_i)_{i=0}^{N-1}$  over  $\Sigma = \{0, 1, \dots, k-1\}$  ( $k \geq 2$ ) is defined by*

$$R_N(\ell; \mathbf{X}, \mathbf{Y}) = \sum_{i=0}^{N-1} \exp\left(\frac{X_i}{k} 2\pi\sqrt{-1}\right) \exp\left(-\frac{Y_{i+\ell \pmod{N}}}{k} 2\pi\sqrt{-1}\right),$$

where  $\ell = 0, 1, \dots, N-1$  and, for integers  $a$  and  $b$  ( $\geq 1$ ),  $a \pmod{b}$  denotes the least residue of  $a$  to modulus  $b$ . The normalized cross-correlation function of time delay  $\ell$  for the sequences  $\mathbf{X}$  and  $\mathbf{Y}$  is defined by

$$r_N(\ell; \mathbf{X}, \mathbf{Y}) = \frac{1}{N} R_N(\ell; \mathbf{X}, \mathbf{Y}).$$

If  $\mathbf{X} = \mathbf{Y}$ , we call  $R_N(\ell; \mathbf{X}, \mathbf{X})$  and  $r_N(\ell; \mathbf{X}, \mathbf{X})$  the auto-correlation function and the *normalized auto-correlation function*, and simply denote them by  $R_N(\ell; \mathbf{X})$  and  $r_N(\ell; \mathbf{X})$ , respectively.

By the definition, we immediately see the following.

**Remark 2** *For any  $X$ , we have*

$$r_N(0; X) = 1.$$



The following basic properties of the normalized auto-correlation functions for the de Bruijn sequences are well known [Zhang & Chen, 1989].

**Theorem 4** *Let  $\mathbf{X}$  and  $\mathbf{Y}$  be the de Bruijn sequences of length  $N = 2^n$  ( $n \geq 1$ ). Then we have*

$$i) \quad \sum_{\ell=0}^{N-1} r_N(\ell; \mathbf{X}, \mathbf{Y}) = 0;$$

$$ii) \quad r_N(\ell; \mathbf{X}) = 0 \quad \text{for } 1 \leq \ell \leq n - 1.$$

*(Zero Correlation Zone (ZCZ))*

**Observation 1** *If  $(Z_n)_{n=0}^{\infty}$  is a sequence of independent and identically distributed (i.i.d.) random variables over  $\{1, -1\}$  with uniform distributions, Theorem 4 ii) implies*

$$r_N(\ell; \mathbf{X}) = \mathbb{E}[Z_0 Z_\ell] \quad \text{for } 0 \leq \ell \leq n - 1.$$

# Correlational Properties of the the Full-Length Sequences Based on the Discretized Golden Mean Transformation

In virtue of symbolic analysis of  $\mathcal{L}_n(X_{\mathcal{F}})$  and  $\mathcal{B}_n$ , we obtain

**Theorem 5** *Let  $X$  and  $Y$  be full-length sequences based on the discretized golden mean transformation of length  $|\mathcal{B}_n|$ . Then we obtain*

$$\sum_{\ell=0}^{|\mathcal{B}_n|-1} r_{|\mathcal{B}_n|}(\ell; X, Y) = \frac{(\beta^n - \bar{\beta}^n)^2}{(\beta - \bar{\beta})^2(\beta^n + \bar{\beta}^n)}.$$

Asymptotically, we obtain

**Remark 3**

$$\lim_{n \rightarrow \infty} \frac{1}{|\mathcal{B}_n|} \sum_{\ell=0}^{|\mathcal{B}_n|-1} r_{|\mathcal{B}_n|}(\ell; X, Y) = \frac{1}{(\beta - \bar{\beta})^2}.$$

