

Avoiding k -abelian powers in words

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Two words u and v are **k -abelian equivalent** if:
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uv is a **k -abelian-square** if u and v are k -abelian equivalent
 uvw is a **k -abelian-cube** if u, v, w are k -abelian equivalent

$u_1 u_2 \dots u_n$ is a **k -abelian- n th-power** if u_1, \dots, u_n are k -abelian equivalent

“ ∞ -abelian-equivalence” = word equality



1-abelian-equivalence = abelian-equivalence

“ ∞ -abelian-equivalence” = word equality



⋮



3-abelian-equivalence



2-abelian-equivalence



1-abelian-equivalence = abelian-equivalence

What about avoidability of powers ?

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Avoidability of k -abelian-powers

Question (Huova, Karhumäki, Saarela, Saari 2011)

- *Is there a k such that k -abelian-squares avoidable on a ternary alphabet?*
- *Is there a k such that k -abelian-cubes avoidable on a binary alphabet?*

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Theorem (Huova, Karhumäki, Saarela, Saari 2011)

2-abelian-squares are not avoidable on a ternary alphabet.

The longest 2-abelian-square-free ternary word has size 537.

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Theorem (R. 2013)

2-abelian-cubes are avoidable over 2 letters.

3-abelian-squares are avoidable over 3 letters.

2-abelian-cube-free binary word

Let $h : A_3^* \rightarrow A_2^*$ be the following 47-uniform morphism.

$$h : \begin{cases} 0 \rightarrow 00100101001011001001010010011001001100101101011 \\ 1 \rightarrow 00100110010011001101100110110010011001101101011 \\ 2 \rightarrow 00110110101101001011010110100101001001101101011 \end{cases}$$

Theorem (R. 2013)

For every abelian-cube-free word $w \in A_3^$, $h(w)$ is 2-abelian-cube-free.*

\Rightarrow 2-abelian-cubes are avoidable on a binary alphabet

3-abelian-square-free ternary word

Let $h : A_4^* \rightarrow A_3^*$ be the following 25-uniform morphism.

$$h : \begin{cases} 0 \rightarrow 0102012021012010201210212 \\ 1 \rightarrow 0102101201021201210120212 \\ 2 \rightarrow 0102101210212021020120212 \\ 3 \rightarrow 0121020120210201210120212 \end{cases}$$

Theorem (R. 2013)

For every abelian-square-free word $w \in A_4^$, $h(w)$ is 3-abelian-square-free.*

\Rightarrow 3-abelian-squares are avoidable on a ternary alphabet

Idea of the proof (1/2)

- Sufficient conditions for a morphism h to be k -abelian- n th-power-free
- i.e. for every abelian- n th-power-free word w , $h(w)$ is k -abelian- n th-power-free

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- Similar to [Carpi 1993] (sufficient conditions for a morphism to be abelian- n th-power-free)
- 1st condition: the prefixes (resp. suffixes) of length $k - 1$ in the images of h are the same.
- 2nd condition: 'generalized' Parikh matrix M of h has full rank:

$$\forall x \in \Sigma, w \in \Sigma^k$$

$$M[x, w] = |h(a)p|_w$$

where p is the prefix of $h(x)$ of length $k - 1$

- Suppose that $h(w)$ has a k -abelian- n th-power

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- Check that for every possibility:
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(e.g. $M'^{-1}V$ is not an integer vector)
 - either the pre-image in w has an abelian- n th-power
- If yes : h is k -abelian- n th-power-free

Avoidability of k -abelian- n th-powers over ℓ letters (sum.)

“ ∞ -abelian-equivalence” = word equality

$\ell \setminus n$	2	3
2	No	Yes (Thue 1906)
3	Yes (Thue 1906)	Yes

k -abelian-equivalence, $3 \leq k$

$\ell \setminus n$	2	3
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2-abelian-equivalence

$\ell \setminus n$	2	3
2	No	Yes (R. 2013)
3	No (Huova, Karhumäki, Saarela, Saari 2011)	Yes
4	Yes	Yes

1-abelian-equivalence = abelian-equivalence

$\ell \setminus n$	2	3	4
2	No	No	Yes (Dekking 79)
3	No	Yes (Dekking 79)	Yes
4	Yes (Keränen 92)	Yes	Yes

What about long k -abelian-repetitions?

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Negative answer to the first Mäkelä's question

Question (Mäkelä 2003)

Can you avoid abelian-cubes of the form uvw where $|u| \geq 2$, over two letters? - You can do this at least for words of length 250.

