

# Localization at $\omega$ -compact types, as sequential colimits.

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**Localization is a process  
of adding inverses  
to an algebraic structure  
in a universal way.**

A type  $X$  is said to be  **$A$ -local** if any map  $A \rightarrow X$  has a unique extension to a map  $\mathbf{1} \rightarrow X$ .

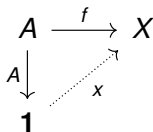
More precisely,  $X$  is  $A$ -local if the map

$$\lambda(x : \mathbf{1} \rightarrow X). x \circ A : (\mathbf{1} \rightarrow X) \rightarrow (A \rightarrow X)$$

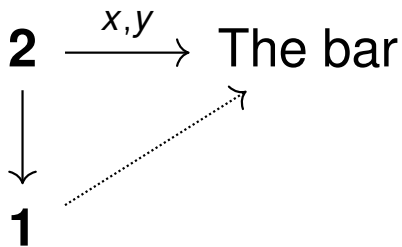
is an equivalence.

Example:

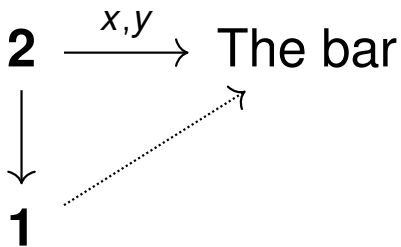
- ▶ The unit type  $\mathbf{1}$  is  $A$ -local for any  $A$ .
- ▶ The mere propositions are  $\mathbf{2}$ -local.



Two terms walk into a **2**-local bar...

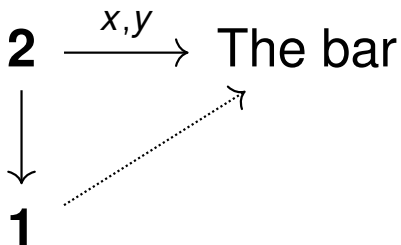


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and they order one beer..

Two terms walk into a 2-local bar...



and they order one beer..

**...Chuck Norris drinks it.**

A **reflective subuniverse** consists of a subuniverse  $P : \mathbf{U} \rightarrow \mathbf{Prop}$  of  **$P$ -types** such that for each type  $X$

- ▶ there is a type  $\circ(X)$  for which  $P(\circ(X))$  holds,
- ▶ there is a map  $\eta_X : X \rightarrow \circ(X)$

such that the map

$$\lambda f. f \circ \eta_X : (\circ(X) \rightarrow Y) \rightarrow (X \rightarrow Y)$$

is an equivalence for each  $P$ -type  $Y$ .

The **sequential colimit**, denoted by either  $\text{colim}_n(A_n)$  or  $A_\infty$ , of a type sequence has constructors

$$i : \prod_{(n:\mathbb{N})} A_n \rightarrow A_\infty$$

$$j : \prod_{(n:\mathbb{N})} \prod_{(x:A_n)} i_n(x) = i_{n+1}(a_n(x)).$$

with the expected universal property.



We say that  $X$  is  **$\omega$ -compact** if the canonical function of type

$$\operatorname{colim}_n (X \rightarrow A_n) \rightarrow (X \rightarrow A_\infty)$$

is an equivalence for *every* type sequence  $(A_n, a_n)_{n:\mathbb{N}}$ .

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Examples:

- ▶ Finite types are  $\omega$ -compact.  $\mathbb{N}$  is not.

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Examples:

- ▶ Finite types are  $\omega$ -compact.  $\mathbb{N}$  is not.
- ▶ If  $B$  is  $\omega$ -compact, and  $E(x, y)$  is  $\omega$ -compact for each  $x, y : B$ , then the higher inductive type  $X$  with constructors

$$b : B \rightarrow X$$

$$e : \prod_{(x, y : B)} E(x, y) \rightarrow (b(x) = b(y))$$

is  $\omega$ -compact.