Moreover, we obtain

**Theorem 6** *Let  $\mathbf{X}$  be a full-length sequence based on the discretized golden mean transformation of length  $|\mathcal{B}_n|$ . Then for  $1 \leq \ell \leq n - 1$ , we obtain*

$$r_{|\mathcal{B}_n|}(\ell; \mathbf{X}) = \frac{1}{(\beta - \bar{\beta})^2} \left\{ 1 + 4 \left( \frac{\bar{\beta}}{\beta} \right)^\ell \cdot \frac{1 + \left( \frac{\bar{\beta}}{\beta} \right)^{n-2\ell}}{1 + \left( \frac{\bar{\beta}}{\beta} \right)^n} \right\}.$$

On the other hand, for the stationary Markov process  $(Z_n)_{n=0}^{\infty}$  over  $\{1, -1\}$  with the transition matrix  $\begin{pmatrix} \frac{1}{\beta} & \frac{1}{\beta^2} \\ 1 & 0 \end{pmatrix}$ , we obtain

$$\mathbb{E}[Z_0 Z_\ell] = \frac{1}{(\beta - \bar{\beta})^2} \left\{ 1 + 4 \left( \frac{\bar{\beta}}{\beta} \right)^\ell \right\} \quad \text{for } \ell \geq 0. \quad (2)$$

For a random variable  $X$ , we use  $\mathbb{E}[X]$  to denote the expected value of  $X$ .

## Discussions

Now let us estimate the error of the normalized auto-correlation function, which is originated from the discretization of the underlying transformations. In view of Theorem 6, (2) leads to

### Observation 2

$$r_{|\mathcal{B}_n|}(\ell; \mathbf{X}) - \mathbb{E}[Z_0 Z_\ell] = \left\{ \left( \frac{\beta}{\bar{\beta}} \right)^\ell - \left( \frac{\bar{\beta}}{\beta} \right)^\ell \right\} \cdot \frac{\left( \frac{\bar{\beta}}{\beta} \right)^n}{1 + \left( \frac{\bar{\beta}}{\beta} \right)^n} \quad (3)$$

*and*

$$\lim_{n \rightarrow \infty} r_{|\mathcal{B}_n|}(\ell; \mathbf{X}) = \mathbb{E}[Z_0 Z_\ell].$$

The equation (3) implies

$$r_{|\mathcal{B}_n|}(\ell; \mathbf{X}) = \mathbb{E}[Z_0 Z_\ell] + O\left(\left(\frac{\bar{\beta}}{\beta}\right)^n\right), \quad (4)$$

where  $O$  is the big  $O$  notation from the Landau symbol.

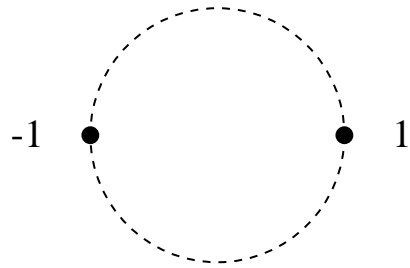
The error  $O\left(\left(\frac{\bar{\beta}}{\beta}\right)^n\right)$  can be regarded as coming from the discretization of the underlying  $\beta$ -transformation.

It is noteworthy that (4) holds even for the de Bruijn sequences in the following sense. If the underlying transformation is the dyadic transformation, we have  $\beta = 2$  and  $\bar{\beta} = 0$ . Thus we obtain  $O\left(\left(\frac{\bar{\beta}}{\beta}\right)^{2^n}\right) = 0$  for the de Bruijn sequences.

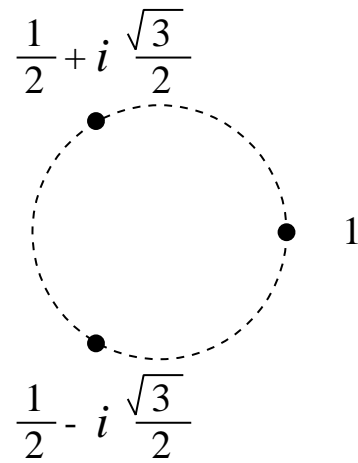
In view of Theorem 4 ii) together with this fact, (4) holds for the de Bruijn sequences if  $(Z_n)_{n=0}^\infty$  is a sequence of independent and identically distributed (i.i.d.) random variables over  $\{1, -1\}$  with uniform distributions.



