

On arithmetic index in the Thue-Morse word

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- Let $\Sigma_t = \{0, 1, \dots, t - 1\}$ be a finite alphabet
A sequence $w = w_0w_1w_2w_3 \dots$, $w_i \in \Sigma_t$ is an infinite word over Σ_t
- A sequence u is *periodic* with the length of a period T ,
if for every i the following holds: $u_i = u_{i+T}$
- An *arithmetic subsequence* of length k with the starting symbol w_c
and difference d in w is a word $w_d^c = w_cw_{c+d}w_{c+2d} \dots w_{c+(k-1)d}$
- If $w_{c+id} = a$ for $i = 0, 1, \dots, k - 1$ and $a \in \Sigma_t$,
then w_d^c is called an *arithmetic progression*.

- An *arithmetic closure* of an infinite word w over Σ_t is a set $A_w = \{w_d^c \mid c \geq 0, d > 0\}$
- An *arithmetic complexity* of w is a function $a_w(n) = |A_w \cap \Sigma_t^n|$
- A word w is *arithmetic universal*, if $a_w(n) = t^n$

Theorem [Van der Waerden, 1927]

An arithmetic closure of every infinite word over a finite alphabet contains an arbitrary long arithmetic progression.

Object of the research

Let $\Sigma = \{0, 1\}$, $x \in \mathbb{N}$

- $B(x) = x_{n-1} \dots x_1 x_0$, where $x_i \in \Sigma$, $x = \sum_{i=0}^{n-1} x_i \cdot 2^i$ is a *binary expansion* of x
- $p(x) = \sum_{i=0}^{n-1} x_i \bmod 2$

The *Thue-Morse word* is a sequence

$w_{TM} = w_0 w_1 w_2 w_3 \dots$, where $w_i = p(i)$.

$$w_{TM} = 01101001100101101001011001101001 \dots$$

Words with $T = 1$

Let $\Sigma = \{0, 1\}$, $w_{TM} = w_0 w_1 w_2 w_3 \dots$ is the Thue-Morse word.

- Define $L(c, d)$ as the length of an arithmetic progression in w_{TM} with the starting symbol w_c and the difference d for $n, d \in \mathbb{N}$
- $L(d) = \max_c L(c, d)$

Theorem [P., 2015]

For every $n \in \mathbb{N}$, $n > 1$ the following holds:

$$\max_{d < 2^n} L(d) = \begin{cases} 2^n + 4, & n \text{ is even,} \\ 2^n, & n \text{ is odd.} \end{cases}$$

Moreover, the maximum is reached with the difference $d = 2^n - 1$ in both cases.

Words with $T > 1$

Consider a word $u = (01)^k, k \in \mathbb{N}$

$u \in A_{TM}$

[Avgustinovich, Frid, Fon-der-Flaas, 2000]

A word w_{TM} is arithmetic universal.

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Questions:

- What is an upper bound on the value of the difference while searching arithmetic subwords in the Thue-Morse word?
- Which words occur in the Thue-Morse word with the maximal difference?

Object of the research

- Let w be an infinite word over Σ_t .
- Let $d \in \mathbb{N}$, then define $A_w(d)$ as a set of arithmetic subsequences with the difference d in w .
- $A_w = \bigcup_{d=1}^{\infty} A_w(d)$
- A finite word u has an *arithmetic index* $l_w(u)$, if $l_w(u)$ is the length of a binary expansion of $i(u) = \min\{d \mid u \in A_w(d)\}$
- If u does not belong to A_w , then $l_w(u) = \infty$
- The object of a research is a function $\max_{u:|u|=n} l_w(u)$

Words with $T = 1$

- Define $L(c, d)$ to be the length of an arithmetic progression in w_{TM} with the starting symbol w_c and the difference d for $c, d \in \mathbb{N}$
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Corollary

Let u be of the form 0^N or 1^N for some $N \in \mathbb{N}$.

Then $I_{TM}(u) = n = \lceil \log N \rceil$.

Upper bound

Let u be a binary word of length N , $2^{n-1} \leq N < 2^n$.

