

# Computable absolutely normal numbers and discrepancies

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Prime Numbers and Automatic Sequences: Determinism and Randomness  
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# Introduction

Let  $b \geq 2$  be an integer. Let  $x$  be a real number and denote its base- $b$  expansion by

$$x = [x] + 0.\epsilon_1\epsilon_2\epsilon_3 \dots$$

with digits  $0 \leq \epsilon_j \leq b - 1$ .

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For all  $k$ , for all digits  $d_1, \dots, d_k$ ,

$$\frac{1}{N} \#\{i \leq N : \epsilon_i\epsilon_{i+1} \dots \epsilon_{i+k-1} = d_1 \dots d_k\} \longrightarrow \frac{1}{b^k}$$

as  $N \rightarrow \infty$ .

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$x$  is called *absolutely normal* if it is normal to all integer bases  $b \geq 2$ .

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- ◇ No algebraic number is known to be normal or not normal, nor is any arithmetical constant known to be normal or not normal.
- ◇ There is no easy example of an absolutely normal number.

Let the base  $\beta > 1$  be a real number. Every real  $x$  can be represented in the form

$$\begin{aligned}x &= [x] + 0.\epsilon_1\epsilon_2\epsilon_3\dots \\ &= [x] + \sum_{i=1}^{\infty} \epsilon_i\beta^{-i}\end{aligned}$$

where the digits  $\epsilon_i$  are integers  $0 \leq \epsilon_i < \beta$  that are chosen in increasing order of  $i$  as large as possible.

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### Definition

$x$  is normal to base  $\beta$ , if  $(T_\beta^n(x))_{n \geq 0}$  is uniformly distributed modulo 1 with respect to  $\mu_\beta$ .

- ◇ There are various constructions of normal numbers to a given base  $\beta$  by concatenating strings.



# Continued fractions

Any real number  $x$  has a continued fraction expansion of the form

$$x = [a_0; a_1, a_2, a_3 \dots]$$

where  $a_0 = \lfloor x \rfloor \in \mathbb{Z}$  and  $a_i \in \mathbb{N}$  for  $i \geq 1$ .

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The  $a_i$  are obtained from  $x$  via  $a_i(x) = \lfloor 1/T_G^{i-1}(x) \rfloor$  where

$$T_G : [0, 1) \rightarrow [0, 1), \quad x \mapsto \frac{1}{x} \bmod 1 \text{ if } x > 0, \text{ and } 0 \mapsto 0,$$

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## Definition

$x$  is called *continued fraction normal*, if  $(T_G^n(x))_{n \geq 0}$  is uniformly distributed modulo 1 with respect to  $\mu_G$ .

E.g., concatenate the partial quotients of  $1/2, 1/3, 2/3, 1/4, 2/4, 3/4, \dots$  (Adler/Keane/Smorodinsky)

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There seems to be a trade-off between the time-complexity and convergence to normality ( $D_N \rightarrow 0$ ) of these algorithms.

Becher/Heiber/Slaman: polynomial time, but  $D_N = O(\frac{1}{\log N})$

Sierpinski, Lebesgue, Turing: double-exp. time, but  $D_N = O(\frac{1}{N^\epsilon})$  for some small  $\epsilon$ .

A string of digits  $\omega$  of length  $n$  is called  $(\epsilon, k)$ -normal, if every word  $d$  of length  $k$  appears at least  $n(1 - \epsilon)\mu(d)$  and at most  $n(1 + \epsilon)\mu(d)$  times in  $\omega$ .

# Some preliminaries

A string of digits  $\omega$  of length  $n$  is called  $(\epsilon, k)$ -normal, if every word  $d$  of length  $k$  appears at least  $n(1 - \epsilon)\mu(d)$  and at most  $n(1 + \epsilon)\mu(d)$  times in  $\omega$ .

A real number  $x$  is called *simply normal* to base  $b$ , if all base- $b$  digits of  $x$  appear with the same asymptotic frequency.

Absolute normality is equivalent to simple normality to all bases.