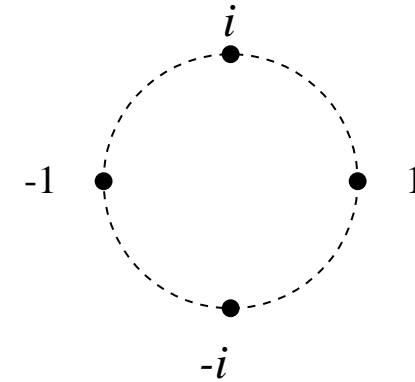
# $k$ -Phase Signals



Binary



3-Phase



Q-Phase

$$\mathbb{E}[Z_0 Z_\ell] = \frac{1}{(k-1)^2(\beta - \bar{\beta})^2} \left( \left\{ \beta + (k-1)\bar{\beta} \right\} + (k-1)k^2 \left( \frac{\bar{\beta}}{\beta} \right)^\ell \right)$$

for  $\ell \geq 0$ .



By using exactly the same manner as above, we generally obtain

### Theorem 7

$$\sum_{\ell=0}^{|\mathcal{B}_n|-1} r_{|\mathcal{B}_n|}(\ell; \mathbf{X}, \mathbf{Y}) = \frac{(\beta^n - \bar{\beta}^n)^2}{(k-1)^2(\beta - \bar{\beta})^2(\beta^n + \bar{\beta}^n)}.$$

Asymptotically, we obtain

### Remark 4

$$\lim_{n \rightarrow \infty} \sum_{\ell=0}^{|\mathcal{B}_n|-1} \frac{1}{|\mathcal{B}_n|} r_{|\mathcal{B}_n|}(\ell; \mathbf{X}, \mathbf{Y}) = \frac{1}{(k-1)^2(\beta - \bar{\beta})^2}.$$

Moreover, we obtain

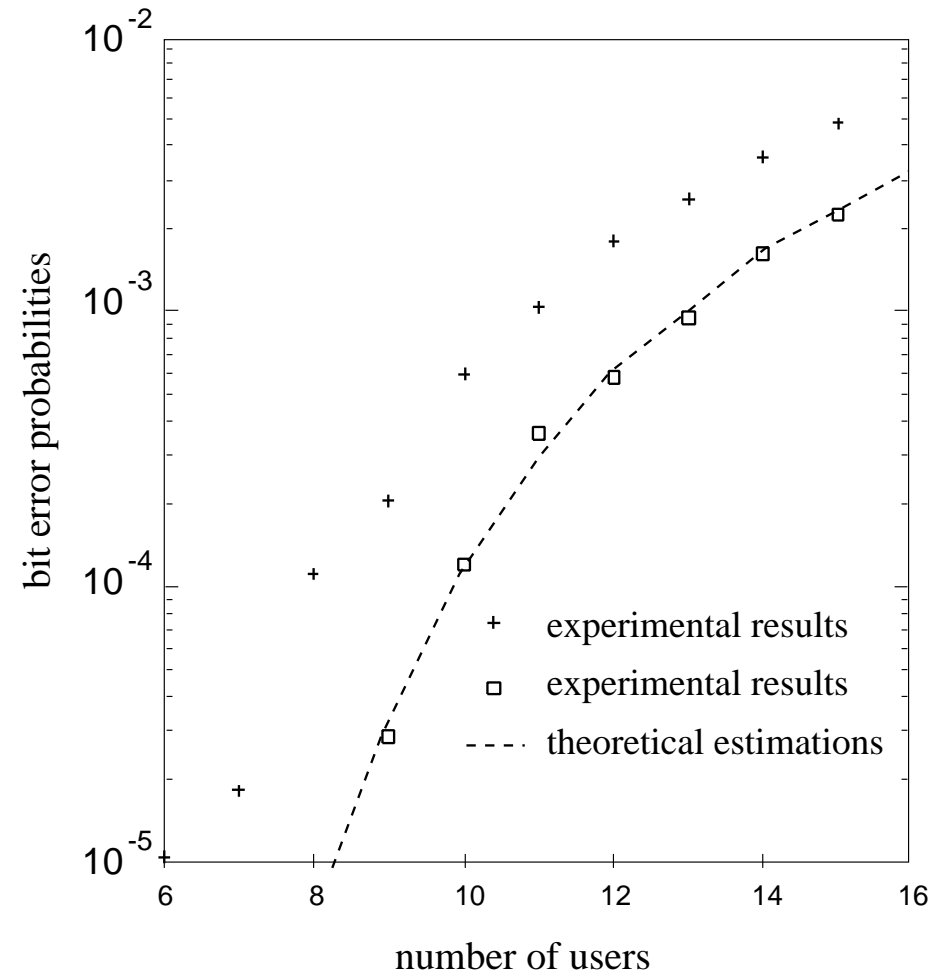
**Theorem 8** *Let  $\mathbf{X}$  be a full-length sequence based on the discretized Markov  $\beta$ -transformation of length  $|\mathcal{B}_n|$  with with the alphabet  $\Sigma = \{0, 1, \dots, k-1\}$  and the set  $\mathcal{F} = \{(k-1)(k-1)\}$  of forbidden blocks. Then for  $1 \leq \ell \leq n-1$ , we obtain*

