

Percolation and isoperimetric inequalities

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Dynamics on random graphs – 2017/10/25



Percolation

Physical phenomenon:

- (i) Models how fluid can spread through a medium;
- (ii) Models how certain epidemics can spread through a network;
- (iii) Many other motivational examples!

Introduced by *Broadbent and Hammersley* in '57 (independent percolation).

Percolation

Ingredients:

- (i) A graph $G = (V, E)$ (we consider only $|V| = \infty$);
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Two types of percolation: *bond* (edges) and *site* (vertices) percolation.

Today we focus on **SITE** percolation.

Example of Percolation

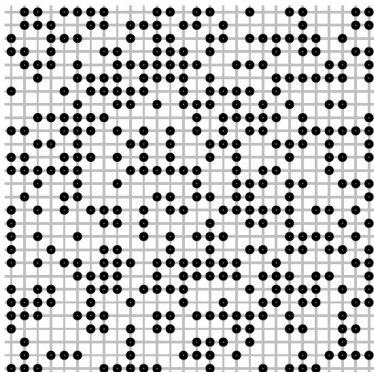


Figure: \mathbb{Z}^2 with $p = 0.5$.

Example of Percolation

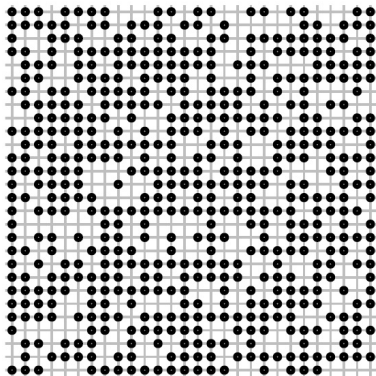


Figure: \mathbb{Z}^2 with $p = 0.7$.

Fundamental questions

Study the connectivity properties of the black (random) subgraph.

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Define

$$\theta(p) := \mathbb{P}_p[\text{vertex } o \text{ is connected to infinity}].$$

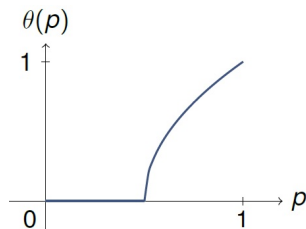


Figure: The function $\theta(p)$.

A critical value

From the previous picture it is then natural to define

$$p_c := \sup\{p \in [0, 1] : \theta(p) = 0\}.$$

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Lots of interesting questions!

- (i) Continuity at p_c ?
- (ii) Behavior of $\theta(p)$ at p_c ?
- (iii) When is p_c non-trivial?

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When is $p_c \in (0, 1)$?

d -dimensional lattices

On \mathbb{Z}^d we know several things, for example:

- If $d \geq 2$ we know that $p_c \in (0, 1)$;
- If $d = 2$ or $d \geq 11$, then we know that $\theta(p_c) = 0$.

We still don't know what happens in the intermediate range of d 's.

Other graphs

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If the degree of the graph G is at most Δ , then $p_c(G) \geq \frac{1}{\Delta} > 0$.

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If the degree of the graph G is at most Δ , then $p_c(G) \geq \frac{1}{\Delta} > 0$.

We do not have such an easy way to investigate upper-bounds for p_c .

The first step in a study of percolation on other graphs [...] will be to prove that the critical probability on these graphs is smaller than one.

Benjamini and Schramm



Other graphs

It is known that $p_c(G) < 1$ when

G has **exponential growth** (Lyons, Benjamini and Schramm, Babson and Benjamini...);

as well as

G is the Cayley graph of the Grigorchuck group, an example of a graph with **intermediate growth** (Muchnik and Pak).

Other graphs

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Vertex-transitive graphs with polynomial growth.

The proof of this fact involves Gromov's theorem (a very difficult and powerful result from **group theory**) and **combinatorial** techniques developed by Babson and Benjamini, and later on simplified by Timar.

Isoperimetric inequalities

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Does the dimension play a role for $p_c(G) < 1$?
How important?