# Sierpinski's construction

Idea:

- ◇ Compute a nested sequence of binary intervals by iteratively halving the previous interval and deciding which of the halves is 'best':
- ◇ do so by successively computing a number in  $[0, 1) \setminus B_N$  for  $N$  increasing with the step of the algorithm where  $B_N$  is an approximation to  $B$ , the set of 'bad' numbers.
- ◇ This will produce a number in  $[0, 1) \setminus B$ , the set of all 'good' numbers.

$B$  is the set of all non- $(\epsilon, k)$ -normal numbers.

Idea:

- ◇ Successively concatenate  $(\epsilon, k)$ -normal words to an increasing set of bases *simultaneously*.
- ◇ Any concatenation of  $(\epsilon, k)$ -normal words, subject to weak conditions, will be normal.
- ◇ Cylinder intervals should be well-behaved.

Need to show explicitly that there is an abundance of  $(\epsilon, k)$ -normal words.

## Theorem (Equipartition)

Let  $T$  be an ergodic transformation on a probability space  $(X, \mathcal{B}, \mu)$ . For any  $\epsilon > 0$  there is an  $n_0 = n_0(\epsilon)$  such that for all  $n \geq n_0$  the set of cylinders of length  $n$  decomposes into two sets  $H$  and  $L$  such that

$$\mu(L) < \epsilon$$

and for any cylinder  $c \in H$

$$\exp(-n(h(T) - \epsilon)) < \mu(c) < \exp(-n(h(T) + \epsilon)).$$

- ◇ Ergodicity: If  $T^{-1}(A) = A$ , then  $\mu(A) \in \{0, 1\}$ .
- ◇  $h(T)$  is the entropy of  $T$ .
- ◇ A cylinder of length  $n$  is a subset of  $[0, 1)$  in which all numbers have the same first  $n$  digits.

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*There is an algorithm that uses only elementary operations to compute an absolutely Pisot normal number.*

Proof:

- ◇ Becher/Heiber/Slaman-construction of choosing simultaneously to several bases long 'good' blocks of digits.
- ◇ Explicit estimate for  $(\epsilon, k)$ -normal numbers.

Levin has also given such a construction using exponential sums estimates.

For  $\beta$ -expansions:

- ◇ The digits are independent, if one looks at two digits that are far enough away from each other.
- ◇ Admissible sequences are exactly sequences all of whose shifts are lexicographically strictly less than the *modified expansion* of 1.
- ◇ Inserting enough zeros in-between two words gives again an admissible word. 'Enough' can be made explicit and depends only on the base, not on the blocks.

# Non- $(\epsilon, k)$ -normal numbers for $\beta$ -expansions

Fix a string  $d$  of  $k$  digits in base  $\beta$ , let  $X_i(x) = 1 - \mu_\beta(c(d))$  if  $d$  appears in  $x$  at  $i$ -th position, and  $X_i(x) = -\mu_\beta(c(d))$  if not.

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## Theorem (A. Siegel)

*Let  $X = X_1 + X_2 + \dots + X_k$  be the sum of  $k$  possibly dependent random variables. Suppose that  $X_i$ , for  $i = 1, 2, \dots, k$ , is the sum of  $n_i$  mutually independent random variables having values in the interval  $[0, 1]$ . Let  $E[X_i] = n_i p_i$ . Then for  $a \geq 0$*

$$P(X - E[X] \geq a) < \text{two explicit exponential terms } (a, k, p_i, n_i)$$

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## Lemma (Madritsch, S., Tichy)

Let  $\beta$  be a Pisot number. The  $\mu_\beta$ -measure of the set of not  $(\epsilon, k)$ -normal words of length  $n$  satisfies

$$\mu_\beta(B(n, \epsilon, k)) \leq 4\beta^k \beta^{-\eta n}$$

for  $n \geq M + k$  with explicit  $\eta > 0$  and  $M$ .

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Has already been greatly improved some weeks ago by Becher and Yuhjtman, who could adapt Becher, Heiber and Slaman's polynomial algorithm.

In particular, they used that  $q_n(x) \sim e^{nL}$  a.e., where  $L = \pi^2/(12 \log(2))$ .