$$r_{|\mathcal{B}_n|}(\ell; \mathbf{X}) = \frac{1}{(k-1)^2(\beta - \bar{\beta})^2} \left( \left\{ \beta + (k-1)\bar{\beta} \right\} + (k-1)k^2 \left( \frac{\bar{\beta}}{\beta} \right)^\ell \cdot \frac{1 + \left( \frac{\bar{\beta}}{\beta} \right)^{n-2\ell}}{1 + \left( \frac{\bar{\beta}}{\beta} \right)^n} \right).$$

This implies

$$r_{|\mathcal{B}_n|}(\ell; \mathbf{X}) = \mathbb{E}[Z_0 Z_\ell] + o\left(\left(\frac{\bar{\beta}}{\beta}\right)^n\right).$$

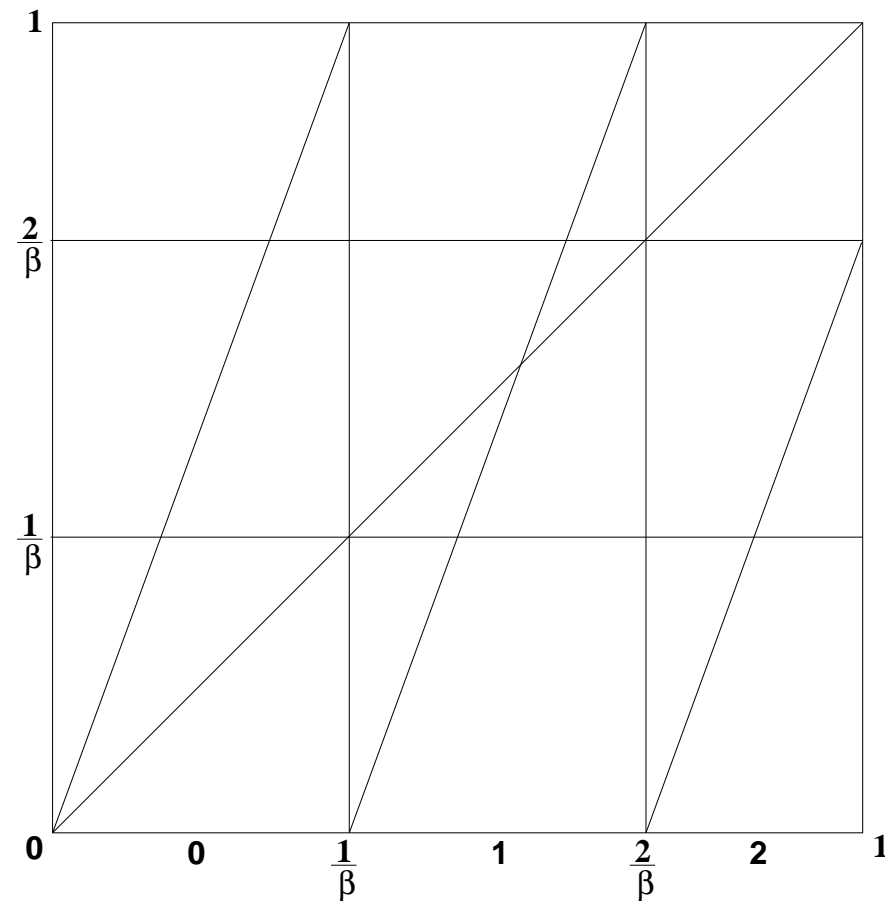
# Optimum Binary Spreading Sequences of Markov Chains