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Proof:

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Proof:

- Exhaustive search.
- One can restrict w.l.o.g. on Lyndon words.
- Largest Lyndon word: 290.

Avoidability of long abelian repetitions

Weak version of Mäkelä's questions:

Question 1

Can we avoid long abelian-cubes over two letters?

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⇒ Yes ! One can avoid abelian squares uv with $|u| \geq 6$

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Weak version of Mäkelä's questions:

Question 1

Can we avoid long abelian-cubes over two letters?

⇒ Still open... but period at least 3

Question 2

Can we avoid long abelian-squares over three letters?

⇒ Yes ! One can avoid abelian squares uv with $|u| \geq 6$

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Can we avoid long abelian-squares over three letters?

$$\text{Let } h_6 : \begin{cases} a \rightarrow ace \\ b \rightarrow adf \\ c \rightarrow bdf \\ d \rightarrow bdc \\ e \rightarrow afe \\ f \rightarrow bce \end{cases} \quad \text{and } \varphi : \begin{cases} a \rightarrow bbbaabaaac \\ b \rightarrow bccacccbcc \\ c \rightarrow ccccbbbcbc \\ d \rightarrow ccccccccaa \\ e \rightarrow bbbbbcabaa \\ f \rightarrow aaaaaabaa. \end{cases}$$

Theorem (R., Rosenfeld 2015)

The sequence obtained by applying φ to the fixed-point of h_6 , $\varphi(h_6^\infty(a))$, does not contain any abelian-square of period more than 5.

What about long k -abelian-repetitions?

Theorem (R., Rosenfeld 2014 & 2015)

For every $k \geq 2$, one can avoid long k -abelian-squares on binary words.

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For every $k \geq 2$, one can avoid long k -abelian-squares on binary words.

Let $g(k)$ be the minimal number of k -abelian-squares in an infinite binary word

Note: $g(1) = \infty$ [Entringer, Jackson & Schatz 1974]

$g(\infty) = 3$ [Fraenkel & Simpson 1995]

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Theorem (R., Rosenfeld 2014 & 2015)

- $5 \leq g(2) \leq 707$ *period* ≤ 60
- $g(3) = g(4) = 4$ $0^2, 1^2, (01)^2, (10)^2$
- $g(k) = 3$ for $k \geq 5$ $0^2, 1^2, (01)^2$

$k = 2$

$$\varphi : \begin{cases} a \rightarrow bbbaabaaac \\ b \rightarrow bccaccbcc \\ c \rightarrow ccccbbcbcb \\ d \rightarrow cccccccaa \\ e \rightarrow bbbbbcabaa \\ f \rightarrow aaaaaabaa. \end{cases} \quad h_2 : \begin{cases} a \rightarrow 11100000000 \\ b \rightarrow 11010001010 \\ c \rightarrow 11111101010. \end{cases}$$

$h_2(\varphi(h_6^\infty(a)))$ has no 2-abelian squares of period more than 60

$k = 3$ and $k = 4$

$$h_3 : \begin{cases} 0 \rightarrow 00011001010011101011000101010001011101011000101 \\ 1 \rightarrow 00011001010011101011001110101011100011101011000101 \\ 2 \rightarrow 0001100101001110001010001100101100011101011000101 \\ 3 \rightarrow 000110010100111001010100111000101100101011000101. \end{cases}$$

For every abelian-square-free word w , $h_3(w)$ contains only 4 distinct 3-abelian-squares: 0^2 , 1^2 , $(01)^2$ and $(10)^2$

Open questions : abelian equivalence

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Question

What is the minimal p such that one can avoid abelian-squares of period at least p over three letters?

$$2 \leq p \leq 6$$

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What is the minimal s such that one can construct infinite ternary word with only s abelian-squares ?

$$3 \leq s \leq 22$$

Open questions : abelian equivalence

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Question

Can we avoid long abelian-cubes over two letters?

If yes:

- What is the minimum period ?
- What is the minimum number of cubes ?

Open questions : 2-abelian equivalence

Question

What is the minimal p such that one can avoid 2-abelian-squares of period at least p over two letters?

$$2 \leq k \leq 60$$

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$$2 \leq k \leq 60$$

Question

What is $g(2)$? (i.e. smallest s such that one can construct an infinite binary word with only s 2-abelian-squares)

$$5 \leq g(2) \leq 707$$

- Fractional k -abelian-repetitions ?
- k -abelian-repetitions patterns ?
- ...