# The circle is $\omega$ -compact

The **circle** is a higher inductive type  $\mathbb{S}^1$  with

$$\text{base} : \mathbb{S}^1$$
$$\text{loop} : \text{base} =_{\mathbb{S}^1} \text{base}$$

**recursion principle of the circle:** for any type  $X$ , a function of type  $\mathbb{S}^1 \rightarrow X$  is determined by

$$x : X$$
$$p : x =_X x.$$

## Each $n$ -sphere is $\omega$ -compact

The  $(n + 1)$ -**sphere** is a higher inductive type  $\mathbb{S}^{n+1}$  with

$$N : \mathbb{S}^{n+1}$$

$$S : \mathbb{S}^{n+1}$$

$$\text{equator} : \mathbb{S}^n \rightarrow (N =_{\mathbb{S}^{n+1}} S).$$

**recursion principle of the  $(n + 1)$ -sphere**: for any type  $X$ , a function of type  $\mathbb{S}^{n+1} \rightarrow X$  is determined by

$$x : X$$

$$y : X$$

$$e : \mathbb{S}^n \rightarrow x =_X y$$

$$\begin{array}{ccc} A & \xrightarrow{f} & X \\ \downarrow & & \downarrow \ell \\ \mathbf{1} & \xrightarrow{j(f)} & L_A(X) \end{array}$$

$$\begin{array}{ccc} A & \xrightarrow{f} & X \\ \downarrow & \nearrow x & \downarrow \ell \\ \mathbf{1} & \xrightarrow{j(f)} & L_A(X) \end{array}$$

$$\begin{array}{ccc}
 A & \xrightarrow{f} & X \\
 \downarrow & \nearrow x & \downarrow \ell \\
 \mathbf{1} & \xrightarrow{j(f)} & L_A(X)
 \end{array}$$

$$\ell : X \rightarrow L_A(X)$$

$$j : (A \rightarrow X) \rightarrow L_A(X)$$

$$J : \prod_{(f:A \rightarrow X)} \prod_{(a:A)} j(f) = \ell(f(a))$$

$$k : \prod_{(f:A \rightarrow X)} \prod_{(x:X)} \prod_{(H:\prod_{(a:A)} x=f(a))} \ell(x) = j(f)$$

$$K : \prod_{(f:A \rightarrow X)} \prod_{(x:X)} \prod_{(H:\prod_{(a:A)} x=f(a))} k(f, x, H)^{-1} \cdot \ell(H(a)) = J(f, a).$$



# Outline of the formalization project with Floris

**We shall give a formal proof (in LEAN) that for any type  $X$  and any  $\omega$ -compact type  $A$ , the sequential colimit  $L_A^\infty(X)$  is the localization of  $X$  at  $A$ , i.e. that  $L_A^\infty(X)$  is the  $A$ -local reflection of  $X$ .**

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**We shall give a formal proof (in LEAN) that for any type  $X$  and any  $\omega$ -compact type  $A$ , the sequential colimit  $L_A^\infty(X)$  is the localization of  $X$  at  $A$ , i.e. that  $L_A^\infty(X)$  is the  $A$ -local reflection of  $X$ .**

- ▶ In particular, we need to show that  $L_A^\infty(X)$  is  $A$ -local, i.e. that maps of type  $A \rightarrow L_A^\infty(X)$  extend uniquely to maps of type  $\mathbf{1} \rightarrow L_A^\infty(X)$ .

$$\begin{array}{ccc} A & \xrightarrow{f} & L_A^\infty(X) \\ A \downarrow & \nearrow & \\ \mathbf{1} & & \end{array}$$

- ▶ The  $\omega$ -compactness comes up when one starts manipulating the type  $A \rightarrow L_A^\infty(X)$ .

Fact: Let  $(P_n, p_n)$  be a sequence in which each  $P_n$  is a mere proposition. Then we have the equivalence

$$\operatorname{colim}(P_n) \simeq \exists_{(n:\mathbb{N})} P_n.$$

### Theorem (Independence of premises)

*If every mere proposition is  $\omega$ -compact, then for any type  $X$ , and any sequence  $(P_n, p_n)$  of mere propositions, one has*

$$(X \rightarrow \exists_{(n:\mathbb{N})} P_n) \leftrightarrow (\exists_{(n:\mathbb{N})} X \rightarrow P_n)$$

**Question: The condition that all mere propositions are  $\omega$ -compact seems to be non-constructive. Is it? How does it relate to other non-constructive principles?**