Isoperimetric inequalities

For every finite set $A \subset V(G)$, define the (internal) vertex-boundary as

$$\partial A := \{x \in A : \exists y \in V(G) \setminus A : \{x, y\} \in E(G)\}.$$

Isoperimetric inequalities (dimension)

Define the (isoperimetric) **dimension** of G as follows: we say that $\dim(G) = d > 1$

if and only if d is the largest value for which

there is a constant $c > 0$ such that

$$\inf_{A \subset V(G), A \text{ finite}} \frac{|\partial A|}{|A|^{(d-1)/d}} \geq c.$$

Isoperimetric inequalities (remarks)

Remark: for every $d \geq 2$, \mathbb{Z}^d has isoperimetric dimension d .

Remark: If G has isoperimetric dimension $d > 1$, then we can say that it satisfies IS_d (d -isoperimetric inequality).

Question (Benjamini and Schramm '96)

Is it true that $\dim(G) > 1$ implies that $p_c(G) < 1$?

Some results

- If G is *planar*, has *polynomial growth* and *no accumulation points* then $\dim(G) > 1 \Rightarrow p_c(G) < 1$. [Kozma]

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- If G is *planar*, has *polynomial growth* and *no accumulation points* then $\dim(G) > 1 \Rightarrow p_c(G) < 1$. [Kozma]
- If G satisfies a *stronger* condition than the isoperimetric inequality (called *local isoperimetric inequality*), and has *polynomial growth* then $\dim_\ell(G) > 1 \Rightarrow p_c(G) < 1$. [Teixeira]

Our results

Definition: A measure \mathbb{P} satisfies the *decoupling inequality* $\mathcal{D}(\alpha, c_\alpha)$ (where $\alpha > 0$ is a fix parameter) if for all $r \geq 1$ and any two decreasing events \mathcal{G} and \mathcal{G}' such that

$$\mathcal{G} \in \sigma(Y_z, z \in B(o, r)) \quad \text{and} \quad \mathcal{G}' \in \sigma(Y_w, w \notin B(o, 2r)),$$

we have

$$\mathbb{P}(\mathcal{G} \cap \mathcal{G}') \leq (\mathbb{P}(\mathcal{G}) + c_\alpha r^{-\alpha})\mathbb{P}(\mathcal{G}').$$

In other words: we admit *dependencies*, as long as they *decay fast* enough in the distance.

Our results

With a completely **probabilistic approach** we showed:

Theorem [C. and Teixeira]: If G is transitive, with polynomial growth, $\dim(G) > 1$, and \mathbb{P} satisfies $\mathcal{D}(\alpha, c_\alpha)$ with α “large enough”, then

- (i) There exists a $p_* < 1$, such that if $\inf_{x \in V} \mathbb{P}[Y_x = 1] > p_*$, then the graph contains almost surely a unique infinite open cluster.
- (ii) Moreover, fixed any value $\theta > 0$, we have

$$\lim_{v \rightarrow \infty} v^\theta \mathbb{P}[v < |\mathcal{C}_o| < \infty] = 0,$$

where $\mathcal{C}_o =$ open connected component containing the origin.

Our results

Moreover, in the dependent case we also need to show that:

Theorem [C. and Teixeira]: If G is transitive, with polynomial growth, $\dim(G) > 1$, and \mathbb{P} satisfies $\mathcal{D}(\alpha, c_\alpha)$ with α “large enough”, then

- (i) There exists a $p_{**} > 0$, such that if $\sup_{x \in V} \mathbb{P}[Y_x = 1] < p_{**}$, then the graph contains almost surely **NO** infinite open cluster.
- (ii) Moreover, fixed any value $\theta > 0$, we have

$$\lim_{v \rightarrow \infty} v^\theta \mathbb{P}[v < |\mathcal{C}_o|] = 0,$$

where $\mathcal{C}_o =$ open connected component containing the origin.

Our results (Remark)

We always assume α to be *large enough*.