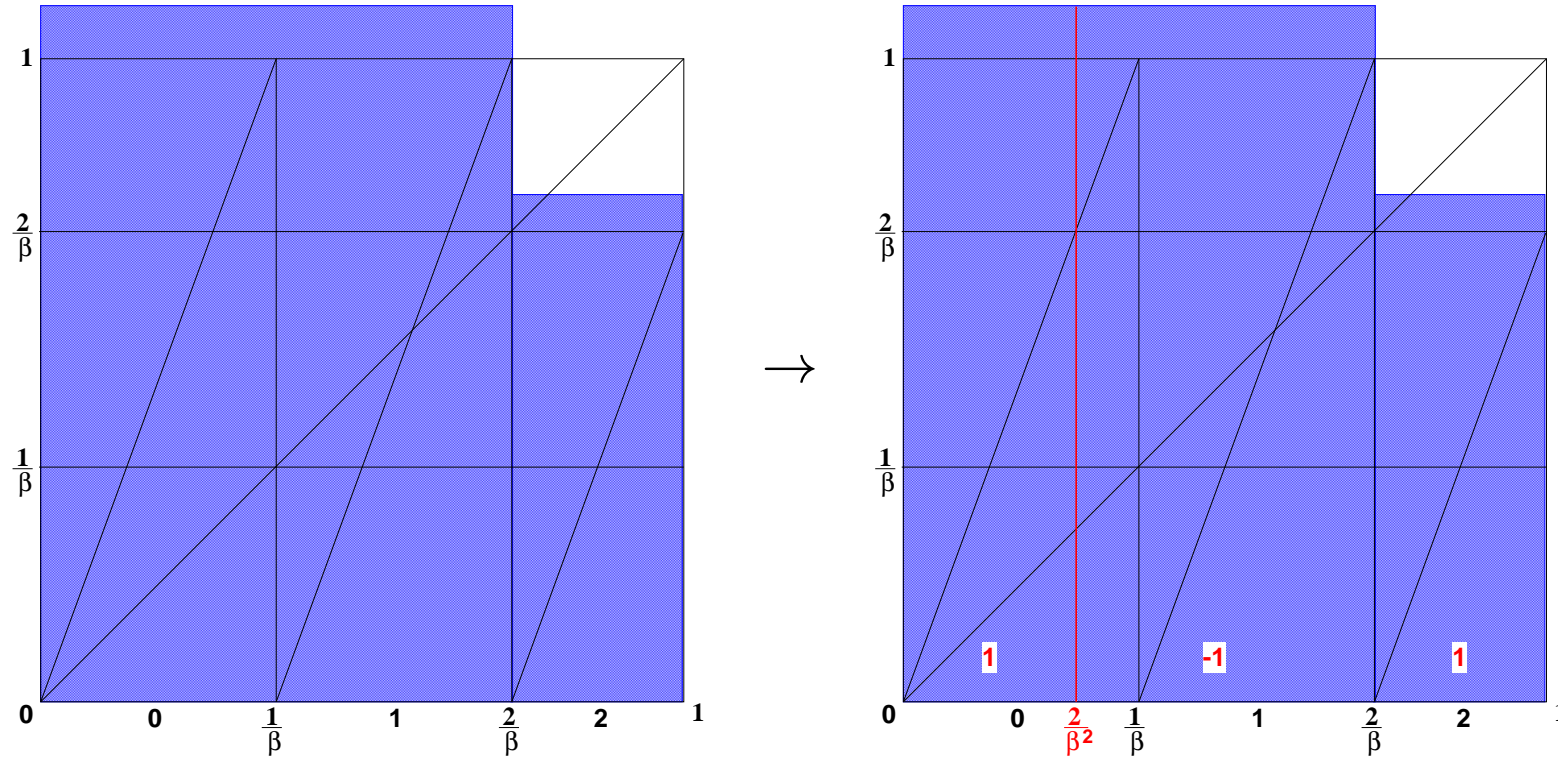
$\{1, -1\}$ -valued Markov chains with

$$\mathbb{E}[Z_n] = 0 \quad \text{and} \quad \mathbb{E}[Z_0 Z_\ell] = (-2 + \sqrt{3})^\ell, \quad \ell \geq 0.$$

We consider the Markov  $\beta$ -transformations with the alphabet  $\Sigma = \{0, 1, 2\}$  and the set  $\mathcal{F} = \{22\}$  of forbidden blocks. Then we have the  $\beta$ -transformation with  $\beta = 1 + \sqrt{3}$ .



# Design of Simple Functions



$$\Theta(x) = \mathbf{1}_{[0, \frac{2}{\beta^2})}(x) - \mathbf{1}_{[\frac{2}{\beta^2}, \frac{2}{\beta})}(x) + \mathbf{1}_{[\frac{2}{\beta}, 1]}(x).$$

$$\mathbb{E}[Z_n] = 0, \text{ (uniform distribution)}$$

and

$$\mathbb{E}[Z_0 Z_\ell] = (-2 + \sqrt{3})^\ell, \quad \ell \geq 0, \text{ negative correlation as desired!}$$