# Non- $(\epsilon, k)$ -normal numbers for continued fraction expansions

Fix a string  $d_1 \dots d_k$  of  $k$  positive integers.

Let  $X_i(x)$  be the normalized random variable that counts whether or not the string  $d_1 \dots d_k$  appears in the continued fraction expansion of  $x$  at position  $i$ .

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The  $X_i$  are not independent, but they satisfy a mixing property that can be derived from mixing properties of the partial quotients  $a_i$ :

W. Philipp:

$|\mu_G(A \cap B) - \mu_G(A)\mu_G(B)| \leq \rho^n \mu_G(A)\mu_G(B)$  for all  
 $A \in \sigma(a_1, \dots, a_k)$ ,  $B \in \sigma(a_{k+n}, a_{k+n+1}, \dots)$  for some  $0 < \rho < 0.8$ .

## Theorem (F. Merlevède, M. Peligrad, E. Rio)

Let  $(X_i)_{i \geq 1}$  be a sequence of centered real-valued random variables bounded by a uniform constant  $M$  and with  $\alpha_n(X_i)$  satisfying  $\alpha_n \leq \exp(-2nc)$  for some positive  $c$ . Then there is an explicit positive constant  $C$  depending only on  $c$  such that for all  $n \geq 4$  and  $x \geq 0$

$$\mathbb{P}(|X_1 + \dots + X_n| \geq x) \leq \exp\left(-\frac{Cx^2}{nM^2 + Mx(\log n)(\log \log n)}\right).$$

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## Lemma

$$\mu_G(B_{CF}(\epsilon, d_1 \dots d_k, n)) \leq \exp\left(-\eta_{CF}(\epsilon, d_1 \dots d_k) \frac{n}{\log n}\right)$$

for some explicit  $\eta_{CF}(\epsilon, \vec{d}) > 0$  and all  $n \geq n_0 = 2(k+1)$ .

This was only known explicitly with linear decay by Vandehey (not enough for this application).

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# Normal numbers and discrepancies

A sequence  $(x_n)_{n \geq 0}$  of real numbers is called *uniformly distributed modulo 1* if for any interval  $I \subseteq [0, 1)$ ,

$$\frac{1}{N} \#\{n \leq N : x_n \bmod 1 \in I\} \longrightarrow \lambda(I)$$

as  $N \rightarrow \infty$ .



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Wall (1949):

$x$  is normal to base  $b$  if and only if  $(b^n x)_{n \geq 0}$  is uniformly distributed modulo 1.

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The quantity

$$D_N(x_n) = \sup_{I \subseteq [0,1)} \left| \frac{1}{N} \#\{n \leq N : x_n \bmod 1 \in I\} - \lambda(I) \right|$$

is called the *discrepancy* of  $(x_n)$ .

# Results on discrepancy

- For  $N$  fixed:  $\frac{1}{N} \leq D_N(x_n) \leq 1$ .

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- Philipp (1974/75), Fukuyama (2013):  
for almost all  $x$ :  $\limsup_{N \rightarrow \infty} \frac{D_N(b^n x) N^{1/2}}{(\log \log N)^{1/2}} = c(b) > 0$ .

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- Levin (1979, 1999):  
For  $b \geq 2$ , he constructed  $x$  with  $D_N(b^n x) = O\left(\frac{(\log N)^2}{N}\right)$ ,  
and  $x$  such that  $D_N(b^n x) = O_b\left(\frac{(\log N)^3}{N^{1/2}}\right)$  for all  $b \geq 2$ .

# Absolutely normal numbers of low discrepancy

## Theorem (with Verónica Becher and Theodore Slaman)

*There is an algorithm that uses only elementary operations to compute an absolutely normal number  $x$  in triple-exp. time, but with  $D_N(b^n x) = O\left(\frac{\sqrt{\log \log N}}{\sqrt{N}}\right)$ , for all  $b \geq 2$ .*



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Proof:

- ◇ Combine Sierpinski's construction with estimates by Philipp on the discrepancy of sequences of the form  $(b^n x)_{n \geq 0}$ .

Thank you for your attention!