Although we don't have sharp bounds on its critical value, if α is too small, there are *counterexamples!*

[One counterexample in paper by Tykesson and Windisch]

Our results

Definition: Two metric spaces (X_1, d_1) and (X_2, d_2) are *roughly isometric* (sometimes called “quasi-isometric”) if there is a map $\varphi : X_1 \rightarrow X_2$ s.t.:

(i) There are $A \geq 1, B \geq 0$ such that for all $x, y \in X_1$

$$A^{-1}d_1(x, y) - B \leq d_2(\varphi(x), \varphi(y)) \leq Ad_1(x, y) + B.$$

(ii) There is $C \geq 0$ such that for all $z \in X_2$ there is $x \in X_1$ s.t.

$$d_2(z, \varphi(x)) \leq C.$$

Our results (Remarks)

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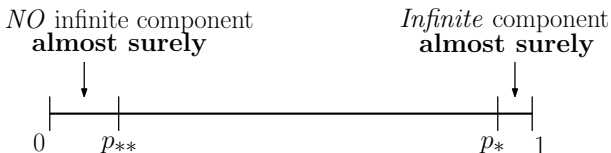
Definition: A graph G is *roughly transitive* if there is a rough isometry between any two vertices of G .

ROUGHLY TRANSITIVE \neq ROUGHLY ISOMETRIC TO TRANSITIVE!
[One counterexample in paper by Elek and Tardos]

Our results

Our proof also works when G is a *roughly transitive graph*:

Theorem [C. and Teixeira]: If G is *roughly-transitive graph*, with polynomial growth, $\dim(G) > 1$, and \mathbb{P} satisfies $\mathcal{D}(\alpha, c_\alpha)$ with α “large enough”, then



and, for every $\theta > 0$,

$$\begin{cases} \lim_{v \rightarrow \infty} v^\theta \mathbb{P}[v < |\mathcal{C}_\infty|] = 0 & \text{if } p < p_{**} \\ \lim_{v \rightarrow \infty} v^\theta \mathbb{P}_p[v < |\mathcal{C}_\infty| < \infty] = 0 & \text{if } p > p_*. \end{cases}$$

Idea of the proof: renormalization (multiscale argument)

On the blackboard.

Main hypothesis I: *polynomial growth*

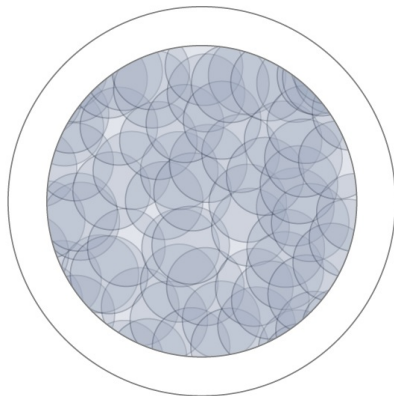


Figure: Polynomial growth allows us to split the graph into cells.



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Main hypothesis II: *Isoperimetric inequality*

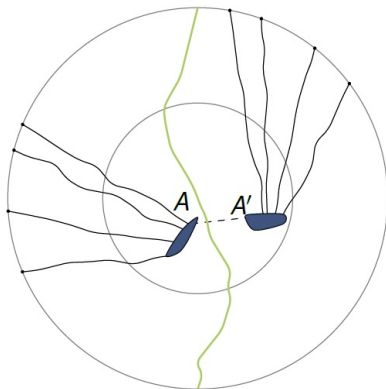


Figure: Isoperimetric inequality implies that there are lots of paths between large connected sets and infinity.

Main hypothesis III: *transitivity* (or rough-transitivity)

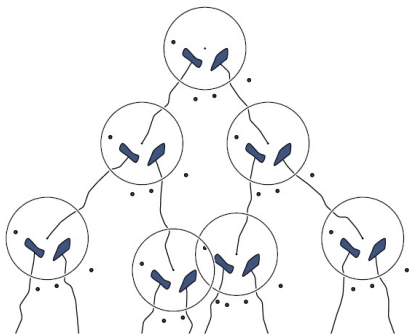


Figure: Transitivity allows us to repeat the same reasoning in different areas of the graph...

Proof

If G satisfies conditions I, II, and III (i.e., polynomial growth, isoperimetric dimension > 1 , rough transitivity), then

assuming $p_c(G) = 1$



it is possible to construct a binary tree inside G .

CONTRADICTION with polynomial growth of G !



Thank you for your attention!