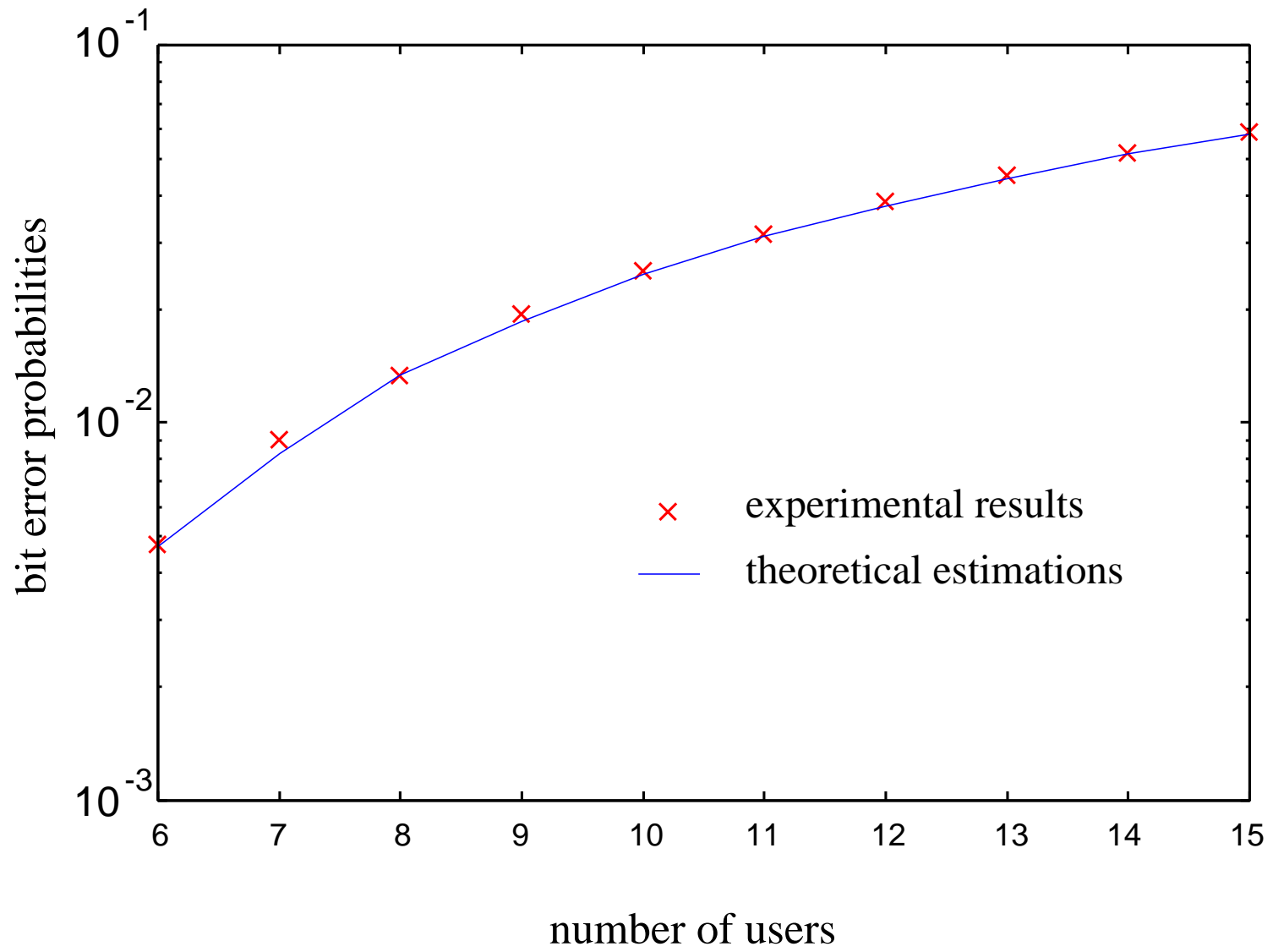
## Experimental Results

$n$	length	# of seq.s	# of seq.s w/ uniform dist.
2	8	12	6
3	20	1728	945

**Example 2** For the order  $n = 3$ , we have

00010020110121021112  $\rightarrow$  11101011001010010001,

where in the right hand side, we use 0 to denote  $-1$  for simplicity.



**Theorem 9** For  $1 \leq \ell \leq n - 1$ , we obtain

$$r_{|\mathcal{B}_n|}(\ell; \mathbf{X}) = (-2 + \sqrt{3})^\ell + \left\{ \left( \frac{\beta}{\bar{\beta}} \right)^\ell - \left( \frac{\bar{\beta}}{\beta} \right)^\ell \right\} \cdot \frac{\left( \frac{\bar{\beta}}{\beta} \right)^n}{1 + \left( \frac{\bar{\beta}}{\beta} \right)^n}.$$

This implies

$$r_{|\mathcal{B}_n|}(\ell; \mathbf{X}) = \mathbb{E}[Z_0 Z_\ell] + O\left(\left(\frac{\bar{\beta}}{\beta}\right)^n\right).$$



## Summary

In this research, we first obtained the topological entropy of the discretized golden mean transformation. We also generalized this result and gave the topological entropy of the discretized Markov  $\beta$ -transformations with the alphabet  $\Sigma = \{0, 1, \dots, k-1\}$  and the set  $\mathcal{F} = \{(k-1)c, \dots, (k-1)(k-1)\}$  ( $1 \leq c \leq k-1$ ) of  $(k-c)$  forbidden blocks.

In view of basic properties of the normalized auto-correlation functions for the de Bruijn sequences that can be regarded as the full-length sequences based on the discretized dyadic transformation, we obtained correlational properties of the full-length sequences based on the discretized golden mean transformation.

We generalized this result and gave the correlational properties of the discretized Markov  $\beta$ -transformations with the alphabet  $\Sigma = \{0, 1, \dots, k-1\}$  and the set  $\mathcal{F} = \{(k-1)(k-1)\}$  of forbidden blocks.

We also applied the generalized result to evaluate the auto-correlation function for the optimum binary spreading sequences of Markov chains based on discretized  $\beta$ -transformations, where  $\beta = 1 + \sqrt{3}$ .