- Define a "basis": b_1, b_2, \dots, b_N ,

$$b_1 = 1\alpha_1\alpha_2 \cdots \alpha_{N-1}$$

$$b_2 = 01\alpha_1\alpha_2 \cdots \alpha_{N-2}$$

.....

$$b_i = 0^{i-1}1\alpha_1\alpha_2 \cdots \alpha_{N-i}$$

.....

$$b_N = 0^{N-1}1, \text{ where } \alpha_i \in \{0, 1\}$$

- Every b_i can be obtained with the difference $d = 2^n - 1$
- A word u can be represented in the following way:

$$u = \sum_{i=1}^n \beta_i \cdot b_i, \beta_i \in \{0, 1\}$$

Upper bound

- Basis: b_1, b_2, \dots, b_N ,
 $b_i = 0^{i-1} 1 \alpha_1 \alpha_2 \dots \alpha_{n-i}, \alpha_i \in \{0, 1\}$
- $u = b_1 + b_2 + \dots + b_N$
- Set C_i to be an index of the first symbol of b_i arithmetic occurrence in the Thue-Morse word, c_i to be a binary expansion of C_i , the length of each c_i is equal to $3n = 3 \cdot \lceil \log N \rceil$
- Build a binary number $c = c_1 c_2 \dots c_N$ of length $3n \cdot N$
- Then take a difference of the form

$$D = \underbrace{0 \dots 0}_{2n} \underbrace{1 \dots 1}_n \underbrace{0 \dots 0}_{2n} \underbrace{1 \dots 1}_n \dots \underbrace{0 \dots 0}_{2n} \underbrace{1 \dots 1}_n$$
- $\max_{u: |u|=N} I_w(u) \leq 3Nn = 3N \cdot \lceil \log N \rceil$

Upper bound. Example

Set $u = 1110$, $N = 4 = 2^2$, $n = 2$

Define a basis:

- $d = 2^2 - 1 = 3$

$$\mathbf{w}_{\text{TM}} = \mathbf{01101001100101101001011001101001} \dots$$

$$b_1 = 1000, C_1 = 21, c_1 = 010101$$

$$b_2 = 0100, C_2 = 18, c_2 = 010010$$

$$b_3 = 0010, C_3 = 15, c_3 = 001111$$

$$b_4 = 0001, C_4 = 12, c_4 = 001100$$

- $u = b_1 + b_2 + b_3$

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- $u = b_1 + b_2 + b_3$

- The initial number:

$$c = c_1 c_2 c_3 = 010101'010010'001111$$

- The difference D :

$$D = 000011'000011'000011$$

- We can obtain u as an arithmetic subword starting with w_C with the difference D :

Upper bound. Example

	c+iD	wt	u
	010101'010010'001111	9	1
\oplus	000011'000011'000011		
	<u>011000'010101'010010</u>	7	1
\oplus	000011'000011'000011		
	<u>011011'011000'010101</u>	9	1
\oplus	000011'000011'000011		
	<u>011110'011011'011000</u>	13	0

Lower bound

A *factor complexity* $p_w(N)$ of an infinite word w is the number of its factors of length N .

Theorem [Avgustinovich, 1994]

A factor complexity of the Thue-Morse word is

$$p_{w_{TM}}(N+1) = 3N + \rho(N),$$

where $\rho(N) = \min\{N - 2^k, 2^k + 1 - N\}$ and $n/2 < 2^k \leq N$.

- $|A_{TM}(1) \cap \Sigma^N| = p_{w_{TM}}(N) < 4N$

Lower bound

- Consider $A(d, N) := |A_{TM}(d) \cap \Sigma^N|$
- $A(1, N) \leq p_{w_{TM}}(N) < 4N$
- $A(3, N) \leq p_{w_{TM}}(3N) < 12N$
-
- $A(d, N) \leq p_{w_{TM}}(d \cdot N) < 4d \cdot N$

The aim is to obtain a lower bound for x s.t. $2^N \leq \sum_{d=1}^x A(d, N)$.

$$x \geq 0.5(N - 2 - \log N)$$

Resume

A word u has an arithmetic index $I_{TM}(u)$ in the Thue-Morse word, if $I_{TM}(u)$ is the length of binary expansion of $i_{TM}(u) = \min\{d | u \in A_{TM}(d)\}$

$$0.5(N - 2 - \log N) < \max_{u: |u|=N} I_w(u) \leq 3N \lceil \log N \rceil$$

Thank you for your attention!